

Complexity of some of Pyramid Graphs Created from a Gear Graph

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Abstract

In mathematics, one always aims to obtain new frameworks from specific ones. This also stratified to the regality of graphs, where one can produce numerous new graphs from a specific set of graphs. In this work we define some classes of pyramid graphs created by a gear graph and we derive straightforward formulas of the complexity of these graphs, using linear algebra matrix analysis techniques and employing knowledges of Chebyshev polynomials.

Keywords: Complexity; Chebyshev Polynomials; Gear graph; Pyramid graphs.

Mathematics Subject Classification: 05C05, 05C50.

1. Introduction

The graph theory is a theory gathering computer science and mathematics, permitting to solve considerable problems in several fields (telecom, social network, molecules, computer network, genetics,...) by designing them by graphs and facilitate them by idealistic cases such as the spanning trees See [1-10].

A spanning tree of a finite connected graph G is a maximal subset of the edges that contains no cycle, equivalently a minimal subset of the edges that connects all the vertices. The history of enumerating number of spanning trees $\tau(G)$ of a graph G dates back into year 1842 in which the physicist of Kirchhoff [11] gave the matrix tree theorem established on the determinants of a certain matrix gained from the Laplacian matrix L defined by the difference between the degree matrix D and adjacency matrix A of a graph G . That is

$$L_{ij} = \begin{cases} a_i & \text{if } i = j \\ -1 & \text{if } i \neq j \text{ and } i \text{ is adjacent to } j, \\ 0 & \text{otherwise} \end{cases}$$

where a_i denotes the degree of vertex i .

This method allows beneficial results for a graph comprising a small number of vertices, but it will be infeasible for large graph. There are one more methods for calculating $\tau(G)$. Let $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_k = 0$ denote the eigenvalues of the matrix L of a graph G with n vertices. Kelmans and Chelnokov [12] have derived that

$$\tau(G) = \frac{1}{k} \prod_{i=1}^{k-1} \lambda_k.$$

Many works have conceived techniques to derive the number of spanning tree of a graph. Now, we give the following Lemma:

Lemma 1.1[13]

$\tau(G) = \frac{1}{k^2} \det(kI - D^c + A^c)$ where A^c and D^c are the adjacency and degree matrices of G^c , the complement of G , respectively, and I is the $k \times k$ identity matrix.

The characteristic of this formula is to express $\tau(G)$ straightway as a determinant rather than in terms of cofactors as in Kirchhoff theorem or eigenvalues as in Kelmans and Chelnokov formula.

2. Chebyshev Polynomial

In this part we insert some relations regarding Chebyshev polynomials of the first and second types which we use it in our calculations.

We start from their definitions, see Yuanping, et. al. [14].

Let $A_n(x)$ be $n \times n$ matrix such that

$$A_n(x) = \begin{pmatrix} 2x & -1 & 0 & \cdots & 0 \\ -1 & 2x & -1 & \ddots & \vdots \\ 0 & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & -1 \\ 0 & \cdots & 0 & -1 & 2x \end{pmatrix}.$$

Furthermore, we rendering that the Chebyshev polynomials of the first type are defined by

$$T_n(x) = \cos(n \cos^{-1} x) \quad (1)$$

The Chebyshev polynomials of the second type are defined by

$$U_{n-1}(x) = \frac{1}{n} \frac{d}{dx} T_n(x) = \frac{\sin(n \cos^{-1} x)}{\sin(\cos^{-1} x)} \quad (2)$$

It is easily confirmed that

$$U_n(x) - 2xU_{n-1}(x) + U_{n-2}(x) = 0 \quad (3)$$

It can then be shown from this recursion that by expanding $\det A_n(x)$ one obtains

$$U_n(x) = \det(A_n(x)), n \geq 1 \quad (4)$$

Moreover by solving the recursion (3), one gets the straightforward formula

$$U_n(x) = \frac{(x + \sqrt{x^2 - 1})^{n+1} - (x - \sqrt{x^2 - 1})^{n+1}}{2\sqrt{x^2 - 1}}, n \geq 1, \quad (5)$$

where the conformity is valid for all complex x (except at $x = \pm 1$, where the function can be taken as the limit).

The definition of $U_n(x)$ easily yields its zeros and it can therefore be confirmed that

$$U_{n-1}(x) = 2^{n-1} \prod_{j=1}^{n-1} (x - \cos \frac{j\pi}{n}) \quad (6)$$

One further notes that

$$U_{n-1}(-x) = (-1)^{n-1} U_{n-1}(x) \quad (7)$$

From Eqs. (6) and (7), we have:

$$U_{n-1}^2(x) = 4^{n-1} \prod_{j=1}^{n-1} (x^2 - \cos^2 \frac{j\pi}{n}) \quad (8)$$

Finally, straightforward manipulation of the above formula gets the following formula (9), which is highly beneficial to us latter:

$$U_{n-1}^2(\sqrt{\frac{x+2}{4}}) = \prod_{j=1}^{n-1} (x - 2 \cos \frac{2j\pi}{n}) \quad (9)$$

Moreover, one can see that

$$U_{n-1}^2(x) = \frac{1 - T_{2n}}{2(1 - x^2)} = \frac{1 - T_n(2x^2 - 1)}{2(1 - x^2)}, \quad (10)$$

$$T_n(x) = \frac{1}{2} x^2 [(\sqrt{1 - \frac{1}{x^2}} + 1)^n + (\sqrt{1 - \frac{1}{x^2}})^n]. \quad (11)$$

Now we introduce the following important two Lemmas.

Lemma 2.1 [15] Let $B_n(x)$ be $n \times n$ Circulant matrix such that

$$B_n(x) = \begin{pmatrix} x & 0 & 1 & \cdots & 1 & 0 \\ 0 & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & x \end{pmatrix}.$$

Then for $n \geq 3, x \geq 4$, one has

$$\det(B_n(x)) = \frac{2(x+n-3)}{x-3} [T_n(\frac{x-1}{2}) - 1].$$

Lemma 2.2 [16]

If $A \in F^{n \times n}$, $B \in F^{n \times m}$, $C \in F^{m \times n}$ and $D \in F^{m \times m}$. Suppose that A and D are nonsingular matrices, then:

$$\det \begin{pmatrix} A & B \\ C & D \end{pmatrix} = \det(A - BD^{-1}C) \det D = \det A \det(D - CA^{-1}B).$$

This Lemma give a type of symmetry for some matrices which simplify our calculations of the complexity of graphs studied in this paper.

3. Main Results

Definition 3.1 The pyramid graph $A_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}$, $\{v_1^2, v_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i = 1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j = 1, 2, \dots, m - 1$, and v_i^m is adjacent to u_1 and u_m . See Fig. (1).

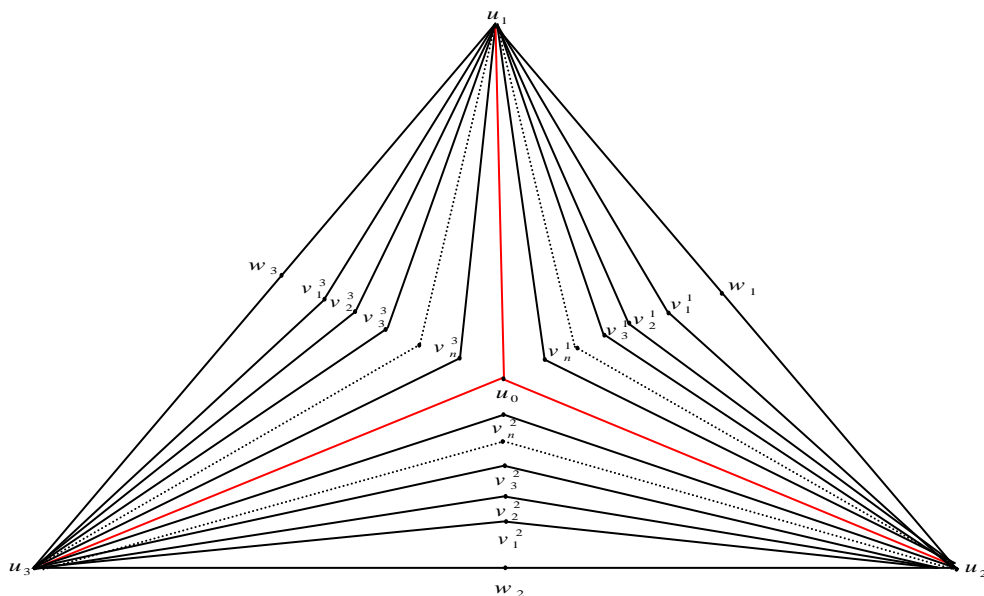


Fig. 1 The pyramid graph $A_n^{(3)}$

Theorem 1 For $n \geq 0, m \geq 3$, $\tau(A_n^{(m)}) = 2^{mn} [(n + 2 + \sqrt{2n + 3})^m + (n + 2 - \sqrt{2n + 3})^m - 2(n + 1)^m]$.

Proof. Using Lemma 1.1, we have

$$\tau(A_n^{(m)}) = \frac{1}{(mn + 2m + 1)^2} \times \det((mn + 2m + 1)I - D^c + A^c) = \frac{1}{(mn + 2m + 1)^2} \times$$

$$= \frac{1}{b} \times \det \begin{pmatrix} a & 1 & \dots & \dots & \dots & 1 & -1 & 0 & \dots & \dots & 0 & -1 & -j & 0 & \dots & \dots & 0 & -j \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots & -1 & -1 & \ddots & \ddots & \ddots & 0 & -j & \ddots & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & -1 & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 1 & \dots & \dots & \dots & 1 & a & 0 & \dots & \dots & 0 & -1 & -1 & 0 & \dots & \dots & 0 & -j & -j \\ 0 & 0 & 1 & \dots & \dots & 1 & 2 & 0 & \dots & \dots & \dots & 0 & 0 & \dots & \dots & \dots & \dots & 0 \\ 1 & 0 & \ddots & \ddots & \ddots & 1 & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 & 0 & \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & 1 & \dots & \dots & 1 & 0 & 0 & \dots & \dots & \dots & 0 & 2 & 0 & \dots & \dots & \dots & \dots & 0 \\ 0 & 0 & j^t & \dots & \dots & j^t & 0 & \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ j^t & 0 & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & j^t & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ j^t & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & j^t & \dots & \dots & j^t & 0 & 0 & \dots & \dots & \dots & \dots & 0 & \dots & \dots & \dots & \dots & \dots & 0 \end{pmatrix}$$

Using Lemma 2.2, yields

$$\tau(A_n^{(m)}) = \frac{1}{b} \times \det \begin{pmatrix} A & B \\ C & 2I_{mn} \end{pmatrix} = \frac{1}{b} \times \det(A - B \frac{1}{2I_{mn}} C) \times 2^{mn}$$

$$= \frac{1}{b} 2^{mn} \times 2^{-2m} \times \det \begin{pmatrix} 2a & n+2 & 2(n+1) & \dots & 2(n+1) & n+2 & -2 & 0 & \dots & \dots & 0 & -2 \\ n+2 & 2a & n+2 & 2(n+1) & \dots & 2(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 2(n+1) & n+2 & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2(n+1) & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2(n+1) & \ddots & \ddots & \ddots & \ddots & n+2 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ n+2 & 2(n+1) & \dots & 2(n+1) & n+2 & 2a & 0 & \dots & \dots & 0 & -2 & -2 \\ 0 & 0 & 2 & \dots & \dots & 2 & 4 & 0 & \dots & \dots & \dots & 0 \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 2 & \dots & \dots & 2 & 0 & 0 & \dots & \dots & \dots & 0 & 4 \end{pmatrix}$$

Using Lemma 2.2 again, yields

$$\tau(A_n^{(m)}) = \frac{2^{m(n-2m)}}{b} \times \det \begin{pmatrix} D & E \\ F & 4I_m \end{pmatrix} = \frac{2^{mn}}{b} \times \det(D - E \frac{1}{4I_m} F)$$

$$\tau(A_n^{(m)}) = \frac{2^{mn}}{b} \times \det \begin{pmatrix} 2a & (n+3) & 2(n+2) & \dots & 2(n+2) & (n+3) \\ (n+3) & 2a & (n+3) & \ddots & \dots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & (n+3) \\ (n+3) & 2(n+2) & \dots & 2(n+2) & (n+3) & 2a \end{pmatrix}$$

Straightforward inducement using properties of determinants, one can obtain

$$\tau(A_n^{(m)}) = \frac{2^{mn}}{b} \times \frac{2b}{m n + m + 2} \times \det \begin{pmatrix} (2a-n-3) & 0 & (n+1) & \cdots & (n+1) & 0 \\ 0 & (2a-n-3) & 0 & \ddots & \cdots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+1) & \cdots & (n+1) & 0 & (2a-n-3) \end{pmatrix}$$

$$= \frac{2^{m n + 1} (n + 1)^m}{m n + m + 2} \times \det \begin{pmatrix} \frac{(2a-n-3)}{(n+1)} & 0 & 1 & \cdots & 1 & 0 \\ 0 & \frac{(2a-n-3)}{(n+1)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & \frac{(2a-n-3)}{(n+1)} \end{pmatrix}$$

Using Lemma 2.1, yields

$$\tau(A_n^{(m)}) = 2^{m n + 1} \times \frac{(n + 1)^m}{m n + m + 2} \times \frac{2(\frac{2a-n-3}{n+1} + m - 3)}{\frac{2a-n-3}{n+1} - 3} \times [T_m(\frac{\frac{2a-n-3}{n+1} - 1}{2}) - 1]$$

$$= 2^{m n + 1} \times (n + 1)^m \times [T_m(\frac{n + 2}{n + 1}) - 1].$$

Using Equation (11), yields the result. □

Definition 3.2 The pyramid graph $B_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ with double internal edges and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}, \{v_1^2, v_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i = 1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j = 1, 2, \dots, m - 1$, and v_i^m is adjacent to u_1 and u_m . See Fig. 2.

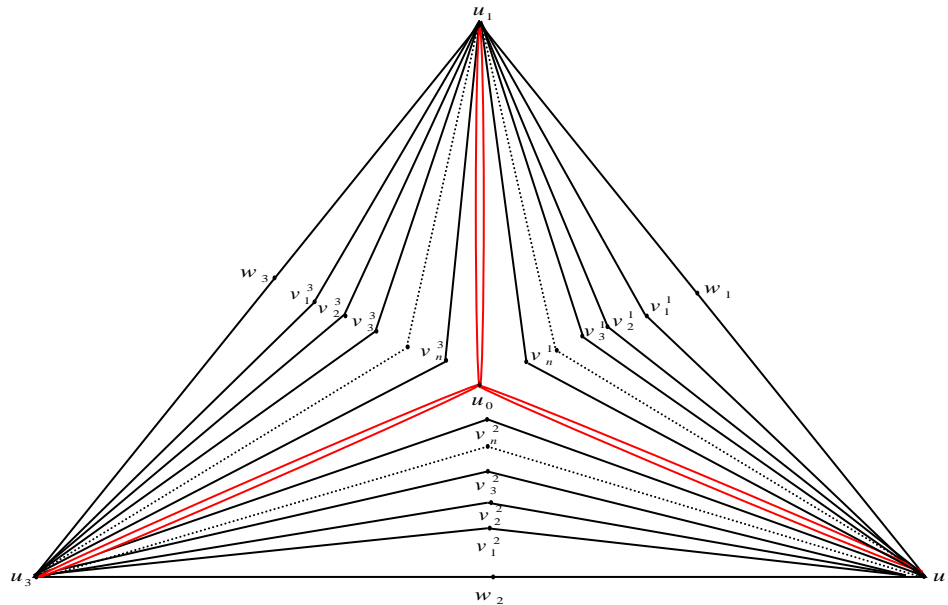


Fig. 2 The pyramid graph $B_n^{(3)}$

$$= \frac{1}{b} \times \det \begin{pmatrix} (a+1) & 2 & \dots & \dots & \dots & 2 & -1 & 0 & \dots & \dots & 0 & -1 & -j & 0 & \dots & \dots & 0 & -j \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & -1 & -1 & \ddots & \ddots & \ddots & 0 & -j & \ddots & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & -1 & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 2 & \dots & \dots & \dots & 2 & (a+1) & 0 & \dots & \dots & 0 & -1 & -1 & 0 & \dots & \dots & 0 & -j & -j \\ 1 & 1 & 2 & \dots & \dots & 2 & 2 & 0 & \dots & \dots & \dots & 0 & 0 & \dots & \dots & \dots & \dots & 0 \\ 2 & 1 & \ddots & \ddots & \ddots & 2 & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 1 & 2 & \dots & \dots & 2 & 1 & 0 & \dots & \dots & \dots & 0 & 2 & 0 & \dots & \dots & \dots & \dots & 0 \\ j' & j' & 2j' & \dots & \dots & 2j' & 0 & \dots & \dots & \dots & \dots & 0 & \dots & \dots & \dots & \dots & \dots & 0 \\ 2j' & j' & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2j' & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2j' & \ddots & \ddots & \ddots & \ddots & j' & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ j' & 2j' & \dots & \dots & 2j' & j' & 0 & \dots & \dots & \dots & \dots & 0 & \dots & \dots & \dots & \dots & \dots & 0 \end{pmatrix}$$

Using Lemma 2.2, yields

$$\tau(B_n^{(m)}) = \frac{1}{b} \times \det \begin{pmatrix} A & B \\ C & 2I_{mn} \end{pmatrix} = \frac{1}{b} \times \det(A - B \frac{1}{2I_{mn}} C) \times 2^{mn}$$

$$= \frac{1}{b} 2^{mn} \times 2^{-2m} \times \det \begin{pmatrix} (2a+2n+2) & 3n+4 & 4(n+1) & \dots & 4(n+1) & 3n+4 & -2 & 0 & \dots & \dots & 0 & -2 \\ 3n+4 & (2a+2n+2) & 3n+4 & 4(n+1) & \dots & 4(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 4(n+1) & 3n+4 & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 4(n+1) & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 4(n+1) & \ddots & \ddots & \ddots & \ddots & 3n+4 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 3n+4 & 4(n+1) & \dots & 4(n+1) & 3n+4 & (2a+2n+2) & 0 & \dots & \dots & 0 & -2 & -2 \\ 2 & 2 & 4 & \dots & \dots & 4 & 4 & 0 & \dots & \dots & \dots & 0 \\ 4 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 4 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 4 & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 2 & 4 & \dots & \dots & 4 & 2 & 0 & \dots & \dots & \dots & 0 & 4 \end{pmatrix}$$

Using Lemma 2.2 again, yields

$$\tau(B_n^{(m)}) = \frac{2^{mn-2m}}{b} \times \det \begin{pmatrix} D & E \\ F & 4I_m \end{pmatrix} = \frac{2^{mn}}{b} \times \det(D - E \frac{1}{4I_m} F)$$

$$\tau(B_n^{(m)}) = \frac{2^{mn}}{b} \times \det \begin{pmatrix} (2a+2n+4) & (3n+7) & 4(n+2) & \dots & 4(n+2) & (3n+7) \\ (3n+7) & (2a+2n+4) & (3n+7) & \ddots & \dots & 4(n+2) \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 4(n+2) \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & (3n+7) \\ (3n+7) & 4(n+2) & \dots & 4(n+2) & (3n+7) & (2a+2n+4) \end{pmatrix}$$

Straightforward inducement using properties of determinants, we obtain

$$\tau(B_n^{(m)}) = \frac{2^{mn}}{b} \times \frac{4b}{mn+m+4} \times \det \begin{pmatrix} (2a-n-3) & 0 & (n+1) & \cdots & (n+1) & 0 \\ 0 & (2a-n-3) & 0 & \ddots & \cdots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+1) \\ (n+1) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+1) & \cdots & (n+1) & 0 & (2a-n-3) \end{pmatrix}$$

$$= \frac{2^{mn+2} \times (n+1)^m}{mn+m+4} \times \det \begin{pmatrix} \frac{(2a-n-3)}{(n+1)} & 0 & 1 & \cdots & 1 & 0 \\ 0 & \frac{(2a-n-3)}{(n+1)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & \frac{(2a-n-3)}{(n+1)} \end{pmatrix}$$

Using Lemma 2.1, yields

$$\tau(B_n^{(m)}) = 2^{mn+2} \times \frac{(n+1)^m}{mn+m+4} \times \frac{2(\frac{2a-n-3}{n+1} + m - 3)}{\frac{2a-n-3}{n+1} - 3} \times [T_m(\frac{\frac{2a-n-3}{n+1} - 1}{2}) - 1]$$

$$= 2^{mn+1} \times (n+1)^m \times [T_m(\frac{n+3}{n+1}) - 1].$$

Using Equation (11), yields the result. □

Definition 3.3 The pyramid graph $C_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ with double external edges and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}, \{v_1^2, v_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i = 1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j = 1, 2, \dots, m-1$, and v_i^m is adjacent to u_1 and u_m . See Fig. 3.

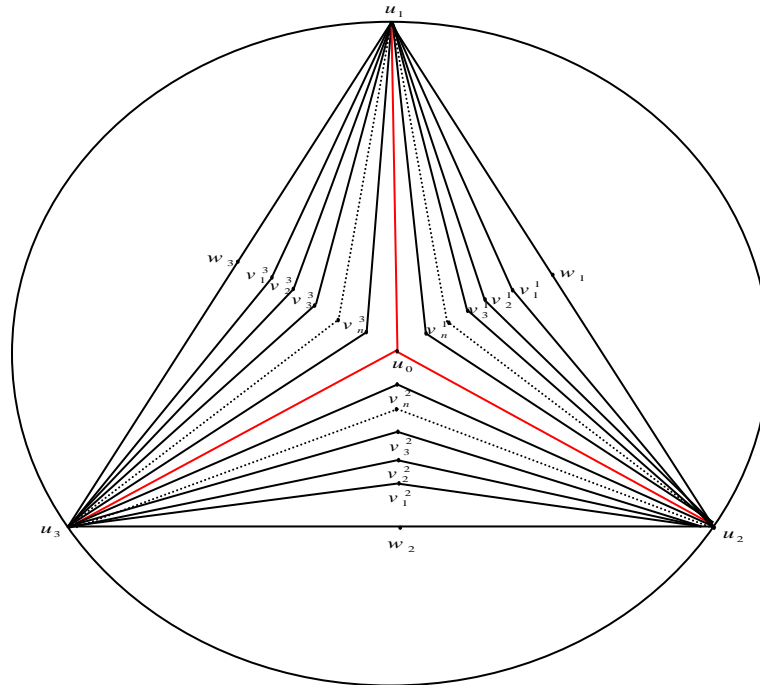


Fig. 3. The pyramid graph $C_n^{(3)}$

$$= \frac{1}{b} \times \det \begin{pmatrix} a & 0 & 1 & \dots & 1 & 0 & -1 & 0 & \dots & \dots & 0 & -1 & -j & 0 & \dots & \dots & 0 & -j \\ 0 & \ddots & \ddots & \ddots & \ddots & 1 & -1 & -1 & \ddots & \ddots & \ddots & 0 & -j & \ddots & \ddots & \ddots & \ddots & 0 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & -1 & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \dots & 1 & 0 & a & 0 & \dots & \dots & 0 & -1 & -1 & 0 & \dots & \dots & 0 & -j & -j \\ 0 & 0 & 1 & \dots & \dots & 1 & 2 & 0 & \dots & \dots & \dots & 0 & 0 & \dots & \dots & \dots & \dots & 0 \\ 1 & 0 & \ddots & \ddots & \ddots & 1 & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & 0 & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & 1 & \dots & \dots & 1 & 0 & 0 & \dots & \dots & \dots & 0 & 2 & 0 & \dots & \dots & \dots & \dots & 0 \\ 0 & 0 & j' & \dots & \dots & j' & 0 & \dots & \dots & \dots & \dots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ j' & 0 & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & j' & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ j' & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & j' & \dots & \dots & j' & 0 & 0 & \dots & \dots & \dots & \dots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \end{pmatrix}$$

Using Lemma 2.2, yields

$$\tau(C_n^{(m)}) = \frac{1}{b} \times \det \begin{pmatrix} A & B \\ C & 2I_{mn} \end{pmatrix} = \frac{1}{b} \times \det \left(A - B \frac{1}{2I_{mn}} C \right) \times 2^{mn}$$

$$= \frac{1}{b} \times 2^{mn} \times 2^{-2m} \times \det \begin{pmatrix} 2a & n & 2(n+1) & \dots & 2(n+1) & n & -2 & 0 & \dots & \dots & 0 & -2 \\ n & 2a & n+2 & 2(n+1) & \dots & 2(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 2(n+1) & n & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2(n+1) & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2(n+1) & \ddots & \ddots & \ddots & \ddots & n & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ n & 2(n+1) & \dots & 2(n+1) & n & 2a & 0 & \dots & \dots & 0 & -2 & -2 \\ 0 & 0 & 2 & \dots & \dots & 2 & 4 & 0 & \dots & \dots & \dots & 0 \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 2 & \dots & \dots & 2 & 0 & 0 & \dots & \dots & \dots & 0 & 4 \end{pmatrix}$$

Using Lemma 2.2 again, yields

$$\tau(C_n^{(m)}) = \frac{2^{mn-2m}}{b} \times \det \begin{pmatrix} D & E \\ F & 4I_m \end{pmatrix} = \frac{2^{mn}}{b} \times \det \left(D - E \frac{1}{4I_m} F \right)$$

$$\tau(C_n^{(m)}) = \frac{2^{mn}}{b} \times \det \begin{pmatrix} 2a & (n+1) & 2(n+2) & \dots & 2(n+2) & (n+1) \\ (n+1) & 2a & (n+3) & \ddots & \dots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2(n+2) \\ 2(n+2) & \ddots & \ddots & \ddots & \ddots & (n+1) \\ (n+1) & 2(n+2) & \dots & 2(n+2) & (n+1) & 2a \end{pmatrix}$$

Using properties of determinants, we have:

$$\tau(C_n^{(m)}) = \frac{2^{mn}}{b} \times \frac{2b}{mn+3m+2} \times \det \begin{pmatrix} (2a-n-1) & 0 & (n+3) & \cdots & (n+3) & 0 \\ 0 & (2a-n-1) & 0 & \ddots & \cdots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+3) & \cdots & (n+3) & 0 & (2a-n-1) \end{pmatrix}$$

$$= \frac{2^{mn+1} (n+3)^m}{mn+3m+2} \times \det \begin{pmatrix} \frac{(2a-n-1)}{(n+3)} & 0 & 1 & \cdots & 1 & 0 \\ 0 & \frac{(2a-n-1)}{(n+3)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & \frac{(2a-n-1)}{(n+3)} \end{pmatrix}$$

Using Lemma 2.1, yields:

$$\tau(C_n^{(m)}) = 2^{mn+1} \times \frac{(n+3)^m}{mn+3m+2} \times \frac{2(\frac{2a-n-1}{n+3} + m - 3)}{\frac{2a-n-1}{n+3} - 3} \times [T_m(\frac{\frac{2a-n-1}{n+3} - 1}{2}) - 1]$$

$$= 2^{mn+1} \times (n+3)^m \times [T_m(\frac{n+4}{n+3}) - 1].$$

Using Equation (11), yields the result. □

Definition 3.4 The pyramid graph $D_n^{(m)}$ is the graph created from the gear graph G_{m+1} with vertices $\{u_0; u_1, u_2, \dots, u_m; w_1, w_2, \dots, w_m\}$ with double internal and external edges and m sets of vertices, say, $\{v_1^1, v_2^1, \dots, v_n^1\}$, $\{v_1^2, v_2^2, \dots, v_n^2\}, \dots, \{v_1^m, v_2^m, \dots, v_n^m\}$, such that for all $i=1, 2, \dots, n$ the vertex v_i^j is adjacent to u_j and u_{j+1} , where $j=1, 2, \dots, m-1$, and v_i^m is adjacent to u_1 and u_m . See Fig.(4).

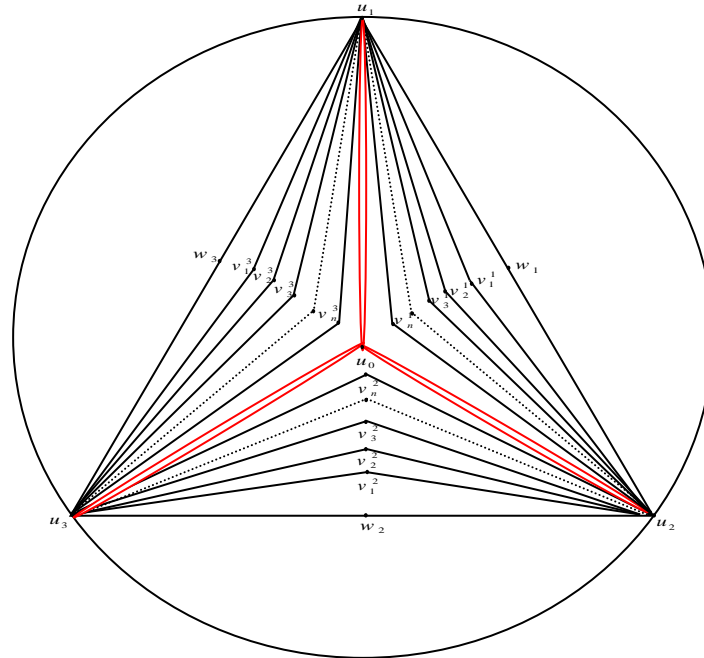


Fig. 4. The pyramid graph $D_n^{(3)}$

$$= \frac{1}{b} \det \begin{pmatrix} (a+1) & 1 & 2 & \dots & 2 & 1 & -1 & 0 & \dots & \dots & 0 & -1 & -j & 0 & \dots & \dots & 0 & -j \\ 1 & \ddots & \ddots & \ddots & \ddots & 2 & -1 & -1 & \ddots & \ddots & \ddots & 0 & -j & \ddots & \ddots & \ddots & \ddots & 0 \\ 2 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & -1 & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ 1 & 2 & \dots & 2 & 1 & (a+1) & 0 & \dots & \dots & 0 & -1 & -1 & 0 & \dots & \dots & 0 & -j & -j \\ 1 & 1 & 2 & \dots & \dots & 2 & 2 & 0 & \dots & \dots & \dots & 0 & 0 & \dots & \dots & \dots & \dots & 0 \\ 2 & 1 & \ddots & \ddots & \ddots & 2 & 0 & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 2 & \ddots & \ddots & \ddots & \ddots & 1 & \vdots & \ddots & \ddots & \ddots & \ddots & 0 & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 1 & 2 & \dots & \dots & 2 & 1 & 0 & \dots & \dots & \dots & 0 & 2 & 0 & \dots & \dots & \dots & \dots & 0 \\ j' & j' & 2j' & \dots & \dots & 2j' & 0 & \dots & \dots & \dots & \dots & 0 & & & & & & \\ 2j' & j' & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & & & & & & \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & & & & & & \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2j' & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & & & & & & \\ 2j' & \ddots & \ddots & \ddots & \ddots & j' & \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & & & & & & \\ j' & 2j' & \dots & \dots & 2j' & j' & 0 & \dots & \dots & \dots & \dots & 0 & & & & & & \end{pmatrix} 2I_{mn}$$

Using Lemma 2.2, yields

$$\tau(D_n^{(m)}) = \frac{1}{b} \times \det \begin{pmatrix} A & B \\ C & 2I_{mn} \end{pmatrix} = \frac{1}{b} \times \det(A - B \frac{1}{2I_{mn}} C) \times 2^{mn}$$

$$= \frac{1}{b} 2^{mn} \times 2^{-2m} \times \det \begin{pmatrix} (2a+2n+2) & 3n+2 & 4(n+1) & \dots & 4(n+1) & 3n+2 & -2 & 0 & \dots & \dots & 0 & -2 \\ 3n+2 & (2a+2n+2) & 3n+2 & 4(n+1) & \dots & 4(n+1) & -2 & \ddots & \ddots & \ddots & \ddots & 0 \\ 4(n+1) & 3n+4 & \ddots & \ddots & \ddots & \ddots & 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & 4(n+1) & \vdots & \ddots & \ddots & \ddots & \vdots \\ 4(n+1) & \ddots & \ddots & \ddots & \ddots & \ddots & 3n+2 & \vdots & \ddots & \ddots & \ddots & 0 \\ 3n+2 & 4(n+1) & \dots & 4(n+1) & 3n+2 & (2a+2n+2) & 0 & \dots & \dots & 0 & -2 & -2 \\ 2 & 2 & 4 & \dots & \dots & 4 & 4 & 0 & \dots & \dots & \dots & 0 \\ 4 & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots & 0 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & 4 & \vdots & \ddots & \ddots & \ddots & \vdots \\ 4 & \ddots & \ddots & \ddots & \ddots & \ddots & 2 & \vdots & \ddots & \ddots & \ddots & 0 \\ 2 & 4 & \dots & \dots & 4 & 2 & 0 & \dots & \dots & \dots & 0 & 4 \end{pmatrix}$$

Using Lemma 2.2, yields

$$\tau(D_n^{(m)}) = \frac{2^{mn-2m}}{b} \times \det \begin{pmatrix} A & B \\ C & 4I_m \end{pmatrix} = \frac{2^{mn}}{b} \times \det(A - B \frac{1}{4I_m} C)$$

$$\tau(D_n^{(m)}) = \frac{2^{mn}}{b} \times \det \begin{pmatrix} (2a+2n+4) & (3n+5) & 4(n+2) & \dots & 4(n+2) & (3n+5) \\ (3n+5) & (2a+2n+4) & (3n+5) & \ddots & \dots & 4(n+2) \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 4(n+2) \\ 4(n+2) & \ddots & \ddots & \ddots & \ddots & (3n+5) \\ (3n+5) & 4(n+2) & \dots & 4(n+2) & (3n+5) & (2a+2n+4) \end{pmatrix}$$

Straightforward inducement using properties of determinants, we get:

$$\tau(D_n^{(m)}) = \frac{2^{mn}}{b} \times \frac{4b}{m n + 3m + 4} \times \det \begin{pmatrix} (2a-n-1) & 0 & (n+3) & \dots & (n+3) & 0 \\ 0 & (2a-n-1) & 0 & \ddots & \dots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & (n+3) \\ (n+3) & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & (n+3) & \dots & (n+3) & 0 & (2a-n-1) \end{pmatrix}$$

$$= \frac{2^{m n+2} (n+3)^m}{m n+3m+4} \times \det \begin{pmatrix} \frac{(2a-n-1)}{(n+3)} & 0 & 1 & \dots & 1 & 0 \\ 0 & \frac{(2a-n-1)}{(n+3)} & 0 & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \dots & 1 & 0 & \frac{(2a-n-1)}{(n+3)} \end{pmatrix}$$

Using Lemma 2.1, yields:

$$\tau(D_n^{(m)}) = 2^{m n+2} \times \frac{(n+3)^m}{m n+3m+4} \times \frac{2^{\binom{2a-n-1}{n+3} + m - 3}}{\frac{2a-n-1}{n+3} - 3} \times [T_m \left(\frac{2a-n-1}{2} \right) - 1] = 2^{m n+1} \times (n+3)^m \times [T_m \left(\frac{n+5}{n+3} \right) - 1].$$

Using Equation (11), yields the result. □

4. Numerical Results

The following table illustrate some values of number of spanning trees of studied pyramid graphs.

Table 4.1. Some values of the number of spanning trees of studied pyramid graphs.

m	n	$\tau(P_n^{(m)})$	$\tau(A_n^{(m)})$	$\tau(B_n^{(m)})$	$\tau(C_n^{(m)})$
3	0	50	196	242	676
3	1	1024	3200	3136	8192
3	2	15488	43264	36992	92416
3	3	200704	524288	409600	991232
3	4	2367488	5914624	4333568	10240000
3	5	26214400	63438848	44302336	102760448

m	n	$\tau(P_n^{(m)})$	$\tau(A_n^{(m)})$	$\tau(B_n^{(m)})$	$\tau(C_n^{(m)})$
4	0	192	1152	1792	6400
4	1	11520	49152	57600	184320
4	2	458752	1638400	1622016	4816896
4	3	14745600	47185920	41746432	117440512
4	4	415236096	1233125376	1006632960	2717908992
4	5	10687086592	30064771072	23102226432	60397977600

m	n	$\tau(P_n^{(m)})$	$\tau(A_n^{(m)})$	$\tau(B_n^{(m)})$	$\tau(C_n^{(m)})$
5	0	722	6724	12482	58564
5	1	123904	739328	984064	3964928
5	2	12781568	59969536	65619968	237899776
5	3	1007681536	4060086272	3901751296	13088325632
5	4	67194847232	243609370624	213408284672	674448277504
5	5	3995393327104	243609370624	10953240346624	33019708571648

5. Concluding remarks

The number of spanning trees $\tau(G)$ in graphs (networks) is an important invariant. The computation of this number is not only pleasant from a mathematical (computational) standpoint, but also, it is an important measure of reliability of a network and electrical circuits layout. Some computationally laborious problems such as the travelling salesman problem can be resolved approximately by using spanning trees. Due to the high reliance of the network design and reliability of the network we gave the above important Theorems and their proofs.

Competing Interests

The authors declare that there are no competing interests regarding the publication of this paper.

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