On sets $X \subseteq \mathbb{N}$ whose finiteness implies that we know an algorithm which for every $n \in \mathbb{N}$ decides the inequality $\max(X) < n$

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Abstract

Let $\Gamma(k)$ denote (k-1)!, and let $\Gamma_n(k)$ denote (k-1)!, where $n \in \{3, ..., 16\}$ and $k \in \{2\} \cup [2^{2^{n-3}}+1,\infty) \cap \mathbb{N}$. For an integer $n \in \{3,...,16\}$, let Σ_n denote the following statement: if a system of equations $S \subseteq \{\Gamma_n(x_i) = x_k : i, k \in \{1,...,n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1,...,n\}\}$ with Γ instead of Γ_n has only finitely many solutions in positive integers x_1, \ldots, x_n , then every tuple $(x_1,\ldots,x_n) \in (\mathbb{N} \setminus \{0\})^n$ that solves the original system S satisfies $x_1,\ldots,x_n \leqslant 2^{2^{n-2}}$. Our hypothesis claims that the statements $\Sigma_3,\ldots,\Sigma_{16}$ are true. The statement Σ_6 proves the following implication: if the equation x(x+1)=y! has only finitely many solutions in positive integers x and y, then each such solution (x,y) belongs to the set $\{(1,2),(2,3)\}$. The statement Σ_6 proves the following implication: if the equation $x!+1=y^2$ has only finitely many solutions in positive integers x and y, then each such solution (x,y) belongs to the set $\{(4,5),(5,11),(7,71)\}$. The statement Σ_9 implies the infinitude of primes of the form x has the infinitude of primes of the form x has the infinitude of primes of the form x has the infinitude of Sophie Germain primes.

Key words and phrases: Brocard's problem, Brocard-Ramanujan equation $x! + 1 = y^2$, composite Fermat numbers, Erdös' equation x(x + 1) = y!, prime numbers of the form $n^2 + 1$, prime numbers of the form n! + 1, Sophie Germain primes, twin primes.

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1 Definitions and lemmas

Definition 1. We say that a non-negative integer m is a threshold number of a set $X \subseteq \mathbb{N}$, if X is infinite if and only if X contains an element greater than m, cf. [17] and [18].

If a set $X \subseteq \mathbb{N}$ is empty or infinite, then any non-negative integer m is a threshold number of X. If a set $X \subseteq \mathbb{N}$ is non-empty and finite, then the all threshold numbers of X form the set $\{\max(X), \max(X) + 1, \max(X) + 2, \ldots\}$.

Definition 2. We say that a set $X \subseteq \mathbb{N}$ is uncomplicated if we know an algorithm which takes as input a non-negative integer n and returns $\mathbf{B} \in \{\text{TRUE}, \text{FALSE}\}$, such that \mathbf{B} equals the truth value of the inequality $\max(X) < n$, if X is finite.

Theorem 1. If a set $X \subseteq \mathbb{N}$ is finite and uncomplicated, then we know an algorithm which computes a threshold number of X.

Proof. We compute the smallest non-negative integer n such that $\max(X) < n + 1$.

Lemma 1. For every positive integers x and y, $x! \cdot y = y!$ if and only if

$$(x + 1 = y) \lor (x = y = 1)$$



Let $\Gamma(k)$ denote (k-1)!.

Lemma 2. For every positive integers x and y, $x \cdot \Gamma(x) = \Gamma(y)$ if and only if

$$(x + 1 = y) \lor (x = y = 1)$$

Lemma 3. For every non-negative integers b and c, b + 1 = c if and only if $2^{2^b} \cdot 2^{2^b} = 2^{2^c}$.

Lemma 4. (Wilson's theorem, [4, p. 89]). For every positive integer x, x divides (x - 1)! + 1 if and only if x = 1 or x is prime.

2 A variant of chess provides an example which explains the title

Let us assume that there are no draws, castlings, and en passant captures. Let us assume that a player with no moves loses. As such, the game may continue forever. Let \mathcal{H} denote the set of all positive integers n such that an appropriate strategy of Black guarantees that White cannot enforce a win in less than n moves.

Lemma 5. ([11, p. 128]). A player who is in a winning position is always able to enforce a win in a number of moves that is less than the number of positions in the game.

Lemma 6. The number of positions does not exceed 13⁶⁴.

Proof. With castlings or en passant captures, a legality of a move depends not only on the positions of the pieces on the board. Without castlings and en passant captures, we observe that 13 corresponds to 12 distinct pieces and the empty square.

Lemmas 5 and 6 imply the following corollary.

Corollary 1. If White have a winning strategy, then $\mathcal{H} \subseteq [1, 13^{64} - 1]$. Otherwise, $\mathcal{H} = \mathbb{N} \setminus \{0\}$. The number $13^{64} - 1$ is a threshold number of \mathcal{H} , and we can decide the equality $\mathcal{H} = \mathbb{N} \setminus \{0\}$. If $\mathcal{H} \neq \mathbb{N} \setminus \{0\}$, then we can compute \mathcal{H} and $\max(\mathcal{H})$. Always, the set \mathcal{H} is uncomplicated.

A Diophantine equation whose non-solvability expresses the consistency of *ZFC*

Gödel's second incompleteness theorem and the Davis-Putnam-Robinson-Matiyasevich theorem imply the following theorem.

Theorem 2. ([3, p. 35]). There exists a polynomial $D(x_1, ..., x_m)$ with integer coefficients such that if ZFC is arithmetically consistent, then the sentences "The equation $D(x_1, ..., x_m) = 0$ is solvable in non-negative integers" and "The equation $D(x_1, ..., x_m) = 0$ is not solvable in non-negative integers" are not provable in ZFC.

Let \mathcal{Y} denote the set of all non-negative integers k such that the equation $D(x_1, \ldots, x_m) = 0$ has no solutions in $\{0, \ldots, k\}^m$. Since the set $\{0, \ldots, k\}^m$ is finite, we know an algorithm which for every $n \in \mathbb{N}$ decides whether or not $n \in \mathcal{Y}$. Let $\gamma \colon \mathbb{N}^{m+1} \to \mathbb{N}$ be a computable bijection, and let $\mathcal{E} \subseteq \mathbb{N}^{m+1}$ be the solution set of the equation $D(x_1, \ldots, x_m) + 0 \cdot x_{m+1} = 0$. Theorem 2 implies Theorems 3 and 4.

Theorem 3. *If* ZFC *is arithmetically consistent, then for every* $n \in \mathbb{N}$ *the sentences* "n is a threshold number of \mathcal{Y} " *and* "n is not a threshold number of \mathcal{Y} " *are not provable in* ZFC.

Theorem 4. We know an algorithm which for every $n \in \mathbb{N}$ decides whether or not $n \in \gamma(\mathcal{E})$. The set $\gamma(\mathcal{E})$ is empty or infinite. In both cases, every non-negative integer n is a threshold number of $\gamma(\mathcal{E})$. If ZFC is arithmetically consistent, then the sentences " $\gamma(\mathcal{E})$ is empty", " $\gamma(\mathcal{E})$ is not empty", " $\gamma(\mathcal{E})$ is finite", and " $\gamma(\mathcal{E})$ is infinite" are not provable in ZFC.

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4 Hypothetical statements $\Psi_3, \ldots, \Psi_{16}$

For an integer $n \ge 3$, let \mathcal{U}_n denote the following system of equations:

$$\begin{cases}
\forall i \in \{1, \dots, n-1\} \setminus \{2\} \ x_i! = x_{i+1} \\
x_1 \cdot x_2 = x_3 \\
x_2 \cdot x_2 = x_3
\end{cases}$$

The diagram in Figure 1 illustrates the construction of the system \mathcal{U}_n .

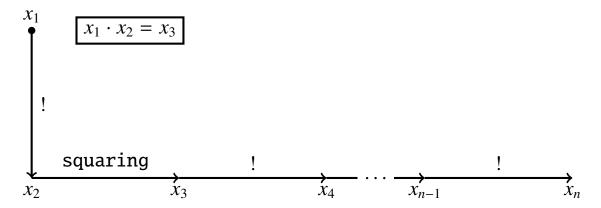


Fig. 1 Construction of the system \mathcal{U}_n

Let g(3) = 4, and let g(n + 1) = g(n)! for every integer $n \ge 3$.

Lemma 7. For every integer $n \ge 3$, the system \mathcal{U}_n has exactly two solutions in positive integers, namely $(1, \ldots, 1)$ and $(2, 2, g(3), \ldots, g(n))$.

Let

$$B_n = \{x_i! = x_k : (i, k \in \{1, \dots, n\}) \land (i \neq k)\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}\}$$

For an integer $n \ge 3$, let Ψ_n denote the following statement: if a system $S \subseteq B_n$ has only finitely many solutions in positive integers x_1, \ldots, x_n , then each such solution (x_1, \ldots, x_n) satisfies $x_1, \ldots, x_n \le g(n)$. The statement Ψ_n says that for subsystems of B_n the largest known solution is indeed the largest possible.

Hypothesis 1. The statements Ψ_3, \dots, Ψ_{16} are true.

Theorem 5. Every statement Ψ_n is true with an unknown integer bound that depends on n.

Proof. For every positive integer n, the system B_n has a finite number of subsystems.

Theorem 6. For every statement Ψ_n , the bound g(n) cannot be decreased.

Proof. It follows from Lemma 7 because $\mathcal{U}_n \subseteq B_n$.

5 The Brocard-Ramanujan equation $x! + 1 = y^2$

Let $\mathcal A$ denote the following system of equations:

$$\begin{cases} x_1! = x_2 \\ x_2! = x_3 \\ x_5! = x_6 \\ x_4 \cdot x_4 = x_5 \\ x_3 \cdot x_5 = x_6 \end{cases}$$

Lemma 1 and the diagram in Figure 2 explain the construction of the system \mathcal{A} .

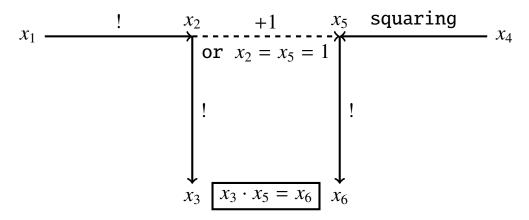


Fig. 2 Construction of the system \mathcal{A}

Lemma 8. For every $x_1, x_4 \in \mathbb{N} \setminus \{0, 1\}$, the system \mathcal{A} is solvable in positive integers x_2, x_3, x_5, x_6 if and only if $x_1! + 1 = x_4^2$. In this case, the integers x_2, x_3, x_5, x_6 are uniquely determined by the following equalities:

$$x_2 = x_1!$$

 $x_3 = (x_1!)!$
 $x_5 = x_1! + 1$
 $x_6 = (x_1! + 1)!$

Proof. It follows from Lemma 1.

It is conjectured that x! + 1 is a perfect square only for $x \in \{4, 5, 7\}$, see [14, p. 297]. A weak form of Szpiro's conjecture implies that there are only finitely many solutions to the equation $x! + 1 = y^2$, see [9].

Theorem 7. If the equation $x_1! + 1 = x_4^2$ has only finitely many solutions in positive integers, then the statement Ψ_6 guarantees that each such solution (x_1, x_4) belongs to the set $\{(4, 5), (5, 11), (7, 71)\}$.

Proof. Suppose that the antecedent holds. Let positive integers x_1 and x_4 satisfy $x_1! + 1 = x_4^2$. Then, $x_1, x_4 \in \mathbb{N} \setminus \{0, 1\}$. By Lemma 8, the system \mathcal{A} is solvable in positive integers x_2, x_3, x_5, x_6 . Since $\mathcal{A} \subseteq B_6$, the statement Ψ_6 implies that $x_6 = (x_1! + 1)! \leq g(6) = g(5)!$. Hence, $x_1! + 1 \leq g(5) = g(4)!$. Consequently, $x_1 < g(4) = 24$. If $x_1 \in \{1, ..., 23\}$, then $x_1! + 1$ is a perfect square only for $x_1 \in \{4, 5, 7\}$.

Corollary 2. Let X denote the set $\{x \in \mathbb{N} \setminus \{0\} : x! + 1 \text{ is a perfect square}\}$. The statement Ψ_6 implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. In order to prove item (a), we consider the algorithm which decides the inequality 7 < n. In order to prove items (b) and (c), we consider the algorithm which returns 7.

6 Are there infinitely many prime numbers of the form $n^2 + 1$?

Let \mathcal{B} denote the following system of equations:

$$\begin{cases} x_2! = x_3 \\ x_3! = x_4 \\ x_5! = x_6 \\ x_8! = x_9 \\ x_1 \cdot x_1 = x_2 \\ x_3 \cdot x_5 = x_6 \\ x_4 \cdot x_8 = x_9 \\ x_5 \cdot x_7 = x_8 \end{cases}$$

Lemma 1 and the diagram in Figure 3 explain the construction of the system \mathcal{B} .

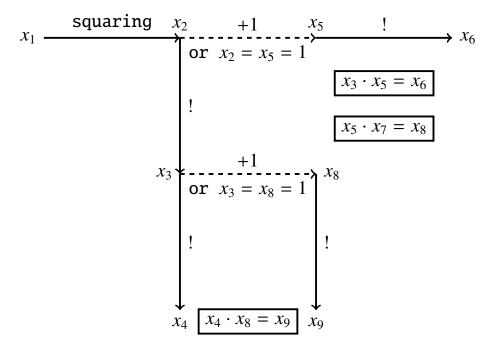


Fig. 3 Construction of the system \mathcal{B}

Lemma 9. For every integer $x_1 \ge 2$, the system \mathcal{B} is solvable in positive integers x_2, \ldots, x_9 if and only if $x_1^2 + 1$ is prime. In this case, the integers x_2, \ldots, x_9 are uniquely determined by the following equalities:

$$x_{2} = x_{1}^{2}$$

$$x_{3} = (x_{1}^{2})!$$

$$x_{4} = ((x_{1}^{2})!)!$$

$$x_{5} = x_{1}^{2} + 1$$

$$x_{6} = (x_{1}^{2} + 1)!$$

$$x_{7} = \frac{(x_{1}^{2})! + 1}{x_{1}^{2} + 1}$$

$$x_{8} = (x_{1}^{2})! + 1$$

$$x_{9} = ((x_{1}^{2})! + 1)!$$

Proof. By Lemma 1, for every integer $x_1 \ge 2$, the system \mathcal{B} is solvable in positive integers x_2, \ldots, x_9 if and only if $x_1^2 + 1$ divides $(x_1^2)! + 1$. Hence, the claim of Lemma 9 follows from Lemma 4.

Lemma 10. There are only finitely many tuples $(x_1, ..., x_9) \in (\mathbb{N} \setminus \{0\})^9$ which solve the system \mathcal{B} and satisfy $x_1 = 1$.

Proof. If a tuple $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$ solves the system \mathcal{B} and $x_1 = 1$, then $x_1, \ldots, x_9 \le 2$. Indeed, $x_1 = 1$ implies that $x_2 = x_1^2 = 1$. Hence, for example, $x_3 = x_2! = 1$. Therefore, $x_8 = x_3 + 1 = 2$ or $x_8 = 1$. Consequently, $x_9 = x_8! \le 2$.

Edmund Landau's conjecture states that there are infinitely many primes of the form $n^2 + 1$, see [8, pp. 37–38].

Theorem 8. The statement Ψ_9 proves the following implication: if there exists an integer $x_1 \ge 2$ such that $x_1^2 + 1$ is prime and greater than g(7), then there are infinitely many primes of the form $n^2 + 1$.

Proof. Suppose that the antecedent holds. By Lemma 9, there exists a unique tuple $(x_2, ..., x_9) \in (\mathbb{N} \setminus \{0\})^8$ such that the tuple $(x_1, x_2, ..., x_9)$ solves the system \mathcal{B} . Since $x_1^2 + 1 > g(7)$, we obtain that $x_1^2 \ge g(7)$. Hence, $(x_1^2)! \ge g(7)! = g(8)$. Consequently,

$$x_9 = ((x_1^2)! + 1)! \ge (g(8) + 1)! > g(8)! = g(9)$$

I sets \(\lambda \leq 1\) whose initieness ...

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Since $\mathcal{B} \subseteq B_9$, the statement Ψ_9 and the inequality $x_9 > g(9)$ imply that the system \mathcal{B} has infinitely many solutions $(x_1, \dots, x_9) \in (\mathbb{N} \setminus \{0\})^9$. According to Lemmas 9 and 10, there are infinitely many primes of the form $n^2 + 1$.

Corollary 3. Let X denote the set of primes of the form $n^2 + 1$. The statement Ψ_9 implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. In order to prove item (a), we compute $\max(X \cap [1, g(7)])$ and consider the algorithm which decides the inequality $\max(X \cap [1, g(7)]) < n$. In order to prove items (b) and (c), we consider the algorithm which computes $\max(X \cap [1, g(7)])$.

7 Are there infinitely many prime numbers of the form n! + 1?

It is conjectured that there are infinitely many primes of the form n! + 1, see [2, p. 443] and [12].

Theorem 9. (cf. Theorem 13). The statement Ψ_9 proves the following implication: if there exists an integer $x_1 \ge g(6)$ such that $x_1! + 1$ is prime, then there are infinitely many primes of the form n! + 1.

Proof. We leave the analogous proof to the reader.

Corollary 4. Let X denote the set of primes of the form n! + 1. The statement Ψ_9 implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. We repeat the proof of Corollary 3.

8 The twin prime conjecture

A twin prime is a prime number that differs from another prime number by 2. The twin prime conjecture states that there are infinitely many twin primes, see [8, p. 39]. Let C denote the following system of equations:

$$x_{1} = x_{2}$$

$$x_{2}! = x_{3}$$

$$x_{4}! = x_{5}$$

$$x_{6}! = x_{7}$$

$$x_{7}! = x_{8}$$

$$x_{9}! = x_{10}$$

$$x_{12}! = x_{13}$$

$$x_{15}! = x_{16}$$

$$x_{2} \cdot x_{4} = x_{5}$$

$$x_{5} \cdot x_{6} = x_{7}$$

$$x_{7} \cdot x_{9} = x_{10}$$

$$x_{4} \cdot x_{11} = x_{12}$$

$$x_{3} \cdot x_{12} = x_{13}$$

$$x_{9} \cdot x_{14} = x_{15}$$

$$x_{8} \cdot x_{15} = x_{16}$$

Lemma 1 and the diagram in Figure 4 explain the construction of the system C.

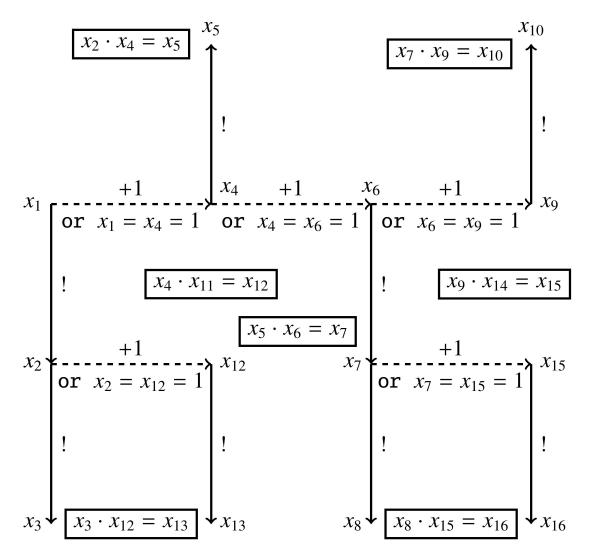


Fig. 4 Construction of the system C

Lemma 11. For every $x_4, x_9 \in \mathbb{N} \setminus \{0, 1, 2\}$, the system C is solvable in positive integers $x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$ if and only if x_4 and x_9 are prime and $x_4 + 2 = x_9$. In this case, the integers $x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$ are uniquely determined by the following equalities:

$$x_{1} = x_{4} - 1$$

$$x_{2} = (x_{4} - 1)!$$

$$x_{3} = ((x_{4} - 1)!)!$$

$$x_{5} = x_{4}!$$

$$x_{6} = x_{9} - 1$$

$$x_{7} = (x_{9} - 1)!$$

$$x_{8} = ((x_{9} - 1)!)!$$

$$x_{10} = x_{9}!$$

$$x_{11} = \frac{(x_{4} - 1)! + 1}{x_{4}}$$

$$x_{12} = (x_{4} - 1)! + 1$$

$$x_{13} = ((x_{4} - 1)! + 1)!$$

$$x_{14} = \frac{(x_{9} - 1)! + 1}{x_{9}}$$

$$x_{15} = (x_{9} - 1)! + 1$$

$$x_{16} = ((x_{9} - 1)! + 1)!$$

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Proof. By Lemma 1, for every $x_4, x_9 \in \mathbb{N} \setminus \{0, 1, 2\}$, the system *C* is solvable in positive integers $x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}$ if and only if

$$(x_4 + 2 = x_9) \wedge (x_4|(x_4 - 1)! + 1) \wedge (x_9|(x_9 - 1)! + 1)$$

Hence, the claim of Lemma 11 follows from Lemma 4.

Lemma 12. There are only finitely many tuples $(x_1, ..., x_{16}) \in (\mathbb{N} \setminus \{0\})^{16}$ which solve the system C and satisfy

$$(x_4 \in \{1, 2\}) \lor (x_9 \in \{1, 2\})$$

Proof. If a tuple $(x_1, \ldots, x_{16}) \in (\mathbb{N} \setminus \{0\})^{16}$ solves the system C and

$$(x_4 \in \{1, 2\}) \lor (x_9 \in \{1, 2\})$$

then $x_1, \ldots, x_{16} \le 7!$. Indeed, for example, if $x_4 = 2$ then $x_6 = x_4 + 1 = 3$. Hence, $x_7 = x_6! = 6$. Therefore, $x_{15} = x_7 + 1 = 7$. Consequently, $x_{16} = x_{15}! = 7!$.

Theorem 10. The statement Ψ_{16} proves the following implication: (*) if there exists a twin prime greater than g(14), then there are infinitely many twin primes.

Proof. Suppose that the antecedent holds. Then, there exist prime numbers x_4 and x_9 such that $x_9 = x_4 + 2 > g(14)$. Hence, $x_4, x_9 \in \mathbb{N} \setminus \{0, 1, 2\}$. By Lemma 11, there exists a unique tuple $(x_1, x_2, x_3, x_5, x_6, x_7, x_8, x_{10}, x_{11}, x_{12}, x_{13}, x_{14}, x_{15}, x_{16}) \in (\mathbb{N} \setminus \{0\})^{14}$ such that the tuple (x_1, \dots, x_{16}) solves the system *C*. Since $x_9 > g(14)$, we obtain that $x_9 - 1 \ge g(14)$. Therefore, $(x_9 - 1)! \ge g(14)! = g(15)$. Hence, $(x_9 - 1)! + 1 > g(15)$. Consequently,

$$x_{16} = ((x_9 - 1)! + 1)! > g(15)! = g(16)$$

Since $C \subseteq B_{16}$, the statement Ψ_{16} and the inequality $x_{16} > g(16)$ imply that the system C has infinitely many solutions in positive integers x_1, \ldots, x_{16} . According to Lemmas 11 and 12, there are infinitely many twin primes.

Corollary 5. Let X denote the set of twin primes. The statement Ψ_{16} implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. In order to prove item (a), we compute $\max(X \cap [1, g(14)])$ and consider the algorithm which decides the inequality $\max(X \cap [1, g(14)]) < n$. In order to prove items (b) and (c), we consider the algorithm which computes $\max(X \cap [1, g(14)])$.

9 Hypothetical statements $\Delta_5, \ldots, \Delta_{14}$ about the Gamma function and their consequences

Let $\lambda(5) = \Gamma(25)$, and let $\lambda(n+1) = \Gamma(\lambda(n))$ for every integer $n \ge 5$. For an integer $n \ge 5$, let \mathcal{J}_n denote the following system of equations:

$$\begin{cases}
\forall i \in \{1, \dots, n-1\} \setminus \{3\} \Gamma(x_i) = x_{i+1} \\
x_1 \cdot x_1 = x_4 \\
x_2 \cdot x_3 = x_5
\end{cases}$$

Lemma 2 and the diagram in Figure 5 explain the construction of the system \mathcal{J}_n .

On sets $X \subseteq \mathbb{N}$ whose finiteness ...

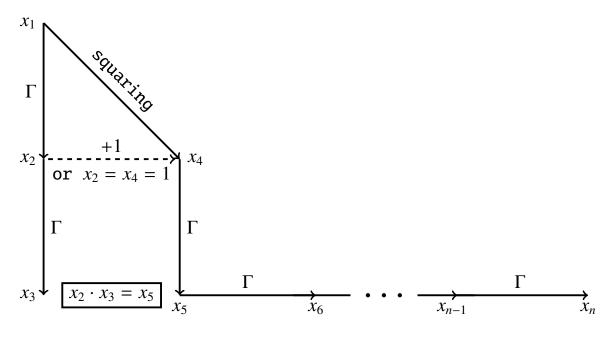


Fig. 5 Construction of the system \mathcal{J}_n

For every integer $n \ge 5$, the system \mathcal{J}_n has exactly two solutions in positive integers, namely $(1, \ldots, 1)$ and $(5, 24, 23!, 25, \lambda(5), \ldots, \lambda(n))$. For an integer $n \ge 5$, let Δ_n denote the following statement: if a system $S \subseteq \{\Gamma(x_i) = x_k : i, k \in \{1, \ldots, n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \ldots, n\}\}$ has only finitely many solutions in positive integers x_1, \ldots, x_n , then each such solution (x_1, \ldots, x_n) satisfies $x_1, \ldots, x_n \le \lambda(n)$.

Hypothesis 2. The statements $\Delta_5, \ldots, \Delta_{14}$ are true.

Lemmas 2 and 4 imply that the statements Δ_n have similar consequences as the statements Ψ_n .

Theorem 11. The statement Δ_6 implies that any prime number $p \ge 25$ proves the infinitude of primes.

Proof. It follows from Lemmas 2 and 4. We leave the details to the reader.

10 Hypothetical statements $\Sigma_3, \ldots, \Sigma_{16}$ about the Gamma function and their consequences

Let $\Gamma_n(k)$ denote (k-1)!, where $n \in \{3, ..., 16\}$ and $k \in \{2\} \cup [2^{2^{n-3}} + 1, \infty) \cap \mathbb{N}$. For an integer $n \in \{3, ..., 16\}$, let

$$Q_n = \{\Gamma_n(x_i) = x_k : i, k \in \{1, \dots, n\}\} \cup \{x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\}\}$$

For an integer $n \in \{3, ..., 16\}$, let P_n denote the following system of equations:

$$\begin{cases} x_1 \cdot x_1 &= x_1 \\ \Gamma_n(x_2) &= x_1 \end{cases}$$

$$\forall i \in \{2, \dots, n-1\} \ x_i \cdot x_i &= x_{i+1} \end{cases}$$

Lemma 13. For every integer $n \in \{3, ..., 16\}$, $P_n \subseteq Q_n$ and the system P_n with Γ instead of Γ_n has exactly one solution in positive integers $x_1, ..., x_n$, namely $\left(1, 2^{2^0}, 2^{2^1}, 2^{2^2}, ..., 2^{2^{n-2}}\right)$.

For an integer $n \in \{3, ..., 16\}$, let Σ_n denote the following statement: if a system of equations $S \subseteq Q_n$ with Γ instead of Γ_n has only finitely many solutions in positive integers $x_1, ..., x_n$, then every tuple $(x_1, ..., x_n) \in (\mathbb{N} \setminus \{0\})^n$ that solves the original system S satisfies $x_1, ..., x_n \leq 2^{2^{n-2}}$.

Hypothesis 3. The statements $\Sigma_3, \ldots, \Sigma_{16}$ are true.

Lemma 14. (cf. Lemma 2). For every integer $n \in \{4, ..., 16\}$ and for every positive integers x and y, $x \cdot \Gamma_n(x) = \Gamma_n(y)$ if and only if $(x + 1 = y) \land (x \ge 2^{2^{n-3}} + 1)$.

Let $\mathbb{Z}_9 \subseteq \mathbb{Q}_9$ be the system of equations in Figure 6.

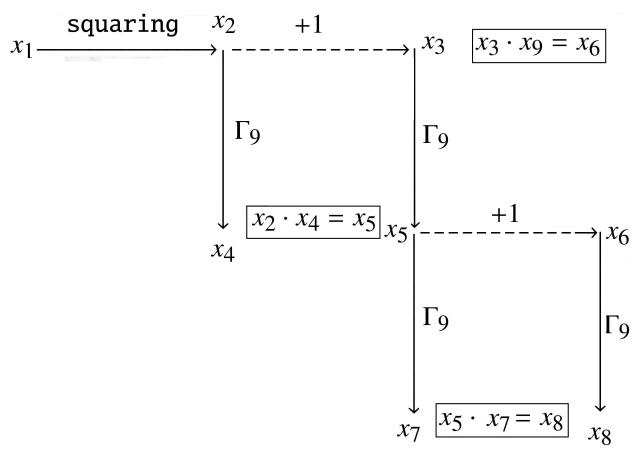


Fig. 6 Construction of the system \mathbb{Z}_9

Lemma 15. For every positive integer x_1 , the system \mathbb{Z}_9 is solvable in positive integers x_2, \ldots, x_9 if and only if $x_1 > 2^{2^{9-4}}$ and $x_1^2 + 1$ is prime. In this case, positive integers x_2, \ldots, x_9 are uniquely determined by x_1 . For every positive integer n, at most finitely many tuples $(x_1, \ldots, x_9) \in (\mathbb{N} \setminus \{0\})^9$ begin with n and solve the system \mathbb{Z}_9 with Γ instead of Γ_9 .

Proof. It follows from Lemmas 2, 4, and 14.

Lemma 16. ([13]). The number $(13!)^2 + 1 = 38775788043632640001$ is prime.

Lemma 17.
$$((13!)^2 \ge 2^{2^{9-3}} + 1 = 18446744073709551617) \land (\Gamma_9((13!)^2) > 2^{2^{9-2}}).$$

Theorem 12. The statement Σ_9 implies the infinitude of primes of the form $n^2 + 1$.

Proof. It follows from Lemmas 15–17.

Theorem 13. (cf. Theorem 9). The statement Σ_9 implies that any prime of the form n! + 1 with $n \ge 2^{2^{9-3}}$ proves the infinitude of primes of the form n! + 1.

Proof. We leave the proof to the reader.

On sets $X \subseteq \mathbb{N}$ whose finiteness ...

Corollary 6. Let X denote the set of primes of the form n! + 1. The statement Σ_9 implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. In order to prove item (a), we compute $\max(X \cap [1, (2^{2^{9-3}} - 1)! + 1])$ and consider the algorithm which decides the inequality $\max(X \cap [1, (2^{2^{9-3}} - 1)! + 1]) < n$. In order to prove items (b) and (c), we consider the algorithm which computes $\max(X \cap [1, (2^{2^{9-3}} - 1)! + 1])$.

Let $\mathcal{Z}_{14} \subseteq Q_{14}$ be the system of equations in Figure 7.

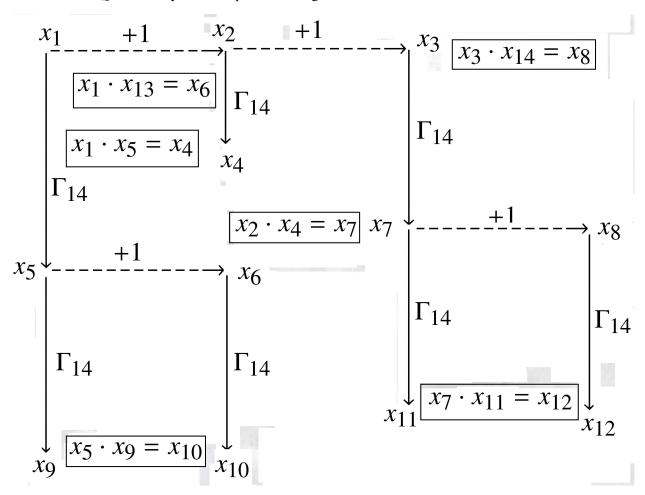


Fig. 7 Construction of the system Z_{14}

Lemma 18. For every positive integer x_1 , the system \mathbb{Z}_{14} is solvable in positive integers x_2, \ldots, x_{14} if and only if x_1 and $x_1 + 2$ are prime and $x_1 \ge 2^{2^{14-3}} + 1$. In this case, positive integers x_2, \ldots, x_{14} are uniquely determined by x_1 . For every positive integer n, at most finitely many tuples $(x_1, \ldots, x_{14}) \in (\mathbb{N} \setminus \{0\})^{14}$ begin with n and solve the system \mathbb{Z}_{14} with Γ instead of Γ_{14} .

Proof. It follows from Lemmas 2, 4, and 14.

Lemma 19. ([16, p. 87]). The numbers $459 \cdot 2^{8529} - 1$ and $459 \cdot 2^{8529} + 1$ are prime (Harvey Dubner). **Lemma 20.** $459 \cdot 2^{8529} - 1 > 2^{2^{14-2}} = 2^{4096}$.

Theorem 14. The statement Σ_{14} implies the infinitude of twin primes.

Proof. It follows from Lemmas 18–20.

On sets $X \subseteq \mathbb{N}$ whose finiteness ...

A prime p is said to be a Sophie Germain prime if both p and 2p + 1 are prime, see [15]. Let $\mathcal{Z}_{16} \subseteq Q_{16}$ be the system of equations in Figure 8.

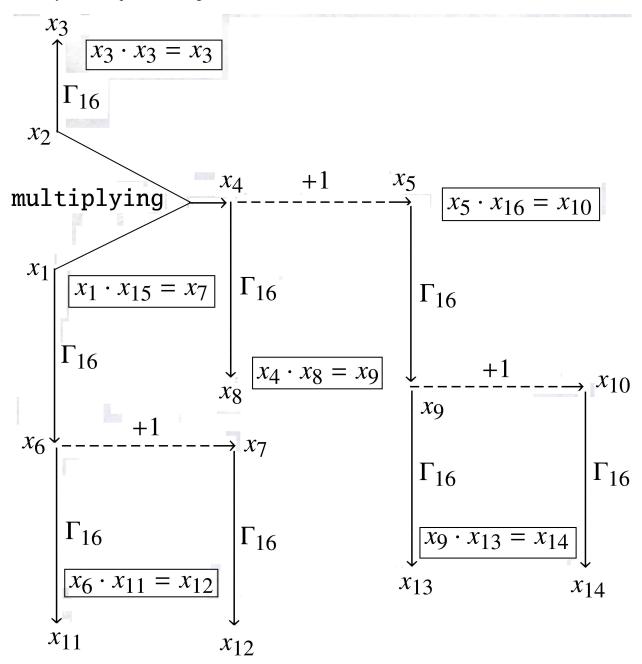


Fig. 8 Construction of the system Z_{16}

Lemma 21. For every positive integer x_1 , the system \mathbb{Z}_{16} is solvable in positive integers x_2, \ldots, x_{16} if and only if x_1 is a Sophie Germain prime and $x_1 \ge 2^{2^{16-3}} + 1$. In this case, positive integers x_2, \ldots, x_{16} are uniquely determined by x_1 . For every positive integer n, at most finitely many tuples $(x_1, \ldots, x_{16}) \in (\mathbb{N} \setminus \{0\})^{16}$ begin with n and solve the system \mathbb{Z}_{16} with Γ instead of Γ_{16} .

Proof. It follows from Lemmas 2, 4, and 14.

Lemma 22. ([10, p. 330]). $8069496435 \cdot 10^{5072} - 1$ is a Sophie Germain prime (Harvey Dubner).

Lemma 23. $8069496435 \cdot 10^{5072} - 1 > 2^{2^{16-2}}$.

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Theorem 15. The statement Σ_{16} implies the infinitude of Sophie Germain primes.

Proof. It follows from Lemmas 21–23.

Theorem 16. The statement Σ_6 proves the following implication: if the equation x(x+1) = y! has only finitely many solutions in positive integers x and y, then each such solution (x,y) belongs to the set $\{(1,2),(2,3)\}$.

Proof. We leave the proof to the reader.

Corollary 7. Let X denote the set $\{y \in \mathbb{N} \setminus \{0\} : y! \text{ is a product of two consecutive positive integers}\}$. The statement Σ_6 implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. In order to prove item (a), we consider the algorithm which decides the inequality 3 < n. In order to prove items (b) and (c), we consider the algorithm which returns 3.

The question of solving the equation x(x + 1) = y! was posed by P. Erdös, see [1]. F. Luca proved that the *abc* conjecture implies that the equation x(x + 1) = y! has only finitely many solutions in positive integers, see [7].

Theorem 17. The statement Σ_6 proves the following implication: if the equation $x! + 1 = y^2$ has only finitely many solutions in positive integers x and y, then each such solution (x, y) belongs to the set $\{(4,5), (5,11), (7,71)\}$.

Proof. We leave the proof to the reader.

Corollary 8. Let X denote the set $\{x \in \mathbb{N} \setminus \{0\} : x! + 1 \text{ is a perfect square}\}$. The statement Σ_6 implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. We repeat the proof of Corollary 2.

11 Are there infinitely many composite Fermat numbers?

Integers of the form $2^{2^n} + 1$ are called Fermat numbers. Primes of the form $2^{2^n} + 1$ are called Fermat primes, as Fermat conjectured that every integer of the form $2^{2^n} + 1$ is prime, see [6, p. 1]. Fermat correctly remarked that $2^{2^0} + 1 = 3$, $2^{2^1} + 1 = 5$, $2^{2^2} + 1 = 17$, $2^{2^3} + 1 = 257$, and $2^{2^4} + 1 = 65537$ are all prime, see [6, p. 1].

Open Problem. ([6, p. 159]). Are there infinitely many composite numbers of the form $2^{2^n} + 1$? Most mathematicians believe that $2^{2^n} + 1$ is composite for every integer $n \ge 5$, see [5, p. 23]. Let

$$H_n = \left\{ x_i \cdot x_j = x_k : i, j, k \in \{1, \dots, n\} \right\} \cup \left\{ 2^{2^{x_i}} = x_k : i, k \in \{1, \dots, n\} \right\}$$

Let h(1) = 1, and let $h(n + 1) = 2^{2h(n)}$ for every positive integer n.

Lemma 24. The following subsystem of H_n

$$\begin{cases} x_1 \cdot x_1 &= x_1 \\ \forall i \in \{1, \dots, n-1\} \ 2^{2^{x_i}} &= x_{i+1} \end{cases}$$

has exactly one solution $(x_1, \ldots, x_n) \in (\mathbb{N} \setminus \{0\})^n$, namely $(h(1), \ldots, h(n))$.

For a positive integer n, let ξ_n denote the following statement: if a system $S \subseteq H_n$ has only finitely many solutions in positive integers x_1, \ldots, x_n , then each such solution (x_1, \ldots, x_n) satisfies $x_1, \ldots, x_n \le h(n)$. The statement ξ_n says that for subsystems of H_n the largest known solution is indeed the largest possible.

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Hypothesis 4. The statements ξ_1, \ldots, ξ_{13} are true.

Theorem 18. Every statement ξ_n is true with an unknown integer bound that depends on n.

Proof. For every positive integer n, the system H_n has a finite number of subsystems.

Theorem 19. The statement ξ_{13} proves the following implication: if $z \in \mathbb{N} \setminus \{0\}$ and $2^{2^z} + 1$ is composite and greater than h(12), then $2^{2^z} + 1$ is composite for infinitely many positive integers z.

Proof. Let us consider the equation

$$(x+1)(y+1) = 2^{2^{z}} + 1$$
 (1)

in positive integers. By Lemma 3, we can transform equation (1) into an equivalent system \mathcal{G} which has 13 variables (x, y, z, and 10 other variables) and which consists of equations of the forms $\alpha \cdot \beta = \gamma$ and $2^{2^{\alpha}} = \gamma$, see the diagram in Figure 9.

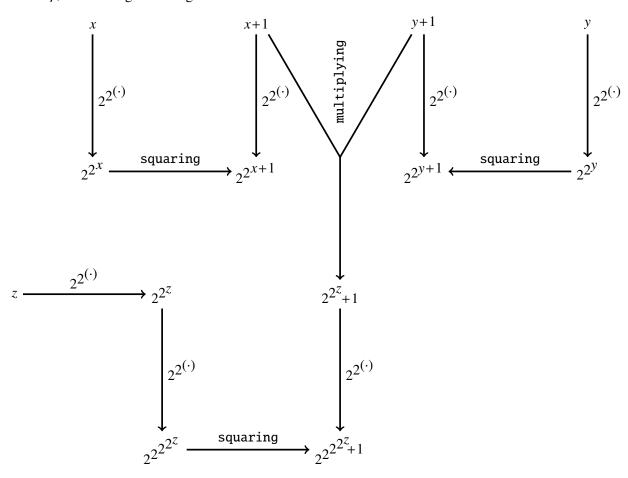


Fig. 9 Construction of the system G

Since $2^{2^{\mathcal{Z}}} + 1 > h(12)$, we obtain that $2^{2^{2^{\mathcal{Z}}} + 1} > h(13)$. By this, the statement ξ_{13} implies that the system \mathcal{G} has infinitely many solutions in positive integers. It means that there are infinitely many composite Fermat numbers.

Corollary 9. Let X denote the set of composite Fermat numbers. The statement ξ_{13} implies the following claims: (a) X is uncomplicated, (b) we know an algorithm which returns a threshold number of X, (c) this algorithm returns $\max(X)$, if X is finite.

Proof. In order to prove item (a), we compute $\max(X \cap [1, h(12)])$ and consider the algorithm which decides the inequality $\max(X \cap [1, h(12)]) < n$. In order to prove items (b) and (c), we consider the algorithm which computes $\max(X \cap [1, h(12)])$.

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