

Article

Interference-Aware Adaptive Beam Alignment for Hyper-Dense IEEE 802.11ax Internet-of-Things Networks

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Abstract: The increased deployment of IoT devices in specific areas results in an interference among them and the quality of communications can be severely degraded. To deal with this interference issue, the IEEE 802.11ax standard has been established in hyper-densely wireless networking systems. The 802.11ax adopts a new candidate technology that is called multiple network allocation vector in order to mitigate the interference problem. In this paper, we point out potential problem in multiple network allocation vector which can cause delays to the communication among IoT devices in hyper-dense wireless networks. Furthermore, this paper introduces an adaptive beam alignment algorithm for interference issue resolution. In addition, we analyze potential delays of communications among IoT devices under interference conditions. Lastly, we simulate our proposed algorithm in densely deployed environment and show that the interference issue can be mitigated and the IEEE 802.11ax-based IoT devices can utilize the air interface more fairly compared to conventional IEEE 802.11 distributed coordination function.

Keywords: IEEE 802.11ax, multiple NAVs, Internet-of-Things (IoT), Beamforming

1. Introduction

According to a recent Cisco technical report [1], there will be 3.5 networked devices per capita by 2021, from 2.3 networked devices per capita in 2016. The trend of ever increasing mobile devices led to a tremendous deployment of access points (APs) for increasing the user satisfaction. As a result, congestion and contention problems among devices in overlapping basic service set (OBSS), as shown in Fig. 1, become important, which severely deteriorates the user experience. For efficiently handling this issue, the task group of IEEE 802.11ax (TGax) is conducting IEEE 802.11ax standardization, in what is called the high-efficiency wireless local area networks (HE-WLAN) fundamentally based on IEEE 802.11ac to address the congestion problems [2–5].

The 802.11ax HE-WLAN adopts a new scheme, called the multiple network allocation vectors (NAV)s, in order to mitigate the interference problem. Each IEEE 802.11ax-based Internet of Things (IoT) stations hold multiple NAVs to avoid wireless congestion. If an STA is located in a certain area which is OBSS created by three APs, the STA hold three different NAVs for the corresponding APs. Specifically, the NAV that is related to the associated AP is called *intra-BSS NAV* and the other NAVs are called *inter-BSS NAVs*. The intra-BSS NAV and inter-BSS NAVs are utilized for avoiding interference among STAs, which are located within associated AP and other alien APs, respectively. If other STAs which are associated with the AP, then the intra-BSS NAV is set and the STA should idle until the reception of *contention free end frame* from the AP. On the other hand, if STAs which are associated with other alien APs are communicating among themselves, the STA should be idle due to the inter-BSS NAV. That is, multiple NAVs consist of intra-BSS NAV and inter-BSS NAVs. In addition, the multiple NAVs behave strategically to avoid interference among STAs in OBSS networks and can be utilized

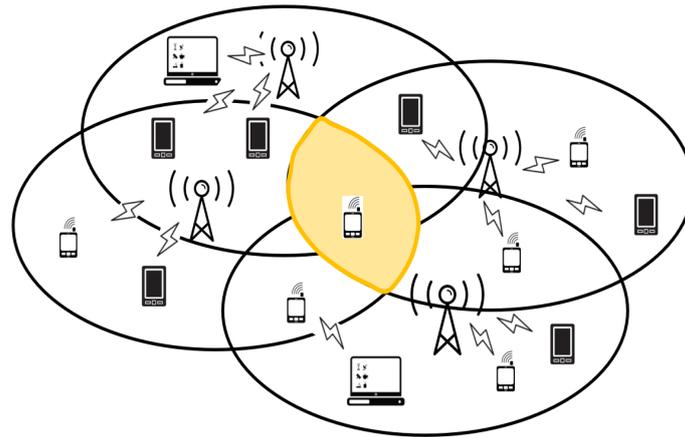


Figure 1. Scenario of hyper-dense APs: multiple inter-BSS NAVs may deteriorate medium utility chance of an STA in OBSS.

32 properly to transmit data for each STA. To sum up, if at least one of inter-BSS NAVs or intra-BSS NAV is set in a
 33 specific subcarrier, the STA cannot transmit data because it assumes the channel is busy. Envisioning a future
 34 WLAN which configures hyper-dense deployment of countless devices in a certain small area, STA may set
 35 multiple inter-BSS NAVs and consequently it could be excessively delayed for even single transmission.

36 For avoiding excessive delay, we propose a novel algorithm to address the delay problem in an STA with
 37 data in its data queue that needs to be transmitted and is blocked due to several inter-BSS NAVs or intra-BSS
 38 NAV settings. This means that the multiple NAV can be helpful for avoiding interference problem but it can
 39 be harmful for delay issues. Therefore, this paper targets a new objective, i.e., a joint optimization of delay
 40 reduction and interference minimization. If the time interval of the idle state exceeds a threshold, the delayed
 41 STA makes a decision which AP to be associated with the STA among its own AP and other alien APs and
 42 swiftly finds a direction towards the candidate AP within a low time complexity. Finally, the STA configures a
 43 beamformed-beam to transmit data to the designated AP to calculated direction. Our algorithm in IEEE 802.11ax
 44 enables an STA to transmit data in densely deployed environment even though it has a set of inter-BSS NAVs or
 45 an intra-BSS NAV. We demonstrate an analytical model of the proposed algorithm and simulate its performance
 46 compared to IEEE 802.11 distributed coordination function (DCF) in terms of fairness and expected transmission
 47 time loss.

48 The main contribution of this paper is the design of an algorithm which mitigates the delay issue of STAs in
 49 the heavily dense deployment scenario of 802.11ax based IoT wireless networks. For this purpose, we firstly
 50 propose low-complexity and swift beam direction selection algorithm that is based on opportunistic beamforming
 51 MAC protocols. The rest of this paper is constituted as follows. Section 2 introduces related work and the
 52 enhanced features of IEEE 802.11ax. Section 3 proposes an opportunistic beamforming algorithm and describes
 53 the analytic model of the proposed algorithm. Section 4 performs intensive simulations and presents discussions
 54 based on the simulation results. Finally, section 5 concludes this paper.

55 2. Related Work and IEEE 802.11ax Features

56 The dense deployment of IoT devices in the aforementioned scenario can be represented as in Fig. 2.
 57 IEEE 802.11ax is designed for high performance distributed networking in dense deployment scenarios, such
 58 as carriages, residential apartments, and auditoriums. Based on the nature of IEEE 802.11ax, it is one of the
 59 most suitable solutions for hyper-dense IoT wireless communications. By adopting the aforementioned multiple
 60 NAV strategy, data transmission in dense deployment scenarios where tens of APs and hundreds of STAs co-exist
 61 simultaneously in a small area can be delayed due to wireless medium resource competition. That is, we can
 62 easily envision that each device should wait until all NAVs are reset when there exists multiple IoT devices which

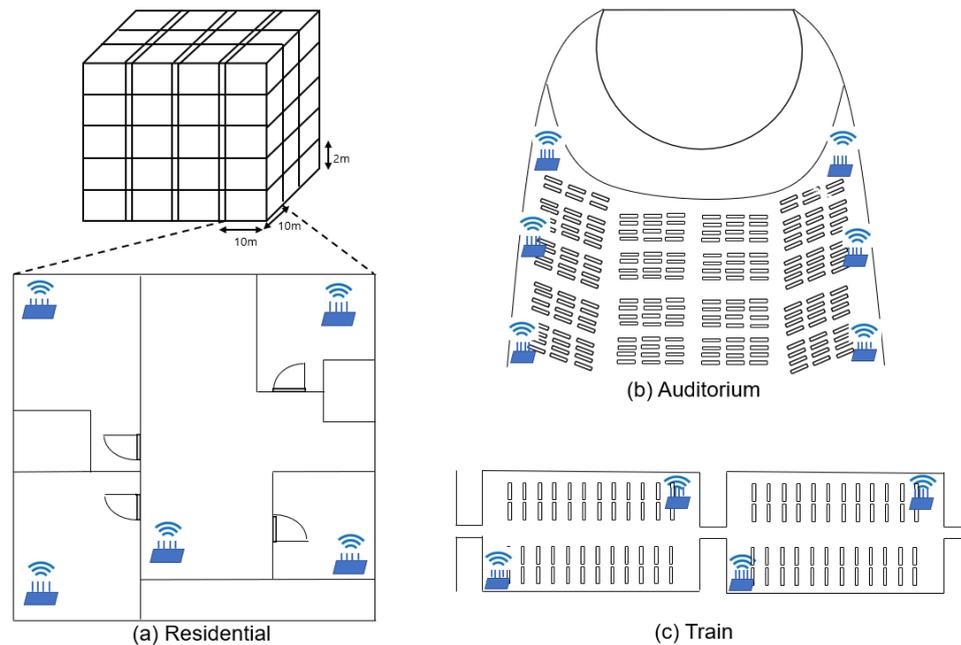


Figure 2. Three densely deployed APs scenarios of wireless devices. Fig. (a) represents a plausible residential scenario which is composed of densely deployed multiple APs in an apartment. Fig. (b) is a scenario for multiple APs in an auditorium. Lastly, Fig. (c) is an example of densely installed APs scenario in a carriage.

63 want to aggressively utilize the air interface. In this section, previous research results and primary features of the
 64 802.11ax standard are introduced.

65 The research results in [6] demonstrated a downlink multi-user MIMO (DL MU-MIMO) along with new
 66 features of IEEE 802.11ax, i.e., uplink MU-MIMO (UL MU-MIMO) transmission in AP-initiated scenarios.
 67 In [6], the differences between legacy IEEE 802.11 wireless local area networks (WLAN) and IEEE 802.11ax
 68 were introduced. In addition, the analytical saturation throughput model of transmission scenarios for both
 69 single-user (SU) and multi-user (MU) were proposed. Furthermore, the IEEE 802.11ax and IEEE 802.11ac were
 70 compared in [7]. The insights of DL/UL orthogonal frequency division multiple access (OFDMA), dynamic clear
 71 channel assessment (CCA), and UL MU-MIMO concepts were also introduced in [7]. Surveys on OFDMA-based
 72 medium access protocols and OFDMA-based concurrent MU medium access control algorithms for IEEE
 73 802.11ax were presented in [8,9].

74 The research results in [10,11] were focused on dynamically tuning of CCA which enhances spatial
 75 reusability for addressing OBSS congestion problems. In terms of implementation, the link-system level
 76 simulator for IEEE 802.11ax with network simulator 3 (NS3) was explored in [12]. In [12], plausible transmission
 77 parameters, dense deployment scenarios, and the simulated throughput of each BSS with various numbers of
 78 STAs were studied. The authors of [13] studied quality-of-service (QoS) support in legacy IEEE 802.11 wireless
 79 networks and summarized the IEEE 802.11ax standardization processes. In addition, they presented an overview
 80 of current and expected features of IEEE 802.11ax in terms of the medium access control. Moreover, emerging
 81 long-term evolution licensed-assisted access which satisfied user needs with low latency and high bandwidth
 82 were jointly considered with IEEE 802.11ax for collaboration between cellular networks and IEEE 802.11
 83 WLAN-based wireless networks. Besides, the method in [14] suggested a next generation medium access control
 84 mechanisms in unlicensed bands (5G-U) for vehicular radio access with licensed-assisted access (LAA) and
 85 IEEE 802.11ax wireless networks. A new medium access control algorithm under the consideration of resource
 86 uncertainty and physical sidelink shared channel (PSSCH) were proposed to deploy LTE vehicle-to-everything
 87 (V2X), channel pre-occupation, and resource binding strategies. The algorithm in [15] proposed co-existence
 88 mechanisms of IEEE 802.11ax with cellular, LTE with unlicensed (LTE-U), and LAA. The authors of [15]
 89 considered single unlicensed frequency band transmissions that the locations of Wi-Fi STAs, Wi-Fi APs, and LTE

evolved node B (eNB) are modeled as three independent homogeneous Poisson point processes. They derived analytical expressions for a set of metrics including signal-to-interference-plus-noise ratio coverage probability, density of successful transmissions, and Shannon throughput probability for both of the UL and DL of IEEE 802.11ax wireless networks and cellular networks.

There exist several improvements in IEEE 802.11ax for various purposes to enhance data rate in order to mitigate interference and improve the efficiency of frequency utilization. Among them, target wake time (TWT) is one of the improvements in IEEE 802.11ax. The TWT has advantage in terms of collision probability reduction in high-dense distributed IEEE 802.11-based WLAN network scenarios with pre-knowledge awake time of STAs with associated AP. The mechanism of TWT is extremely simple yet effective by calculating and analyzing when STA wakes up to transmit data over time. Furthermore, the TWT enables STAs to operate energy efficiently with the doze mode. For example, authors in [16] proposed a scheduling scheme in IEEE 802.11ax with TWT to efficiently handle multi-user transmissions and multi-AP cooperation. The authors in [17] explored several new IEEE 802.11ax UL scheduling mechanisms and compared them among the maximum throughput of MU. They conducted evaluations on MIMO and OFDMA transmissions in IEEE 802.11ax versus the carrier-sense multiple access with collision avoidance (CSMA/CA) MAC of IEEE 802.11ac with the SU and MU modes for various number of STAs scenarios in both reliable and unreliable channel environments. The authors of [18] suggested advanced IEEE 802.11ax TCP-aware scheduling strategies to optimize the performance using MU UL and MU DL transmissions. They were based on transmission opportunities (TXOP) and can control achieved goodput versus delays. The authors showed that minimal goodput degradation strategy can avoid tremendous delay. In [19], the authors proposed a semi-matching based load balancing scheme for hyper-dense IEEE 802.11 wireless networks. The scheme is operated in a centralized controller which makes decision whether the load is unevenly distributed among APs and controls the overall network throughput to be maximized. By taking advantage of IEEE 802.11ax new feature considerations, simultaneous transmit and receive (STR) which is also called full-duplex, is likely to be applied in legacy IEEE 802.11-based wireless standards. One of the key challenges is to integrate the STR mode with minimal modifications into legacy standards. In [20], the authors proposed a simple yet practical approach to enable the STR mode in IEEE 802.11-based wireless networks with co-existing full-duplex and half-duplex STAs. TGax for IEEE 802.11ax standardization plans to adopt some new features such as dynamic CCA tuning, multiple network allocation vectors, and TWT. The TWT is an energy efficient strategy in terms of the operation of STA considering that the associated AP of the STA has full knowledge of awake schedule of the STA and thus the AP properly transmits data by calculating the awake time of STA which is not in doze mode. In addition, a multi-NAVs strategy for efficient frequency utilization is considered for mitigating wireless interference in hyper-dense situations. The STA set NAV for an AP where the AP initiates transmission. If at least one of the inter-BSS NAVs or intra-BSS NAV of STA is set, the STA should be idle for avoiding congestion until all NAVs are reset. In the following subsections, IEEE 802.11ax specific wireless frame format and the error correction and modulation of IEEE 802.11ax are proposed.

2.1. PPDU frame structure

IEEE 802.11ax TGax standard changed the physical protocol data unit (PPDU) frame format compared to IEEE 802.11ac for improving frequency, channel, and medium utilizations. As shown in Fig. 3, the IEEE 802.11ax PPDU frame formats keep legacy preambles for backward compatibility and add *RL-SIG* field behind the legacy preambles for automatic detection of IEEE 802.11ax frames. In addition, IEEE 802.11ax IoT devices detect channel state information (CSI) by referring to the *L-LTF* field in Fig. 3. If the result of modulo three operations of the length field of *L-SIG* is 1, the frame is for SU. If the result is 2, then the frame is for MU transmission. A null data packet (NDP) frame for CSI exchange is in Fig. 3(c) and is transmitted by AP. When an STA receives NDP announce (NDPA) frame followed by NDP, the STA transmits CSI report frame to AP and the AP transmits trigger frame as a response to CSI report to initiate UL transmission.

2.2. Error correction and modulation

The legacy IEEE 802.11 standard adopted block check character (BCC) for forward error correction (FEC), and not the low density parity check (LDPC) due to its high computational cost. However, IEEE 802.11ax

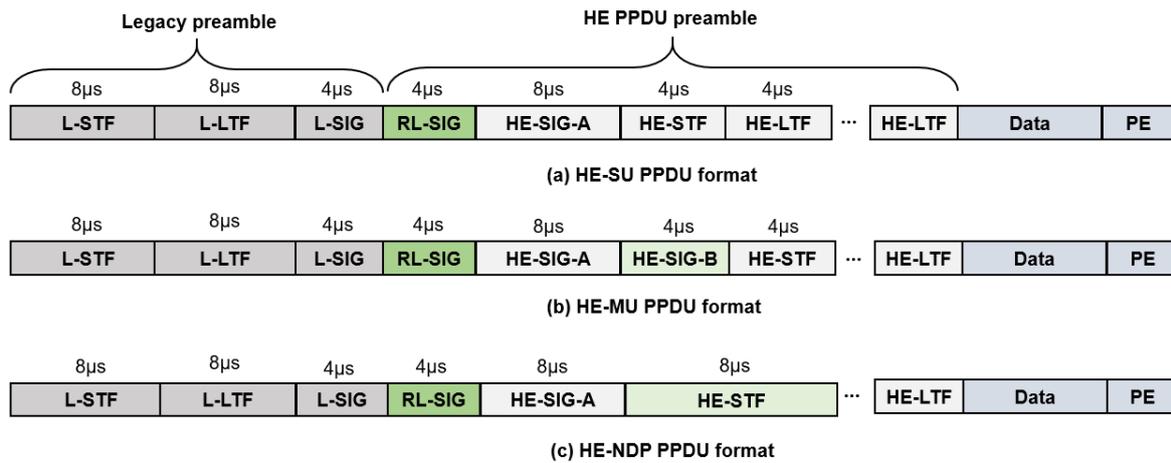


Figure 3. IEEE 802.11ax PPDU frame format.

Table 1. A comparison between the IEEE 802.11ac and the IEEE 802.11ax

	IEEE 802.11ac	IEEE 802.11ax
Band [GHz]	5	2.4 and 5
Channel bandwidth [MHz]	20, 40, 80, 80+80, 160	20, 40, 80, 80+80, 160
Modulation	BPSK, QPSK, 16QAM, 64QAM, 256QAM	1024QAM is newly added
FFT size	64, 128, 256, 512	256, 512, 1024, 2048
Subcarrier spacing [KHz]	312.5	78.12
Symbol duration [us]	3.2	12.8
CP [us]	0.4 and 0.8	0.8, 1.6, and 3.2
FEC	BCC, LDPC (optional)	LDPC
Spatial stream (SS)	Up to 8 SS for each AP Up to 4 SS for each STA	Up to 8 SS for each AP Up to 4 SS for each STA
MU-MIMO	DL MU-MIMO	UL/DL MU-MIMO

138 selects LDPC for error correction as mandatory because it performs better in terms of capacity compared to
 139 IEEE 802.11ac. In addition, IEEE 802.11ax selectively utilizes dense constellation, such as 1024-QAM, so that
 140 each symbol contains more information and achieves an improved throughput performance. Furthermore, IEEE
 141 802.11ax exploits four-fold larger fast Fourier transform (FFT) size compared to the IEEE 802.11ac so that the
 142 spectral efficiency can be improved for appropriate dense scenarios, as shown in Table 1. Although its subcarrier
 143 is narrower than before, the inter symbol interference (ISI) problem is relieved by setting longer guard interval
 144 (GI). By utilizing these improvements in terms of medium access control, data transmission under severe delay
 145 spread in outdoor environments can be successful.

146 3. Proposed Opportunistic Medium Access

147 In this section, we discuss an excessively long delay problem that may occur in IEEE 802.11ax and propose
 148 a swift and low-complexity beam direction selection algorithm based on opportunistic beamforming medium
 149 access control. First, the simplified mechanism of the DL and UL transmissions of IEEE 802.11ax is as illustrated
 150 in Fig. 4. In case of DL MU transmission, the data packets from an AP are transmitted simultaneously to multiple
 151 STAs and each STA replies the corresponding *block ACK* to the associated AP. On the other hand, STAs, which
 152 have packets to transmit simultaneously, send their packets as they receive the *trigger frame*. In the following
 153 subsections, the mechanism and analytical model of the proposed medium access control algorithms and related
 154 assumptions are presented.

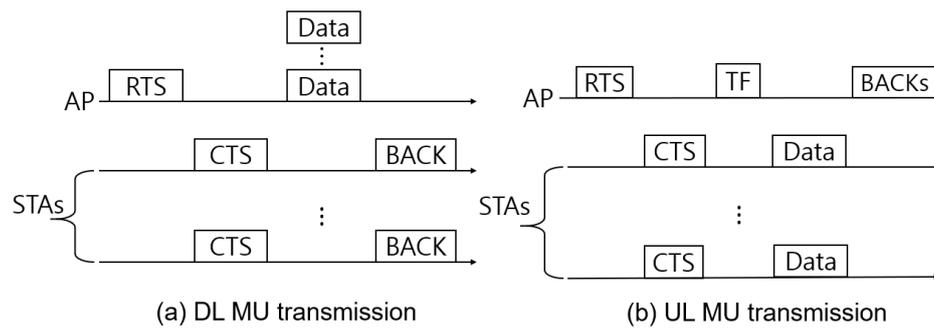


Figure 4. Simplified DL and UL transmissions of IEEE 802.11ax.

Table 2. Comparison with other OFDMA protocols

	Proposed	[21]	[22]	[23]	[24]
MU transmission	✓	✓	✓	✓	✓
MU access	✓	✓	✓	✓	✓
MU diversity	✓		✓		
Simple signal exchange	✓	✓	✓	✓	
CSI measurement	✓				

In terms of the various evaluation criteria used for our algorithms and previous research, and used for comparison between the previous research [21–24] and our algorithm is proposed as Table 2. The MU transmission and MU access represent feature of simultaneous channel access and concurrent data transmission of STAs by taking advantage of the frequency orthogonality of OFDMA. Hence, channel access efficiency of STAs can be enhanced by utilizing this feature. In addition, MU diversity stands for the characteristic that each STAs can be allocated only high channel gain resource blocks (RBs) and can improve their throughput performance. Simple signal exchange shows that simple control message exchange procedure in communication system. In case of 802.11ax and our proposed protocol, *trigger* frame, which is one of the control message among data transmission procedure, is related to this feature. Finally, channel state information (CSI) measurement item refers to the measurement process of channel state between AP and STAs. The AP and STAs exchange the result of this procedure and utilize the information of channel to appropriately tune their transmission parameters.

3.1. System model and assumptions

In order to understand our proposed algorithm, consider a simple IEEE 802.11ax network which consists of multiple OBSSes with five APs and twenty STAs in a narrow area as presented in Fig. 5 (only two APs and five STAs are represented for simplification). In particular, the user experience of STA_2 , which is our main focus, is deteriorated during the red-lined duration in Fig. 5 due to multiple inter-BSS NAVs setting. It is straightforward to envision that the red-lined duration will get longer while the more affected APs are increased. Thus, this phenomenon can be much more severe and the STA hardly transmits data due to increased delays in hyper-dense wireless networking scenarios. If affecting alien APs are relaying to transmit for each transmission cycle, it is obvious that the STA gets the worst chance to transmit data. To address this issue, we propose a low-complexity and swift beam direction selection algorithm based on *opportunistic beamforming* data transmission for letting delayed STA transmit data. Therefore, the proposed algorithm consists of two modules: i) appropriate beamformee designation and direction configuration procedure and ii) data transmission through the direction with beamforming.

In order to analyze our proposed algorithm in hyper-dense IEEE 802.11ax wireless networks, the following assumptions are made.

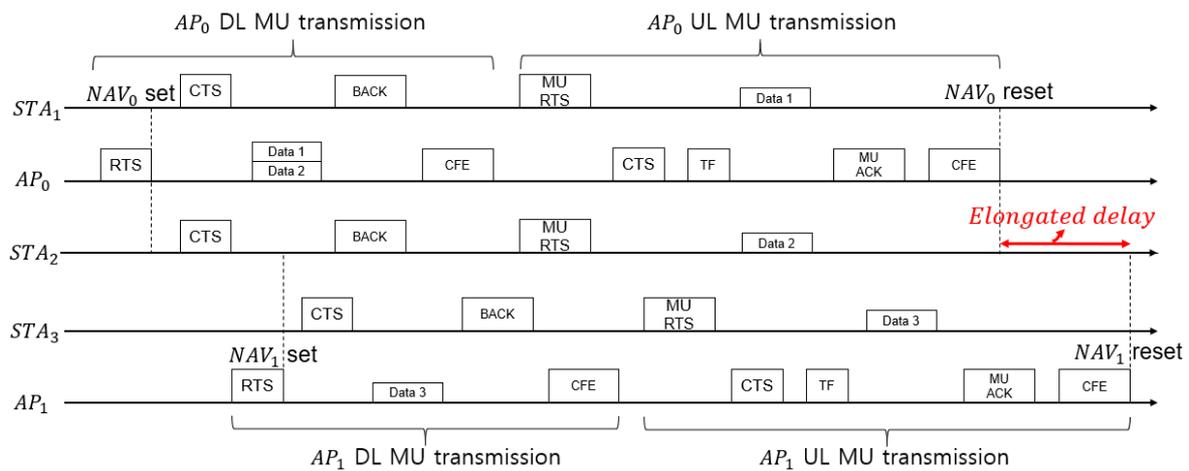


Figure 5. Elongated delay scenario of STA_2 due to inter-BSS NAVs NAV_1 setting.

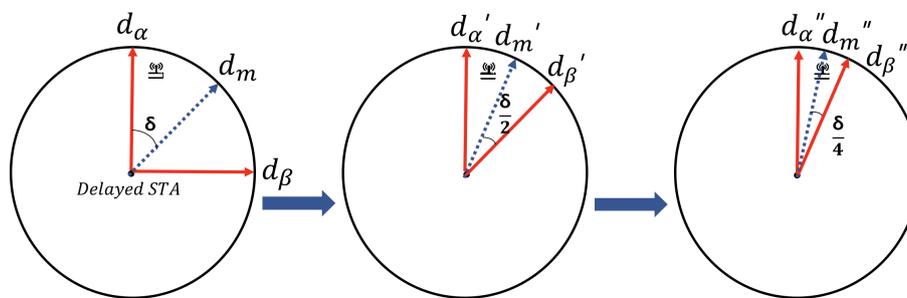


Figure 6. Detailed mechanism of $Search_{HR}^e((d_\alpha, d_\beta), BSR)$ in Algorithm 1.

- 181 • Each STA has exactly the same performances with OFDMA and MU-MIMO for mathematical performance
 182 analysis.
 183 • In addition, antennas which are installed in each IEEE 802.11ax IoT device are full-duplex vouching
 184 simultaneous U/DL transmission.
 185 • Detailed parameter setting of antennas for beamforming including azimuth, half power beam width
 186 (HPBW), and antenna gains are not associated with the proposed algorithm.
 187 • U/DL MU transmissions are considered in this paper.
 188 • In our proposed algorithm, NAV is set to 0 if the corresponding AP is idle and vice versa.

189 3.2. Beam direction selection and beamforming algorithm

190 This section introduces a fast and lightweight beam direction selection based beamforming algorithms to
 191 address the aforementioned issue in IEEE 802.11ax wireless networks. The main mechanisms of our algorithms
 192 can be summarized as follows. i) If an STA attempts to transmit data, it considers its inter-BSS NAVs are set. If
 193 all of them are reset then it immediately transmits data to its associated AP (line 14 of Algorithm 2). However,
 194 if at least one of them is set, the STA should not act selfish to transmit data to avoid inter-BSS congestion
 195 deterioration. In addition, the delayed STA due to inter-BSS NAV setting investigates a new AP to associate after
 196 a specific duration θ_{ed} , i.e., the STA has *wait threshold* variable ω to avoid excessive delays and only finds the
 197 appropriate direction of beamformee after the ω exceeds waiting threshold θ_{ed} (line 3 of Algorithm 2). ii) After
 198 the STA makes decision to find the beamformee, it chooses the associated AP or alien APs by sorting them based
 199 on a descending order of their signal strengths (line 2 of Algorithm 1) and finds the most appropriate direction
 200 (line 3 to 13 of Algorithm 1). iii) STA beamfoms toward its beamformee in the direction of dir , which is the
 201 result of Algorithm 1 (line 7 to 11 of Algorithm 2). Notice that the beam alignment procedure is recursively

Algorithm 1 Joint searching for direction and CSI**Input:** $LR_{dir}^a, HR_{dir}^a, AP_{CS}^a, CSI_{AP_{CS}^e}, HR_{dir}^e, d_\alpha, d_\beta$ **Output:** Appropriate *direction* for beamforming

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1: Initialize Beam search rate (BSR)
2: Listen channel
3: if  $AP^a$  is dominant then
4:    $LR_{dir}^a = Search_{LR}^a(AP_{CS}^a, BSR)$ 
5:    $HR_{dir}^a = Search_{HR}^a(LR_{dir}^a, BSR)$ 
6:   return  $HR_{dir}^a$ 
7: else
8:   Sort candidate  $AP^e$ s based on signal strength
9:   Select dominant alien  $AP^e$ 
10:  Set  $\langle d_\alpha, d_\beta \rangle$ 
11:   $HR_{dir}^e = Search_{HR}^e(\langle d_\alpha, d_\beta \rangle, BSR)$ 
12:  Set  $CSI_{AP_{CS}^e}$ 
13:  return  $\langle HR_{dir}^e, CSI_{AP_{CS}^e} \rangle$ 
14: end if

```

202 called as in Fig. 6 for finding the suitable beam direction angle (azimuth and elevation angle of beamforming)
 203 between the delayed STA and target AP which is within the range of the HPBW of STA. The STA actively
 204 searches with a beam search rate (BSR) between the originally associated AP or other the alien APs based on the
 205 received signal strength indicators (RSSI) of each AP given that the wireless data traffic loads on each AP are
 206 closely similar so that the RSSI is a proper criterion for selecting the AP.

Algorithm 2 Proposed beamforming algorithm**Input:** HR_{dir}^a or $\langle HR_{dir}^e, CSI_{AP_{CS}^e} \rangle, \omega, \theta_{ed}, dir$

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1: Check  $\forall inter - BSSNAV_A[i]$ 
2: if  $\exists inter-NAV_A[i]$  is set then
3:   if  $\omega < \theta_{ed}$  then
4:     Increase  $\omega$ 
5:   else
6:      $dir = \text{Result of Algorithm 1}$ 
7:     if  $dir$  is  $HR_{dir}^a$  then
8:       Beamform to  $HR_{dir}^a$ 
9:     else
10:      Beamform with  $\langle HR_{dir}^e, CSI_{AP_{CS}^e} \rangle$ 
11:    end if
12:  end if
13: else
14:  Immediately transmit DATA through  $AP^a$ 
15: end if

```

207 Based on our assumptions, if the STA chooses to transmit data via its associated AP^a , it utilizes the
 208 corresponding information which already has and exploits it for searching direction set of AP in low resolution
 209 (LR_{dir}^a) within (AP_{CS}^a) direction set under BSR (line 4 of Algorithm 1). After that, the STA searches a beamformee
 210 with a higher resolution (HR_{dir}^a) by exploiting the result of low resolution direction (line 5 of Algorithm 1). The
 211 main difference between low resolution and high resolution is the searching direction size in a circular sector
 212 (CS) scale, which is the range of direction searching procedure where low resolution is wider than the high

Table 3. Parameters used in analysis

Parameter	Description
N_A	Number of APs
N_S	Number of STAs
U_{mu}	Number of MU STAs
$\tau(x)$	Return time when $x = 0$
NAV_i	i -th inter-BSS NAV
NAV_m	intra-BSS NAV
L_c	Length of control frame
L_d	Length of data frame
B_r	Beamforming search rate
x_i^t	Throughput of i -th STA in time t
n	Required count to find appropriate beamformees
V_s	Number of SU-MIMO spatial streams per STA
V_m	Number of MU-MIMO spatial streams per STA
$R(V_s, B_{ru})$	Data rate
B_{ru}	Bandwidth of RU
L_D	Packet size
P_a	Number of aggregated packets in A-MPDU

213 resolution's one. On the other hand, if the STA chooses to transmit among one of the alien AP^e s, the STA sorts
 214 them considering the signal strength with a descending order and selects one of them (line 8 to 9 of Algorithm
 215 1). Later, the STA sets CSI between d_α and d_β , which is the initial range of searching direction area (line 10
 216 of Algorithm 1), and finds the direction of targeted alien AP (HR_{dir}^e) by calling $Search_{HR}^e$ function recursively
 217 while dynamically tuning the d_α and d_β per every function call (line 11 of Algorithm 1). The detailed operation
 218 of $Search_{HR}^e(\langle d_\alpha, d_\beta \rangle, BSR)$ is represented in Fig. 6. After these processes, the STA finally sets up CSI between
 219 the selected AP^e and gets the joint beamforming information (line 12 to 13 of Algorithm 1). Finally, it can be
 220 observed that the STA addresses delay issue and the STA can transmit data to either AP^a or AP^e in the direction
 221 of the calculated result in order to improve the performance in terms of the interference reduction.

222 To sum up, the STA with our proposed beamforming algorithm in hyper-dense environments does not act
 223 selfishly but little tenacious to cooperate with other STAs in OBSS to avoid wireless congestion. In addition,
 224 as illustrated in Fig. 6, if the STA makes a decision to associate with AP^e s, then finding a dominant AP^e
 225 and searching appropriate beam direction are required with the complexity of $O(n \log(n))$, according to the fact
 226 that finding dominant AP^e among AP^e s requires $O(n)$ and searching appropriate beam direction requires
 227 $O(\log(n))$.

228 3.3. Analytical model

229 In this section, the total elapsed time for D/UL MU transmissions are formulated. This measurement is used
 230 for calculating the lost time of the denied transmission of STA and can be utilized for computing the expected
 231 lost throughput. The analysis regarding the expected lost throughput can be conducted with the expected lost
 232 time as red-lined in Fig. 5. Furthermore, fairness on wireless medium access is also used to compare DCF with
 233 our proposed algorithm using Jain's index [25]. The used parameters and descriptions for the mathematical
 234 analysis are summarized in Table 3.

235 3.3.1. Total elapsed time of D/UL MU transmission

236 STAs which communicate with AP for D/UL follow specific transmission procedure as defined in the IEEE
 237 802.11ax TGax standardization.

238 In case of DL MU transmission, AP initiates IEEE 802.11ax-based transmission procedure by sending
 239 request-to-send (RTS) frame to its associated STAs. After short inter frame space (SIFS), STAs responds with a
 240 clear-to-send (CTS) to the AP announcing that they are ready for receiving DL MU transmission data. Then,

241 the AP concurrently transmits DL MU data to STAs. After receiving all data from the AP, STAs responds
 242 with block acknowledgement (BACK) to the AP in order to inform successful reception. Finally, the AP sends
 243 contention-free-end (CFE) to reset its NAV of STAs.

244 In case of UL MU transmission, STAs which want to upload data to AP send MU RTS frame to the AP.
 245 The AP responds with CTS and sends NDP/NDPA frames before transmitting *trigger* frame which initiates
 246 STAs' UL MU transmission. After STAs trigger frame reception, the STAs concurrently transmit their data.
 247 After the termination of STAs' data transmission procedure, the AP sends MU ACK back to STAs for informing
 248 successful UL data transmission. Finally, it transmits CFE to STAs, same as in the DL procedure and STAs reset
 249 corresponding NAV.

250 Based on this description, the total elapsed time for a single downlink and uplink MU transmission can be
 251 formulated as (1) and (2), respectively.

$$\mathcal{T}_{\text{mu}}^d(U_{\text{mu}}, V_s, B_{\text{ru}}) = \mathcal{T}_{\text{mu-RTS}}(U_{\text{mu}}) + \mathcal{T}_{\text{SIFS}} + \mathcal{T}_{\text{CTS}} + \mathcal{T}_{\text{SIFS}} + \mathcal{T}_{\text{mu,d}}^{\text{DATA}}(U_{\text{mu}}, V_s, B_{\text{ru}}) + \mathcal{T}_{\text{SIFS}} + \mathcal{T}_{\text{BA}} + \text{AIFS} + \mathcal{T}_e \quad (1)$$

$$\mathcal{T}_{\text{mu}}^u(U_{\text{mu}}, V_m, V_s, B_{\text{ru}}) = \mathcal{T}_{\text{mu-RTS}}(U_{\text{mu}}) + \mathcal{T}_{\text{SIFS}} + \mathcal{T}_{\text{CTS}} + \mathcal{T}_{\text{SIFS}} + \mathcal{T}_{\text{trigger}}(U_{\text{mu}}) + \mathcal{T}_{\text{SIFS}} + \mathcal{T}_{\text{mu,u}}^{\text{DATA}}(V_s, B_{\text{ru}}) + \mathcal{T}_{\text{SIFS}} + \mathcal{T}_{\text{mu-ACK}}(V_m) + \mathcal{T}_e \quad (2)$$

252 The above two equations, i.e., (1) and (2), include the elapsed time of backoff, data transmission time, and
 253 control frames exchange times in D/UL transmission. $\mathcal{T}_{\text{mu}}^d(U_{\text{mu}}, V_s, B_{\text{ru}})$ and $\mathcal{T}_{\text{mu}}^u(U_{\text{mu}}, V_m, V_s, B_{\text{ru}})$ refer to the
 254 total duration of DL and UL MU transmission respectively. They include commonly $\mathcal{T}_{\text{mu-RTS}}(U_{\text{mu}}) + \mathcal{T}_{\text{SIFS}} +$
 255 $\mathcal{T}_{\text{CTS}} + \mathcal{T}_{\text{SIFS}}$ for initial control message exchange procedure. $\mathcal{T}_{\text{mu,d}}^{\text{DATA}}(U_{\text{mu}}, V_s, B_{\text{ru}})$ in DL MU transmission
 256 represent required time for data transmission for U_{mu} with $R(V_s, B_{\text{ru}})$ per each STA. After the transmission
 257 of data through DL after SIFS, STAs transmit BACK, which consumes \mathcal{T}_{BA} and idle for $\text{AIFS} + \mathcal{T}_e$ which
 258 is the duration of arbitrary inter symbol space (AIFS) and CFE. $\mathcal{T}_{\text{trigger}}(U_{\text{mu}})$ in UL MU transmission and
 259 $\mathcal{T}_{\text{mu,u}}^{\text{DATA}}(V_s, B_{\text{ru}})$ stand for the duration of trigger frame to U_{mu} UL MU STAs and the required time for MU UL
 260 data transmission, respectively. Finally, after SIFS, $\mathcal{T}_{\text{mu-ACK}}(V_m) + \mathcal{T}_e$ is required for terminating entire UL
 261 MU transmission procedures. These are utilized to calculate the cost values of expected lost time l_t (refer to
 262 Section 3.3.2) and the expected lost throughput $E_{L_{Th}}$ (refer to Section 3.3.3) in the following subsections.

263 3.3.2. Expected lost time l_t

264 The τ function that is defined as

$$l_t = \tau \left(\sum_{i=0}^{N_A} \text{NAV}_i \right) - \tau(\text{NAV}_m) \quad (3)$$

265 is used for calculating the total time difference depending on NAV values when intra-BSS NAV and all other
 266 inter-BSS NAVs are reset to zero. The (1) and (2) are used to calculate the time of each NAV duration. Thus,
 267 the (3) represents the delayed period of an STA affected by the inter-BSS NAVs set which are generated by
 268 neighboring APs in dense deployment environments.

269 3.3.3. Expected throughput loss, E_{Th}

Based on (3), the expected throughput loss of STA during delayed period can be calculated under the
 consideration of missed control frames and data frames. Ahead of the detailed analysis, note that all data rates
 for control frames are statically set to 6 Mbps which is the data rate of a single spatial stream with 20MHz
 channel in IEEE 802.11ax (i.e., $R(1, 20\text{MHz})$). In addition, the n is proportional to N_A where $n = E_S N_A$ and

E_S represents the average neighboring N_S . The expected loss of control and data frames can be analytically formulated as in the following; (4) and (5).

$$E_{L_c} = \underbrace{l_t \cdot \left[\frac{l_c}{R(1, 20\text{MHz})} \right]}_{T_c^*} \cdot \underbrace{l_t \cdot \left(1 - \frac{B_r}{l_t} \right)^{\frac{n-1}{n}} \cdot \left(\frac{B_r}{l_t} \right)^{\frac{1}{n}}}_{P_c^*} \quad (4)$$

270 where E_{L_c} stands for the expected throughput loss of control frame. In (4), T_c^* denotes the throughput loss during
271 l_t with $\left[\frac{l_c}{R(1, 20\text{MHz})} \right]$ size of control frame per unit time and P_c^* stands for the successful probability to find the
272 appropriate beamformee with beam search rate B_r in time period l_t in n counts.

Similarly, the expected throughput loss of data frame, i.e., E_{L_d} , can be formulated as follows:

$$E_{L_d} = \underbrace{l_t \cdot \left[\frac{l_d}{R(V_s, B_{ru})} \right]}_{T_d^*} \cdot \underbrace{l_t \cdot \left(1 - \frac{B_r}{l_t} \right)^{\frac{n-1}{n}} \cdot \left(\frac{B_r}{l_t} \right)^{\frac{1}{n}}}_{P_d^*} \quad (5)$$

273 where T_d^* and P_d^* denote the throughput loss duration l_t with data rate $R(V_s, B_{ru})$ and successful probability for
274 finding the appropriate beamformee with beam search rate B_r for the duration l_t in n counts, respectively.

275 3.3.4. Jain's index \mathcal{J} for fairness

Finally, the total expected loss of throughput can be denoted by:

$$E_{L_{Th}} = E_{L_c} + E_{L_d}. \quad (6)$$

In addition, the medium access fairness can be evaluated by the following Jain's index, i.e.,

$$\mathcal{J}(x_1^t, x_2^t, \dots, x_i^t) \triangleq \frac{\left(\sum_{i=1}^{N_S} x_i^t \right)^2}{N_S \cdot \sum_{i=1}^{N_S} (x_i^t)^2} \quad (7)$$

276 which each of x_i^t represents throughput of i -th STA in time step t .

277 4. Performance Evaluation

278 In this section, we discuss the simulation setting and results regarding delay scenario which is considered in
279 the previous section. We simulate the denied transmission period based on the analytic model which is presented
280 in the previous section for both U/DL MU transmission. Moreover, we evaluate our proposed medium access
281 control protocol in terms of fairness using Jain's index. We compare the fairness values of the ordinary DCF and
282 the proposed algorithm under different options and various numbers of STAs.

283 4.1. Simulation setting and overview

284 Our simulation is designed and implemented based on the assumptions and formulations in Section 3 and
285 Table 3, in order to precisely operate with obeying these settings. The parameters and corresponding values used
286 in our simulation are summarized in Table 4. There are 8 APs and STAs from 4, 8, 16, 32, and 64, which always
287 have data in the data queue to transmit to the associated AP (i.e., saturated conditions). Among the STAs, we
288 specifically designed a delay scenario where the inter-BSS NAVs of specific target STA is successively set due to
289 neighboring APs which successively initiate transmission with their associated STAs. Thus, the STA is delayed
290 because of successive set of inter-BSS NAVs, and has to wait until all of them are reset. The delay duration of the
291 STA can be minimized to nearly wait time threshold θ_{ed} in Algorithm 2 by taking advantage of our proposed
292 algorithms above.

Table 4. Simulation parameters

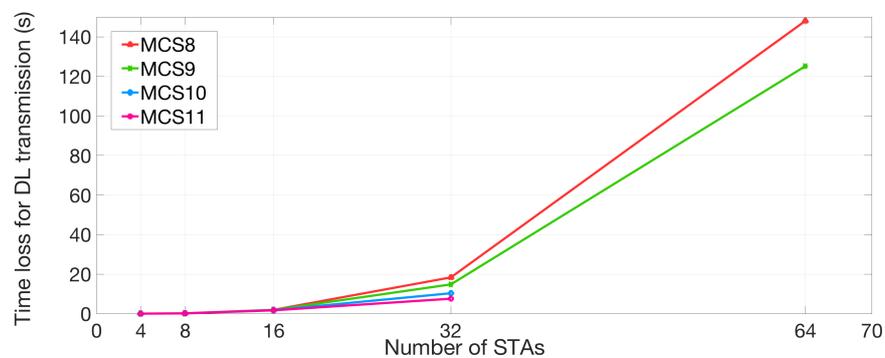
Parameter	Description
B	160MHz
FFT	256
L_d	1460 bytes
P_a	256
CW_{\min}	32
CW_{\max}	1024
SIFS	16 μs
aSlotTime	9 μs
ρ	16 μs
N_A	8
N_S	from 8 up to 64

293 The total bandwidth of a carrier is 160MHz and FFT size is adopted as 256. In addition, the aggregated
 294 medium access control protocol data unit (A-MPDU) consists of P_a packets with symbol duration in the length
 295 of ρ (16 μs) is used for transmission simulations. The size of l_c including ACK, SIFS, and other control frames is
 296 set to 123bytes and l_d is 1460bytes. In addition, P_a frames in the size of l_d are aggregated as an A-MPDU for
 297 simulations. Moreover, the AP with 8 antennas serves STAs that each of them is equipped with 2 antennas for
 298 enabling transmission through multiple spatial streams.

299 First, we simulate the time loss in U/DL transmission under hyper-dense deployment scenarios varying
 300 modulation and coding scheme (MCS) levels and the number of STAs which delay an STA such as STA_2 in
 301 Fig. 5. Furthermore, dense constellations which supports up to 32 STAs are considered to evaluate the stability of
 302 time loss for the delayed STA. In addition, the fairness of the medium access for each STA is compared with
 303 DCF by calculating Jain's index.

304 4.2. Simulation results and discussions

305 The l_t values of DL and UL MU transmission are simulated in dense deployment environments under the
 306 various MCS and STAs numbers, as presented in Fig. 7 and Fig. 8. Based on (3), the time loss of STA is measured
 307 using the parameters in Table 4. The time loss of STA is sharply increased as the number of neighboring STAs
 308 is added. As the number of interfering STAs reaches 64, the STA which uses MCS9 and MCS10 cannot transmit
 309 any data in effect. The reason for this phenomenon is caused by multi-NAV's policy of IEEE 802.11ax where the
 310 inter-BSS NAVs are increased to manage the possible congestion for STA. Furthermore, the inter-BSS NAVs are
 311 proportional to the number of neighboring or interfering APs and associated STAs.

**Figure 7.** Expected time loss l_t for DL transmission.

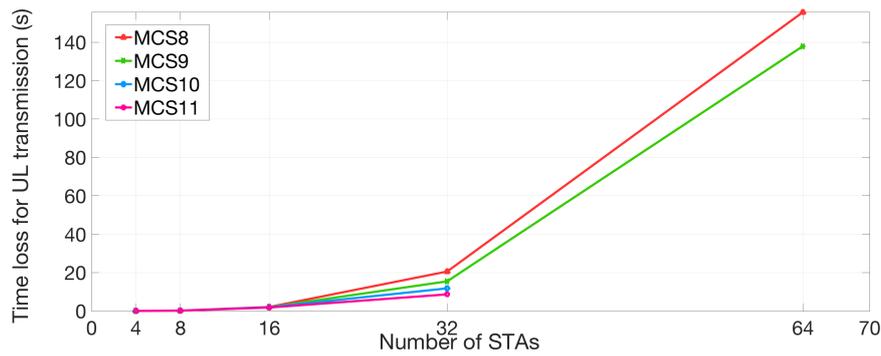


Figure 8. Expected time loss l_t for UL transmission.

312 Finally, Fig. 9 and Fig. 10 represent a comparison of fairness index values for medium access between DCF
 313 and our proposed algorithm. The fairness of our proposed algorithm is improved to 44% and 31% compared to
 314 DCF when the numbers of STAs are 16 and 64 respectively in DL transmission. In case of UL transmission,
 315 the fairness is improved up to 50% and 36% when the numbers of interfering STAs are 16 and 64 respectively.
 316 In conclusions, our proposed algorithm enhances the fairness of the channel access efficiency in hyper-densely
 317 deployed environments in order to minimize the delay time.

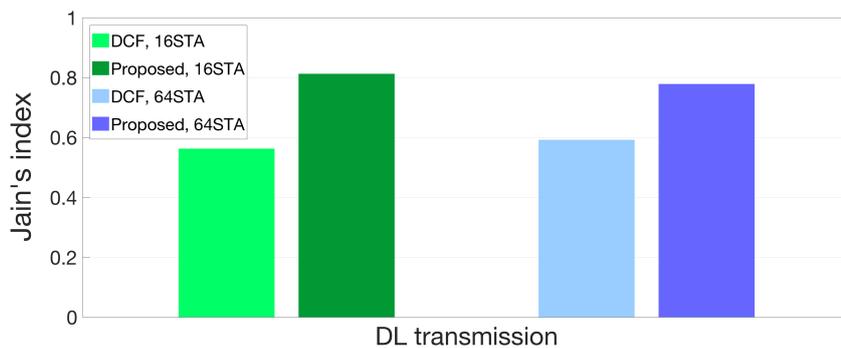


Figure 9. Fairness comparison between DCF and proposed protocol in DL scenario

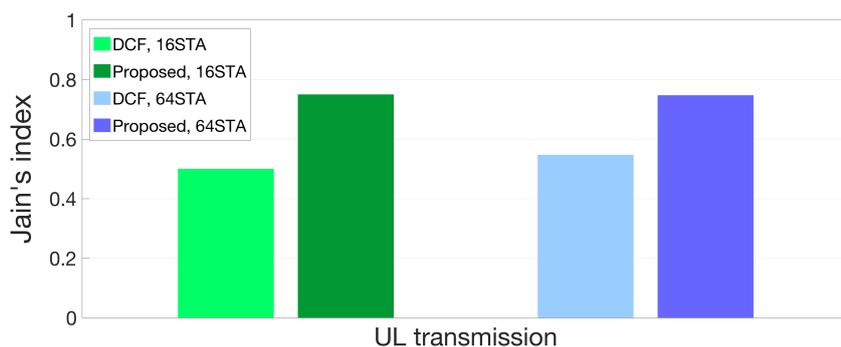


Figure 10. Fairness comparison between DCF and proposed protocol in UL scenario

318 5. Conclusion Remarks

319 In this paper, we demonstrated the possible delay problem of IEEE 802.11ax in a hyper-dense deployment of
 320 IoT wireless devices due to multiple inter-BSS NAVs settings and analyzed the expected lost time and throughput
 321 of victim STAs. To address this issue, we first propose a swift and low-complexity beam direction selection
 322 algorithm based on an opportunistic beamforming for IEEE 802.11ax-based IoT systems in terms of the joint
 323 optimization of delay reduction and interference mitigation. In addition, we evaluated the performance of our

324 proposed algorithm in terms of the throughput and fairness, and clearly show that our proposed algorithm can
325 achieve a desired performance.

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331 **Conflicts of Interest:** The authors declare no conflict of interest.

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