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A Study of Roots of a Certain Class of Counting Polynomials

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Abstract: Suppose \mathcal{G} is a graph with vertex-orbits $\mathbb{O}_1, \mathbb{O}_2, \dots, \mathbb{O}_t$, and $|\mathbb{O}_i|$ denotes the cardinality of \mathbb{O}_i . Then $\mathbb{O}_{\mathcal{G}}(x) = \sum_{i=1}^t x^{|\mathbb{O}_i|}$ is called as orbit polynomial. It is well-known that this polynomial has a unique positive zero δ in the interval $[0, 1]$. The aim of this paper is to study the specific properties of this polynomial and then we determine the location of this root for several classes of complex networks to compare with other graphical measures. Besides, we compare the unique positive zero measure with several well-known centrality graph measures.

Keywords: orbit polynomial; complex networks; automorphism group; entropy

1. Introduction

The orbit polynomial of a graph is a polynomial that encodes information about the orbits of its automorphism group. Its roots and their properties have become an important tool for studying the symmetry structure and other structural features of networks in many different fields, including graph theory, physics, chemistry, and computer science.

The unique positive root of the orbit polynomial, denoted δ , has gained particular attention due to its ability to serve as a structural descriptor for networks. δ has been shown to provide insight into the symmetry structure of a graph, with its value being closely related to the size and complexity of its automorphism group. This line of research has given rise to a range of applications of δ in network analysis, including the classification and comparison of different types of networks.

Overall, the study of the roots of the orbit polynomial of networks has opened up new avenues for research into the structural properties of networks, with many potential applications in various fields. Ongoing efforts are addressing open questions related to the computation, properties, and applications of the orbit polynomial and its roots, with the aim of furthering our understanding of the structure of networks and their applications in many different areas of science.

Recently, the concept of orbit polynomial [13] for a connected network \mathcal{N} was defined as follows:

$$\mathbb{O}_{\mathcal{G}}(x) = \sum_{i=1}^t x^{|\mathbb{O}_i|},$$

where $\mathbb{O}_1, \mathbb{O}_2, \dots, \mathbb{O}_t$ are all vertex-orbits of \mathcal{N} and the cardinality of an orbit is denoted by $|\mathbb{O}_i|$. Moreover, the modified version of this polynomial or M -orbit polynomial is defined as

$$\mathbb{O}_{\mathcal{N}}^*(x) = 1 - \sum_{i=1}^t x^{|\mathbb{O}_i|}. \quad (1)$$

The definition of this polynomial is based on two concepts: the automorphism group and the vertex-orbits. It is well-known that the orbit polynomial has a unique positive zero δ in the interval $[0, 1]$, see [13]

Ghorbani et al. in [20] generalized the concept of orbit polynomial as follows: Consider the action of permutation group Γ acting on the set X . Then the generalized orbit polynomial is $\mathbb{O}_\Gamma(x) = \sum_{i=1}^t c_i x^{|\mathbb{O}_i|}$, where $\mathbb{O}_1, \dots, \mathbb{O}_t$ are all orbits of Γ . It is clear that for a graph \mathcal{N} with automorphism group $\Gamma = \text{Aut}(\mathcal{N})$, \mathbb{O}_Γ equals with the ordinary orbit polynomial.

It should be noted that in computing the orbit polynomial of a network, not only the group Γ is important but also the orbit-set \mathcal{X} has a vital role in computing the orbit polynomial. In other words, for two isomorphic but not permutationally isomorphic permutation groups Γ_1 and Γ_2 , we may obtain distinct orbit polynomials. This means that the cycle-type of all permutations or at least the cycle-type of the generators of the automorphism group has a significant role in the structure of orbit polynomial.

We proceed as follows. Section 2 outlines the concepts and definitions that will be used in this paper. In Section 3, we compute the several bounds for the roots of M -orbit polynomial. Finally, in Section 4, we compute some topological indices for certain networks like *Human B Cell*, *BioGRID Drosophila*, *US Airports* and *Email* with QuACN-package [28].

2. Preliminaries

Our notation is standard and mainly taken from standard books of graph theory such as [25]. The vertex and edge sets of a network \mathcal{N} are denoted by $V(\mathcal{N})$ and $E(\mathcal{N})$, respectively. All networks considered in this paper are simple, connected, and finite. For a network \mathcal{N} with automorphism group $\mathbb{A}_{\mathcal{N}} = \text{Aut}(\mathcal{N})$ and an arbitrary vertex $v \in V(\mathcal{N})$, the vertex-orbit of v (or orbit of v) is the set of all $\alpha(v)$'s, where α is an automorphism of \mathcal{N} . Finding the automorphism group of a network could take exponential time, since, for example, the complete graph K_n has S_n as its automorphism group. Let $A(\mathcal{N})$ be the adjacency matrix of network \mathcal{N} and P_σ be a permutation matrix, corresponding to the permutation $\sigma \in S_n$. Then, σ is an automorphism of network \mathcal{N} if and only if $P_\sigma^t A P_\sigma = A$. Here, computing the automorphism group, as well as the orbits of considered networks was done with the aid of the *i-graph* package [12].

3. Location of Roots

In [18, 21–24] several bounds for the positive root δ are given and some properties of the orbit polynomial are investigated. In this section, we obtain several results concerning the location of zeros of the modified orbit polynomial. The following proposition sharpens previous results along with some of the other known results which were based on Walsh classical theorem. Moreover, an *R-code* is developed to construct polynomials, and compare the bounds obtained by our result with these known results.

Let $f(z) = \sum_{k=0}^n a_k z^k$, ($a_k \neq 0$) be a non-constant polynomial with complex coefficients. Then all its zeros lie in disc $C = \{z \in \mathbb{C} : |z| \leq r\}$, where

$$i) \quad [9] \quad r < 1 + \max_{0 \leq k \leq n-1} \{|a_k|\},$$

$$ii) \quad [31] \text{ (Walsh)} \quad r = \sum_{k=0}^{n-1} |a_k|^{1/(n-k)},$$

$$iii) \quad [10] \text{ (Carmichael et. al.)} \quad r = \sqrt{1 + \sum_{k=0}^{n-1} |a_k|^2}.$$

It is not difficult to see that, if G is a vertex-transitive graph, then $\mathbb{O}_G(x) = x^n$, $\mathbb{O}_G^*(x) = 1 - x^n$ and thus its positive root is one. Also, if $\mathbb{A}_G \cong id$ (namely, G is asymmetric graph with identity automorphism group), then $\mathbb{O}_G(x) = nx$, $\mathbb{O}_G^*(x) = 1 - nx$ and its positive root is $\frac{1}{n}$. This yields that for a given graph G on n vertices, which is neither vertex-transitive nor identity graph, then $\delta \in (\frac{1}{n}, 1)$. In the case that G is not vertex-transitive and $\mathbb{A}_G \neq id$, then clearly we have $0 \leq a_i \leq n - 2$ and according to

Theorem 3(i), it can be verified that $r < n - 1$ and $0 < |z| < n - 1$. By Theorem 3(ii), we obtain

$$r = a_0^{1/n} + a_1^{1/n-1} + \dots + a_{n-2}^{1/2} + a_{n-1}. \quad (1)$$

For $i > [n/2]$, it is clear that $a_i \in \{0, 1\}$ and only one of a_i 's is 1. On the other hand, if the order of orbits is smaller than $[n/2]$, the number of terms in the Eq.(1) will be increased and thus the value of r . The maximum number of terms holds if all a_i 's are one. Since $\sum_{i=1}^n ia_i = n$, a graph in which $n = 1 + 2 + 3 + \dots + m$ (for some integer m), has the largest value of r . Hence, $m = \frac{-1+\sqrt{8n}}{2}$ and so $r \leq 1 + m = \frac{1+\sqrt{8n}}{2}$, or equivalently $|z| < \frac{1+\sqrt{8n}}{2}$.

If the number of terms in the Eq.(1) decreases, the value of r will be decreased and if the orders of the orbits of graph are equal, then Eq.(1) will be a monomial. It is not difficult to see that by Theorem 3 (ii), the minimum value of r is $r \geq 1 + 2^{2/n} > 2$, in which $\mathbb{O}_G^*(x) = 1 - 2x^{n/2}$. By Theorem 2.1 (iii), we yield

$$r = \sqrt{1 + \sum_{k=0}^{n-1} |a_k|^2} = \sqrt{1 + (1^2 + a_1^2 + a_2^2 + \dots + a_{n-1}^2)}.$$

Knowing that the graph is not vertex-transitive and $\mathbb{A}_G \neq id$, it can be inferred that $2 \leq \sum a_i \leq n - 1$ and thus

$$\sum a_i^2 < (\sum a_i)^2 \leq (n - 1)^2,$$

or equivalently,

$$r < \sqrt{2 + (n - 1)^2} = (n - 1) \sqrt{1 + \frac{2}{(n - 1)^2}}.$$

Thus, if $n \geq 3$, we may conclude that $r < (n - 1) \sqrt{\frac{3}{2}}$. The lower bound for r using Theorem 3 (iii) is (G is not vertex-transitive and $\mathbb{A}_G \not\cong id$)

$$r \geq \sqrt{1 + (1^2 + 1^2 + 1^2)} = 2 \quad (a_{n-1}, a_1 = 1).$$

Hence, we proved the following theorem.

Theorem 1. Suppose G is a graph on n vertices, then all zeros of modified orbit polynomial \mathbb{O}_G^* lie in disc $C = \{z \in \mathbb{C} : |z| < \frac{1+\sqrt{8n}}{2}\}$.

Corollary 2. For the star graph S_n with n vertices, $\mathbb{O}_{S_n}(x) = x + x^{n-1}$ and all the zeros of $\mathbb{O}_{S_n}^*(x) = 1 - (x + x^{n-1})$ lie in $[-2, 2]$.

Theorem 3. [10] Let

$$f(z) = a_n z^n + a_{n-1} z^{n-1} + \dots + a_0, a_n a_{n-1} \neq 0,$$

be a complex polynomial. All zeros of $f(z)$ lie in

$$|z| \leq \frac{1 + \phi}{2} + \frac{\sqrt{(\phi - 1)^2 + 4M_2}}{2},$$

where

$$\phi := \left| \frac{a_{n-1}}{a_n} \right|,$$

and

$$M_2 := \max_{0 \leq j \leq n-2} \left| \frac{a_j}{a_n} \right|.$$

Theorem 4. Let \mathcal{N} be a network on $n \geq 3$ vertices that is not vertex-transitive and $\mathbb{A}_{\mathcal{N}} \neq id$, with the orbit polynomial $\mathbb{O}_{\mathcal{N}}(x) = \sum_{i=1}^t a_i x^i$. Then all zeros lie in $\left(\frac{-1}{n-2}, n-2\right]$.

Proof. We know that $a_i \leq n-2$, therefore Theorem 3 implies that

$$\text{Max}\phi = \frac{n-2}{1} = n-2.$$

Hence, the orbit polynomial is of form $x^2 + (n-2)x$ and so $M_2 = 0$. Consequently,

$$z \leq \frac{1+(n-2)}{2} + \frac{\sqrt{(n-3)^2}}{2} = n-2.$$

Also

$$\text{Min}\phi = \frac{1}{n-2}.$$

Thus, the orbit polynomial is $(n-2)x^2 + x$. Since $M_2 = 0$ and $\sum_{i=1}^t ia_i = n$, we conclude that $n=3$. This means that if $n>3$ there is no graph with $\text{Min}\phi = \frac{1}{n-2}$.

□

Theorem 5. Let \mathcal{N} be a network on $n \geq 3$ vertices and $(n-1)$ orbits. Then $\mathbb{O}_{\mathcal{N}}(x) = x^2 + (n-2)x$ and $\mathbb{A}_{\mathcal{N}} \cong \mathbb{Z}_2$.

Proof. Since \mathcal{N} has $(n-1)$ orbits, one can easily conclude that there are $(n-2)$ singleton orbits and an orbit of order 2. This completes the proof. □

Example 6. We utilized the Sage software [30] to analyze all networks up to the order of 10 and made an observation about the structure of graphs with orbit polynomial $\mathbb{O}_{\mathcal{N}}(x) = x^2 + (n-2)x$. Our finding revealed that the structure of such graphs falls under one of the categories illustrated in Figure 6.

3.1. The Location of Positive Root δ

In last section, we showed that $\delta \in [\frac{1}{n}, 1]$. In addition, in [3], it has been shown that as δ increases towards one, the graph becomes more symmetric, and as δ tends to zero, the graph will be less symmetric, especially for sufficiently large values of n .

Based on the above discussions, this section aims to develop a framework to discuss the behavior of the positive root of the modified orbit polynomial, which can capture meaningful structural information.

Theorem 7. Let \mathcal{N} be a network on $n \geq 3$ vertices with the orbit polynomial $\mathbb{O}_{\mathcal{N}}(x) = \sum_{i=1}^t a_i x^i$ and $\mathbb{O}_{\mathcal{N}}^*(x) = 1 - \sum_{i=1}^t a_i x^i$. If for $i \in \{1, 2, \dots, n\}$, $m \leq \sqrt[i]{a_i}$, then $\delta \in (0, \frac{1}{m}]$.

Proof. $\mathbb{O}_{\mathcal{N}}^*(\frac{1}{m}) = 1 - (\frac{1}{m}a_1 + \frac{1}{m^2}a_2 + \dots + \frac{1}{m^i}a_i + \dots + \frac{1}{m^t}a_t) \leq 0$, Considering that $\mathbb{O}_{\mathcal{N}}^*(0) > 0$ and $\mathbb{O}_{\mathcal{N}}^*(1) < 0$ it can be concluded $\delta \in (0, \frac{1}{m}]$. □

Example 8. Let \mathcal{N} be a network with the orbit polynomial $\mathbb{O}_{\mathcal{N}}(x) = 4x + 7x^2 + 25x^3$ and $\mathbb{O}_{\mathcal{N}}^*(x) = 1 - (4x + 7x^2 + 25x^3)$. Since $a_1 = 4$ we may conclude that $m \leq 4$ and $a_2 = 7$ gives $m^2 \leq 7$. Also $a_3 = 25$ yields $m^3 \leq 25$ and thus $\max(m) = 4$ and $\delta \in (0, \frac{1}{4}]$.

Corollary 9. The value $\delta = \frac{1}{2}$ holds for a network with orbit polynomial $\mathbb{O}_{\mathcal{N}}(x) = 2^i x^i$ and if the network \mathcal{N} has k singleton orbits, then the positive root $\mathbb{O}_{\mathcal{N}}^*(x)$ lies in $(0, \frac{1}{k}]$.

Theorem 10. For any rational number α in the interval $(-\infty, 0]$ there is a network \mathcal{N} such that $\mathbb{O}_{\mathcal{N}}(\alpha) = 0$. Moregenerally, the set of all roots of $\mathbb{O}_{\mathcal{N}}$ is dense.

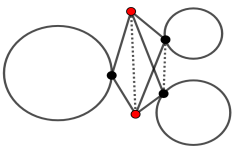
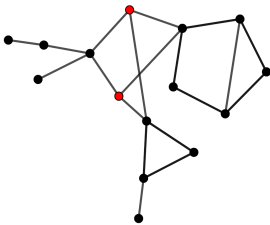


Figure 1. Structure(1).



An example of structure(1).

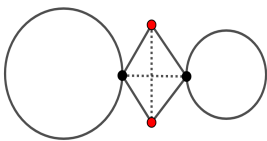
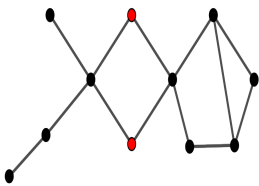


Figure 2. Structure(2).



An example of structure(2).

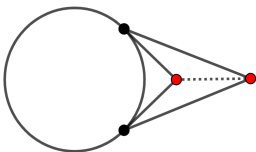
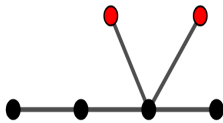


Figure 3. Structure(3).



An example of structure(3).

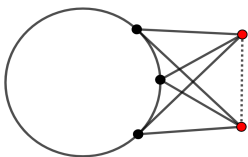
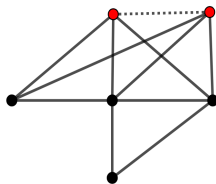


Figure 4. Structure(4).



An example of structure(4).

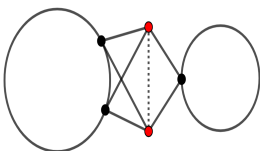
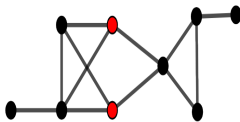


Figure 5. Structure(5).



An example of structure(5).

Figure 6. All structures of graphs with $O_N(x) = x^2 + (n - 2)x$.

Proof. Let $a, b \in \mathbb{N}$. There is a network of order $n = a + 2b$ such that $\mathbb{O}_{\mathcal{N}}(x) = ax + bx^2$. An example of such networks is shown in Figure 7. It is enough to have a network with a vertices and putting two pendant vertices in the position of the b vertices, such that no pair of vertices on the path P_b can be permuted to each other. \square

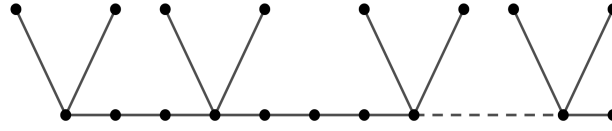


Figure 7. An example of a graph with orbit polynomial $\mathbb{O}_{\mathcal{N}}(x) = ax + bx^2$.

Theorem 11. If the network \mathcal{N} has an orbit of order $n - i$ ($i \leq \frac{n}{2}$), then $\delta \in (\frac{1}{i+1}, 1)$.

Proof. It is clear that $a_{n-i} = 1$ and i vertices lie in orbits by order smaller than $\frac{n}{2}$. So the orbit polynomial as follows: $\mathbb{O}_{\mathcal{N}}(x) = a_1x^1 + a_2x^2 + \dots + a_kx^k + x^{n-i}$ ($k < \frac{n}{2}$). The maximum value of $\mathbb{O}_{\mathcal{N}}(\delta)$ holds if $\mathbb{O}_{\mathcal{N}}(x) = ix$ ($a_1 = i$) and $\mathbb{O}_{\mathcal{N}}^*(\frac{1}{i+1}) = i(\frac{1}{i+1}) < 1$. This means that $\delta \in (\frac{1}{i+1}, 1)$.
The maximum value of $\mathbb{O}_{\mathcal{N}}(\delta)$ holds if $\mathbb{O}_{\mathcal{N}}(x) = ix + x^{n-i}$ ($a_1 = i$) and $\mathbb{O}_{\mathcal{N}}^*(\frac{1}{i+1}) = 1 - (i(\frac{1}{i+1}) + (\frac{1}{i+1})^{n-i}) > 0$. This means that $\delta \in (\frac{1}{i+1}, 1)$.
 \square

Theorem 12. Suppose the tree T has an orbit of order 2 ($|\mathbb{O}_i| = 2$) and the other orbits are singleton, then $\langle \mathbb{O}_i \rangle \cong P_3$.

Proof. If $\mathbb{O}_i = \{y, z\}$ it is clear that there is an automorphism like φ so that $\varphi(y) = z$. Now if y and z lie in different branches of T such as B_1 and B_2 , then it follows from $\varphi(y) = z$ that $\varphi(B_1) = B_2$ which is contradiction. Therefore y and z are in the same branch like B . Let's assume that the root of the branch B is x , and the paths from y and z to x are denoted by p_1 and p_2 , respectively. So $\varphi(p_1) = p_2$ and if t is a vertex on p_1 , then there is a vertex like t' on p_2 which $\varphi(t) = t'$ that is contradiction. Hence, $d(x, y) = d(x, z) = 1$ and the issue is confirmed. \square

Corollary 13. Let T be a tree on $n \geq 3$ vertices and $\mathbb{O}_{\mathcal{T}}(x) = x^2 + (n-2)x$, then

$$\delta = \frac{2}{\sqrt{(n-2)^2 + 4} + (n-2)}.$$

Proof. Suppose $\mathbb{O}_{\mathcal{T}}^*(x) = 1 - (x^2 + (n-2)x)$, then clearly the roots are

$$x_{1,2} = \frac{-(n-2) \pm \sqrt{(n-2)^2 + 4}}{2}$$

and thus

$$\delta = \frac{2}{\sqrt{(n-2)^2 + 4} + (n-2)}.$$

This completes the proof. \square

Corollary 14. Let T be a tree on n vertices:

- 1) If $\mathbb{A}_T \cong id$, then $\delta = \frac{1}{n}$.
- 2) If $\mathbb{A}_T \neq id$, then $\delta \in [\alpha, 1]$ so that

$$\alpha = \frac{\sqrt{(n-2)^2 + 4} - (n-2)}{2}.$$

137 **4. The Role of δ in the Study of Real Network Structures**

Complex networks have been extensively studied in various fields in recent years. The concept of entropies has been used to characterize and quantify the structure of networks, and has been investigated in many studies [1,2,4–8,11]. As a specific definition for graph entropies, we use the definition proposed by Dehmer [15–17,29], which involves a probability vector $p = (p_1, \dots, p_n)$ that satisfies two conditions: $0 \leq p_i \leq 1$ and $\sum_i p_i = 1$. The Shannon’s entropy is then defined as

$$I(p) = - \sum_{i=1}^n p_i \log(p_i).$$

138 Mowshowitz and Dehmer [27] introduced the symmetry index $S(G)$, which is defined
139 as

$$\begin{aligned} S(G) &= (\log n - I_a(G)) + \log |\mathbb{A}_G| \\ &= \frac{1}{n} \left(\sum_{i=1}^t |O_i| \log |O_i| \right) + \log |\mathbb{A}_G|. \end{aligned}$$

140 If $I(G)$ and $I(H)$ are two graph invariants of graphs G and H , respectively, the graph
141 distance measure between $I(G)$ and $I(H)$ is defined as $d_I(G, H) = 1 - e^{-\left(\frac{I(G)-I(H)}{\sigma}\right)^2}$,
142 according to [14]. The values of orbit entropy I_a , size of the automorphism group,
143 symmetric index $S(G)$, unique positive root δ , distance measure between δ and I_a , and
144 distance measure between δ and $S(G)$ of biological and technological networks are
145 reported in Table 1.

Table 1:

	n	$ \mathbb{A} $	δ	I_a	S	$d(\delta, I_a)$	$d(\delta, S)$
Human B Cell Genetic	5930	5.94×10^{13}	1710×10^{-7}	3.77	13.8	14.20	189.83
Caenorhabditis elegans Genetic	2060	6.99×10^{161}	6258×10^{-7}	3.20	161.96	10.22	26.23×10^3
BioGRID Human	7013	1.26×10^{485}	1764×10^{-7}	3.73	485.2	13.93	23.54×10^4
BioGRID Saccharomyces cerevisiae	5295	6.86×10^{64}	1958×10^{-7}	3.70	64.86	13.71	4206.4
BioGRID Drosophila	7371	3.07×10^{493}	1690×10^{-7}	3.76	493.6	14.10	24.36×10^4
BioGRID Mus musculus	209	5.35×10^{125}	221440×10^{-7}	1.46	126.59	2.06	16.02×10^3
Yeast Protein Interactions	1458	1.07×10^{254}	11599×10^{-7}	2.88	254.30	8.31	64.67×10^3
c. elegans metabolic	453	1.93×10^{10}	25702×10^{-7}	2.60	10.33	6.78	106.79
Internet	22332	$1.28 \times 10^{11,298}$	1035×10^{-7}	3.67	11.3×10^3	13.44	12.77×10^7
US Power Grid	4941	5.18×10^{152}	2380×10^{-7}	3.63	152.77	13.20	23.34×10^3
US Airports	332	2.59×10^{24}	40472×10^{-7}	2.39	24.54	5.71	602.07
www California search subnet	5925	$1.24 \times 10^{1,298}$	2820×10^{-7}	3.45	1298.4	11.93	16.86×10^5
www EPA.gov subnet	4253	$1.28 \times 10^{2,321}$	4992×10^{-7}	2.91	2321.8	8.47	53.91×10^5
www Political Blogs	1222	2.40×10^{35}	8741×10^{-7}	3.04	35.43	9.25	1254.9
Email	1133	1.53×10^9	9216×10^{-7}	3.04	9.20	9.23	84.62
Media ownership	4475	$3.36 \times 10^{4,818}$	13278×10^{-7}	2.16	4820.01	4.68	23.23×10^6
Geometry Co-authorship	3621	1.90×10^{320}	4419×10^{-7}	3.38	320.45	11.45	10.27×10^4
Erdős Collaboration	6927	$3.46 \times 10^{4,222}$	5491×10^{-7}	2.956	4223.4	8.73	17.84×10^6
PhD network	1025	2.98×10^{292}	25245×10^{-7}	2.55	292.93	6.49	85.81×10^3

146 Table 2 shows that the correlation values of δ and I_a in biological and technological
147 networks are very close. However, it appears that there may be differences in the effective

Table 2: The correlation values.

Biological Networks				Technological Networks				Social Networks			
δ	I_a	S		δ	I_a	S		δ	I_a	S	
1	-0.88	-0.20		1	-0.86	-0.37		1	-0.64	-0.19	
-0.88	1	0.33		-0.86	1	0.48		-0.64	1	-0.55	
-0.20	0.33	1		-0.37	0.48	1		-0.19	-0.55	1	

variables that contribute to the structure of social networks compared to other types of networks. This suggests that further research should be conducted to investigate these differences.

Table 3 presents a collection of well-known real-world networks with distinct topologies. Analysis of the reported data shows that the symmetry measure δ is strongly correlated with the orbit entropy I_a (with a correlation coefficient $R = -0.73$), but does not appear to be correlated with $S(G)$ (with a correlation coefficient $R = -0.15$).

For the networks listed in Table 3, various topological indices, including the first Zagreb index (M_1), second Zagreb index (M_2), spectral radius (ρ), Randic index (R), Laplacian Estrada index (LEE), Laplacian energy (LE), Harary index (H), Estrada index (EE), energy (E), Balaban ID(BI), and atom-bond connectivity (ABC), were calculated, as reported in [28]. The results indicate that among all the above indices, Laplacian energy has the strongest correlation with δ , with a coefficient of greater than 0.91.

δ	M_1	M_2	ρ	R	LEE	LE	H	EE	E	BI	ABC
-0.07	0.47	-0.19	-0.17	-0.02	0.91	-0.17	-0.16	-0.20	-0.23	-0.20	-0.20

For this study, all graphs of orders 4-7, all trees of orders 7-20, together with 470 randomly generated trees of order 21 and 1248 randomly generated trees of order 25 were considered, and two measures based on the automorphism group and Laplacian eigenvalues of a graph were established. The computed correlation values between the unique positive root δ and Laplacian energy of a graph, denoted by $LE(G)$, are reported in Table 5. For trees of order 21 and 25, the correlation values were 0.556 and 0.589, respectively, while for other classes they were less than 0.5.

Although the correlation between δ and $LE(G)$ in real networks is meaningful, it appears that in the class of trees, the correlation value between δ and $LE(G)$ increases as the order of the tree increases. In other words, an appropriate analogy is the case of two functions that have very different functional forms on the same set of variables.

Table 4:

#Graphs	Graph Order	$R(\delta, LE)$
9	2 to 4	-0.414
21	5	-0.009
121	6	0.065
853	7	0.029
33	unicycle graph of order 7	0.031
11	Tree of order 7	0.290
3159	Tree of order 14	0.381
7741	Tree of order 15	0.441
19320	Tree of order 16	0.429
48629	Tree of order 17	0.466
123867	Tree of order 18	0.466
317955	Tree of order 19	0.493
823065	Tree of order 20	0.499

Table 5: The correlation values between randomly graphs.

5. Small-World Graphs

A social network is called a small-world network, if most nodes are not neighbors of one another, but most nodes can be reached from every other by a small number of hops. The clustering coefficient of a graph is a measure to find which vertices in the

Table 3: A set of well-known real-world networks with distinct topologies and their graph invariants.

Networks	n	δ	M_1	M_2	ρ	R	LEE	LE	H	EE	E	BI	ABC
Human B Cell Genetic	5930	1710×10^{-7}	50502	63.7×10^7	158.8	159.04	-	50.5×10^3	40.3×10^6	9.48×10^{68}	327.75	-	-
Caenorhabditis elegans Genetic	2060	6258×10^{-7}	13410	6803	1.62	66.8×10^2	56.6×10^3	13.3×10^3	67.3×10^2	20.6×10^3	13.3×10^3	13362	66.47
BioGRID Saccharomyces cerevisiae	5295	1958×10^{-7}	12032	27.3×10^5	1	0.11	142.61	11998	17	52.46	34	34	1.61
BioGRID Drosophila	7371	1690×10^{-7}	61434	18.7×10^7	2	1.58	224.7	61410	15	35.21	20	21	4.19
BioGRID Mus musculus	209	221440×10^{-7}	18.5×10^4	47.3×10^7	3.87	1.08	88.9×10^5	18.5×10^4	89.5	130	51.74	60	9.89
Yeast Protein Interactions	1458	11599×10^{-7}	23710	19.4×10^6	65.75	1076.3	7.11	23.8×10^3	61.5×10^4	3.60	4141.3	-	4349.1
US Power Grid	4941	2380×10^{-7}	8	16	1	0.25	8.39	6	1	3.09	2	6	0.61
US Airports	332	40472×10^{-7}	26100	14.5×10^7	53.51	148.6	7.9×10^{72}	21947	90490	1.7×10^{23}	1419	-	-
Email	1133	9216×10^{-7}	10902	19.7×10^5	20.7	460.5	2.54×10^{31}	8399.9	19.2×10^4	10.5×10^8	2429.9	-	2203.4

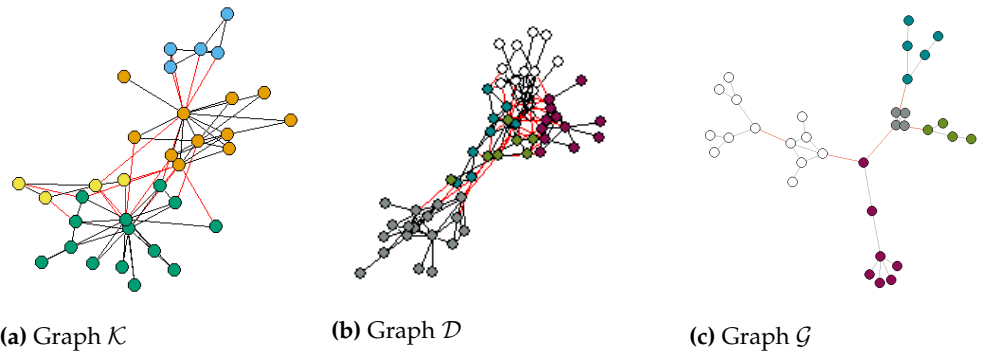


Figure 8. Three small-world graphs (a) Karate Club \mathcal{K} , (b) Dolphin graph \mathcal{D} , and (c) the network \mathcal{G} .

graph tend to form a clique. There are two types of clustering coefficient, namely local and global. The local clustering coefficient $C(v)$ of vertex v is defined as $\frac{2e(v)}{d_v(d_v-1)}$, where $e(v)$ is the number of edges $\langle N_G(v) \rangle$. The global clustering coefficient of graph G is $C^\Delta = \frac{3t}{l_2}$, where t is the number of triangles and l_2 is the number of paths of length two [32]. Assume that C_{rand}^Δ is the clustering coefficient of an equivalent ER random graph with the same order and size. In [32], it is proved that $L_{rand} = \frac{\ln(n)}{\ln(2k)}$ and $C_{rand}^\Delta = \frac{2k}{n}$. Additionally, the small-worldness of a graph is expressed by

$$S^\Delta = \frac{C_g^\Delta L_{rand}}{C_{rand}^\Delta L_g}.$$

Example 15. Consider three graphs Karate Club \mathcal{K} , Dolphin graph \mathcal{D} , and the network \mathcal{G} , as depicted in Figure 8. The following R programs determine their small-worldness which are respectively 1.48, 2.10, and 4.58.

The aim of continuing this paper is to determine whether removing a vertex or an edge can decrease or increase the small-worldness of a graph. In the second step, our goal is to determine the small-worldness of two small-world graphs linked by a new edge. To do this, we consider 1000 random small-world networks of the same order. Our results show that by removing a random edge, the resulting graph is again a small-world graph with small-worldness greater than one. We repeated this way by removing a random vertex to conclude a similar result. We applied the following program written by R to obtain our results.

An R program for computing smallworldness of networks:

```
library(NetworkToolbox)
library(igraph)
library(Matrix)
count<-0
repeat{
  a<-sample(50:60,1)
  b<-sample(3:8,1)
  g<-sample_smallworld(1,a,b,0.05)
  H<-as_adj(g)
  L<-matrix(H,nrow=a,ncol=a)
  x<-smallworldness(L,method="rand")
  p<-sample(V(g),1)
  g1<-delete_vertices(g,p)
  A<-as_adjacency_matrix(g1)
  B<-matrix(A,nrow=a-1,ncol=a-1)
  y<-smallworldness(B,method="rand")
  t<-cbind(x$swm,y$swm)
```

```

203 write.table(t,file="G-v.csv",append=T,sep=","col.names=F,row.names=F)
204 count<-count+1
205 if(count==1000){break}
206 }

```

For the second step, we also generate 1000 random small-world graphs joined by a new edge and the resulted graph is again small-world by almost the same small-worldness. The following program do this.

Algorithm 2:

```

211 library(igraph)
212 library(Matrix)
213 library(NetworkToolbox)
214 library(rgl)
215 count<-0
216 repeat{
217   a<-sample(50:100,1)
218   b<-sample(3:8,1)
219   g1<-sample_smallworld(1,a,b,0.05)
220   H<-as_adj(g1)
221   L<-matrix(H,nrow=a,ncol=a)
222   x<-smallworldness(L,method="rand")
223   d<-sample(50:100,1)
224   e<-sample(3:8,1)
225   g2<-sample_smallworld(1,d,e,0.05)
226   K<-as_adj(g2)
227   M<-matrix(K,nrow=d,ncol=d)
228   y<-smallworldness(M,method="rand")
229   g3<-(g1+g2)%<% add_edegs(c(1,a+d))
230   A<-as_adjacency_matrix(g3)
231   B<-matrix(A,nrow=a+d,ncol=a+d)
232   z<-smallworldness(B,method="rand")
233   write.table(c(x,y,z),file="new.csv",append=T,sep=","col.names=F,row.names=F)
234   count<-count+1
235   if(count==1000){break}
236 }

```

To summarize, the results of the investigation indicate that the δ metric is not redundant, and offers unique information about the symmetry structure of graphs that cannot be obtained from standard measures such as small-worldness. Additionally, the successful use of δ as a structural descriptor for chemical structures highlights its potential applications in diverse fields beyond network analysis.

However, further research is necessary to determine the full range of structural properties that can be captured by the orbit polynomial, and to better understand the relationship between δ and other graph metrics. Overall, the results suggest that δ is a valuable and versatile tool for analyzing complex systems and may have important implications for a variety of fields.

6. Conclusion

Dehmer et al. proposed a novel approach to compare graphs based on their symmetry structure using the modified orbit polynomial and the unique positive root, δ . The characteristics of δ have been demonstrated and it has been applied successfully to several network classes to investigate correlations with other graph indices.

Furthermore, the successful application of δ as a structural descriptor for chemical structures is significant because it opens up new possibilities for studying the symmetry structure of various complex systems beyond networks, such as molecules and crystals.

This demonstrates the versatility and potential of δ in a range of fields beyond network analysis.

In conclusion, the orbit polynomial and its associated metric, $\delta(G)$, provide a novel approach to comparing graphs based on their symmetry structure.

However, further research is needed to fully understand the range of structural properties captured by the orbit polynomial and to investigate the relationship between $S(G)$ and I_a , as both are defined based on the number and sizes of orbits belonging to a graph's automorphism group. This would provide a more comprehensive understanding of the relationship between symmetry and other graph metrics, and could potentially lead to new insights and applications for δ and related metrics.

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