

Article

Rényi Entropy Power Inequalities via Normal Transport and Rotation

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Abstract: Following a recent proof of Shannon's entropy power inequality (EPI), a comprehensive framework for deriving various EPIs for the Rényi entropy is presented, that uses transport arguments from normal densities and a change of variable by rotation.

Keywords: Rényi entropy; Entropy power inequalities; Transportation Arguments; Normal distributions; Escort distributions; Log-concave distributions

1. Introduction

The entropy power inequality (EPI) dates back to Shannon's seminal paper [1] and has a long history [2]. The link with the Rényi entropy was first made by Dembo, Cover and Thomas [3] in connection with Young's convolutional inequality with sharp constants, where Shannon's EPI is obtained by letting the Rényi entropy orders tend to one [4, Theorem 17.8.3].

Recently there has been significant interest in Rényi EPIs for several independent variables. Bobkov and Chistyakov [5] extended the classical Shannon's EPI to the Rényi entropy by incorporating a multiplicative constant that depends on the order of the Rényi entropy. Ram and Sason [6] improved the constant by making it depend also on the number of variables. Even more recently, Bobkov and Marsiglietti [7] proved another modification of the EPI for the Rényi entropy for two independent variables, with a power exponent parameter α which was improved by Li [8]. All these EPIs were found for Rényi entropies of orders > 1 . The α -modification of the Rényi EPI was extended to orders < 1 for two independent variables having log-concave densities by Marsiglietti and Melbourne [9]. The starting point of all the above works was Young's strengthened convolutional inequality.

Also recently, Shannon's original EPI was proved anew [10] using a simple transport argument from normal variables and a change of variable by rotation. In this paper, we exploit these ingredients, described in the following lemmas, to establish all previously known Rényi EPIs and derive new ones.

Notation 1. Throughout this article the considered zero-mean random variables $X \in \mathbb{R}^n$ admit a density which is implicitly assumed continuous inside its support. Write $X \sim f$ if X has density f .

Lemma 1 (Normal Transport [10]). Let $X^* \sim \mathcal{N}(0, \sigma^2 \mathbf{I})$. There exists a diffeomorphism $T : \mathbb{R}^n \rightarrow \mathbb{R}^n$ such that $X = T(X^*) \sim f$. Moreover, T can be chosen such that its Jacobian matrix T' is (lower) triangular with positive diagonal elements.

Two different proofs of Lemma 1 are given in [10]. The proof is very simple for one-dimensional variables [11], where T is just an increasing function with continuous derivative $T' > 0$.

Lemma 2 (Normal Rotation [10]). If X^*, Y^* are i.i.d. normal, then for any $0 < \lambda < 1$, the rotation

$$\begin{pmatrix} \tilde{X} \\ \tilde{Y} \end{pmatrix} = \begin{pmatrix} \sqrt{\lambda} & \sqrt{1-\lambda} \\ -\sqrt{1-\lambda} & \sqrt{\lambda} \end{pmatrix} \begin{pmatrix} X^* \\ Y^* \end{pmatrix} \quad (1)$$

31 yields i.i.d. normal variables \tilde{X}, \tilde{Y} .

32 Notice that the the starred variables can be expressed in terms of the tilde variables by the inverse
33 rotation

$$\begin{pmatrix} X^* \\ Y^* \end{pmatrix} = \begin{pmatrix} \sqrt{\lambda} & -\sqrt{1-\lambda} \\ \sqrt{1-\lambda} & \sqrt{\lambda} \end{pmatrix} \begin{pmatrix} \tilde{X} \\ \tilde{Y} \end{pmatrix}. \quad (2)$$

34 The proof of Lemma 2 is trivial considering covariance matrices. A deeper result states that this
35 property of remaining i.i.d. by rotation characterizes the normal distribution—this is known as
36 Bernstein’s lemma (see e.g., [12, Chap. 5], [11, Lemma 4]). This explains why one obtains equality in
37 the EPI only for normal variables.

38 This article is a revised, full version of what was presented in part in a previous conference
39 communication [13]. It is organized as follows. Preliminary definitions and known properties are
40 presented in Section 2. Section 3 derives a crucial “information inequality” for Rényi entropies that
41 enjoys a transformational invariance. The central result is in Section 4, where the first version of the
42 Rényi EPI by Dembo, Cover and Thomas is proved using the ingredients of Lemmas 1 and 2. All
43 previously known Rényi EPIs—and new ones—are then derived using a simple method in Section 5.
44 Section 6 concludes.

45 2. Preliminary Definitions and Properties

46 Throughout this article we consider exponents $p > 0$ with $p \neq 1$. The following definition is well
47 known and used e.g., in Hölder’s inequality.

48 **Definition 1** (Conjugate Exponent). *The conjugate exponent of p is*

$$p' = \frac{p}{p-1}, \quad (3)$$

49 that is, the number p' such that

$$\frac{1}{p} + \frac{1}{p'} = 1. \quad (4)$$

50 **Remark 1.** *There are two situations depending on whether p' is positive or negative, as summarized in the*
51 *following table.*

$$\begin{array}{c|c} p > 1 & 0 < p < 1 \\ \hline p' > 1 & p' < 0 \end{array}$$

53 **Definition 2** (Rényi Entropy). *If X has density $f \in L^p(\mathbb{R}^n)$, its Rényi entropy of order p is defined by*

$$h_p(X) = \frac{1}{1-p} \log \int_{\mathbb{R}^n} f^p(x) dx \quad (5)$$

$$= -p' \log \|f\|_p \quad (6)$$

54 It is known that the limit as $p \rightarrow 1$ is the Shannon entropy

$$h_1(X) = - \int_{\mathbb{R}^n} f(x) \log f(x) dx. \quad (7)$$

55 The Rényi entropy enjoys well-known properties similar to those of the Shannon entropy, which are
56 recalled here for completeness.

57 **Lemma 3** (Scaling Property). For any $a \in \mathbb{R}$,

$$h_p(aX) = h_p(X) + n \log |a|. \quad (8)$$

58 **Proof.** Making a change of variables, $h_p(aX) = \frac{1}{1-p} \log \int \left(\frac{1}{|a|^n} f\left(\frac{x}{a}\right) \right)^p dx = \frac{1}{1-p} \log \int f^p\left(\frac{x}{a}\right) \frac{dx}{|a|^n} +$
59 $\frac{1}{1-p} \log |a|^{n(1-p)} = h_p(X) + n \log |a|. \quad \square$

60 One recovers the usual scaling property for the Shannon entropy by letting $p \rightarrow 1$.

61 **Lemma 4** (Rényi Entropy of the Normal). If $X^* \sim \mathcal{N}(0, \mathbf{K})$ for some nonsingular covariance matrix \mathbf{K} ,
62 then

$$h_p(X^*) = \frac{1}{2} \log((2\pi)^n |\mathbf{K}|) + \frac{n}{2} p' \frac{\log p}{p} \quad (9)$$

63 where $|\cdot|$ denotes the determinant. In particular for $X^* \sim \mathcal{N}(0, \sigma^2 \mathbf{I})$,

$$h_p(X^*) = \frac{n}{2} \log(2\pi\sigma^2) + \frac{n}{2} p' \frac{\log p}{p}. \quad (10)$$

64 **Proof.** By direct calculation, $h_p(X^*) = \frac{1}{1-p} \log \int \left(\frac{\exp(-\frac{1}{2}x^t \mathbf{K}^{-1}x)}{\sqrt{(2\pi)^n |\mathbf{K}|}} \right)^p dx = \frac{1}{1-p} \log \frac{\sqrt{(2\pi)^n |\mathbf{K}|} p^{-n}}{\sqrt{(2\pi)^n |\mathbf{K}|}^p} =$
65 $\frac{1}{2} \log((2\pi)^n |\mathbf{K}|) - \frac{n}{2} \frac{\log p}{1-p}. \quad \square$

66 Again one recovers the Shannon entropy of a normal variable by letting $p \rightarrow 1$ (then $p' \frac{\log p}{p} \rightarrow$
67 $\log e$).

68 The following notion of escort distribution [14,15] is useful in the sequel.

69 **Definition 3** (Escort Density [15, § 2.2]). If $f \in L^p(\mathbb{R}^n)$, its escort density of exponent p is the density
70 defined by

$$f_p(x) = \frac{f^p(x)}{\int_{\mathbb{R}^n} f^p(x) dx}. \quad (11)$$

71 In other words $f_p = \frac{f^p}{\|f\|_p^p}$ where $\|f\|_p$ denotes the L^p norm of f . We also use the notation X_p to denote the
72 corresponding escort random variable with density f_p .

73 **Lemma 5** (Monotonicity Property). If $p < q$ then $h_p(X) \geq h_q(X)$ with equality if and only if X is uniformly
74 distributed.

75 **Proof.** Let $p \neq 1$ and assume that $f \in L^q(\mathbb{R}^n)$ for all q in a neighborhood of p so that one can freely
76 differentiate under the integral sign:

$$\frac{\partial}{\partial p} h_p(X) = \frac{1}{(1-p)^2} \log \int f^p + \frac{1}{1-p} \frac{\int f^p \log f}{\int f^p} \quad (12)$$

$$= \frac{1}{(1-p)^2} \left(\log \int f^p + \int f_p \log f^{1-p} \right) \quad (13)$$

$$= \frac{1}{(1-p)^2} \int f_p \log \frac{f}{f_p} \quad (14)$$

$$= -\frac{D(f_p \| f)}{(1-p)^2} \leq 0 \quad (15)$$

77 where $D(\cdot \| \cdot)$ denotes the Kullback-Leibler divergence. Equality $D(f_p \| f) = 0$ can hold only if $f = f_p$
 78 a.e., which since $p \neq 1$ implies that f is constant over some measurable subset $A \subset \mathbb{R}^n$, and is zero
 79 elsewhere. It follows that $h_p(X) > h_q(X)$ for any $p < q$ if X is not uniformly distributed. Conversely,
 80 if X is uniformly distributed over some measurable subset $A \subset \mathbb{R}^n$, its density can be written as
 81 $f(x) = 1/\text{vol}(A)$ for all $x \in A$ and $= 0$ elsewhere. Then $h_p(X) = \frac{1}{1-p} \log \frac{\text{vol}(A)}{\text{vol}(A)^p} = \log \text{vol}(A)$ is
 82 independent of p . \square

83 **Remark 2.** *The identity established in the proof:*

$$\frac{\partial}{\partial p} h_p(X) = -\frac{D(f_p \| f)}{(1-p)^2} \quad (16)$$

84 *seems new, although a similar formula for discrete variables can be found in [16, § 5.3].*

85 3. An Information Inequality

86 The Shannon entropy satisfies a fundamental “information inequality” [4, Theorem 2.6.3] from
 87 which many classical information-theoretic inequalities can be derived. This can be written as

$$h_1(X) \leq -\mathbb{E} \log \varphi(X) \quad (17)$$

88 for any density φ , with equality if and only if $\varphi = f$ a.e. The following Theorem can be seen as the
 89 natural extension of the information inequality to Rényi entropies and is central in the following
 90 derivations of this paper. J.F. Bercher has pointed out to the author that it is similar to an inequality for
 91 discrete distributions established by Campbell [17] in the context of source coding (see also [14]).

92 **Theorem 1** (Information Inequality). *For any density φ ,*

$$h_p(X) \leq -p' \log \mathbb{E}(\varphi^{1/p'}(X)) \quad (18)$$

93 *with equality if and only if $\varphi = f_p$ a.e.*

94 By letting $p \rightarrow 1$ one recovers the classical information inequality (17) for the Shannon entropy.

95 **Proof.** By definition (6),

$$h_p(X) = -p' \log \|f\|_p \quad (19)$$

$$= -p' \log \left(\int (f^p \varphi^{-1}) \varphi \right)^{1/p} \quad (20)$$

$$\leq -p' \log \int (f^p \varphi^{-1})^{1/p} \varphi \quad (21)$$

$$= -p' \log \int f \varphi^{1/p'} \quad (22)$$

96 where the inequality follows from Jensen’s inequality applied to the function $x \mapsto x^{1/p}$, which is
 97 strictly concave if $p > 1$ (that is, $p' > 0$) and strictly convex if $p < 1$ (that is, $p' < 0$). Equality holds if
 98 and only if $f^p \varphi^{-1}$ is constant a.e., which means that φ and f^p are proportional a.e. Normalizing gives
 99 the announced condition $\varphi = f_p$ a.e. \square

100 **Remark 3.** *An alternate proof is obtained using Hölder’s inequality or its reverse applied to f and $\varphi^{1/p'}$. Notice*
 101 *that the equality case for $\varphi = f_p$ gives*

$$h_p(X) = -p' \log \mathbb{E}(f_p^{1/p'}(X)) \quad (23)$$

102 as can be easily checked directly.

103 The following conditional version of Theorem 1 involves a more complicated relation for
104 dependent variables.

105 **Corollary 1** (Conditional Information Inequality). *For any two random variables $X, Y \in \mathbb{R}^n$,*

$$-p' \log \mathbb{E}_Y \exp(-h_p(X|Y)/p') \leq -p' \log \mathbb{E}(\varphi^{1/p'}(X|Y)) \quad (24)$$

106 where $h_p(X|y)$ denotes the Rényi entropy of X knowing $Y = y$ and the expectation on the l.h.s. is taken over Y
107 (the expectation in the r.h.s. is taken over (X, Y)).

108 In particular, when X and Y are independent,

$$h_p(X) \leq -p' \log \mathbb{E}(\varphi^{1/p'}(X|Y)). \quad (25)$$

109 with equality if and only if $\varphi(x|y)$ does not depend on y and equals $f_p(x)$ a.e.

110 **Proof.** From (18) for fixed y , one has $\mathbb{E}(\varphi^{1/p'}(X|y)) \leq \exp(-h_p(X|y)/p')$ for $p' > 0$ ($p > 1$) and the
111 opposite inequality for $p' < 0$ ($p < 1$), with equality if and only if $\varphi(x|y) = f_p(x|y)$ a.e. Taking the
112 expectation over Y yields $\mathbb{E}(\varphi^{1/p'}(X|Y)) \leq \mathbb{E}_Y \exp(-h_p(X|Y)/p')$ for $p' > 0$ ($p > 1$) and the opposite
113 inequality for $p' < 0$ ($p < 1$). The result follows by taking the logarithm and multiplying by $-p'$.
114 When X and Y are independent, equality holds if and only if $\varphi(x|y) = f_p(x)$ a.e. for all y . \square

For the Shannon entropy, the difference between the two sides of the information inequality (17) is the Kullback-Leibler divergence:

$$D(f\|\varphi) = \mathbb{E} \log \frac{f(X)}{\varphi(X)} \geq 0$$

115 which can also be noted $D(X\|Z)$ where $X \sim f$ and $Z \sim \varphi$. It is known (and easy to check) that
116 the divergence is invariant by reversible transformations T . This means that when $X = T(X^*)$ and
117 $Z = T(Z^*)$, one has $D(X\|Z) = D(X^*\|Z^*)$. A natural extension to Rényi entropies can be obtained on
118 the difference

$$-p' \log \mathbb{E}(\varphi^{1/p'}(X)) - h_p(X) \geq 0 \quad (26)$$

119 between the two sides of the information inequality (18).

120 **Theorem 2** (Transformational Invariance). *Let $T : \mathbb{R}^n \rightarrow \mathbb{R}^n$ be a diffeomorphism and suppose that*

$$X_p = T(X_p^*) \quad (27)$$

$$Z = T(Z^*) \quad (28)$$

121 where $Z \sim \varphi$ and $Z^* \sim \varphi^*$. Then

$$-p' \log \mathbb{E}(\varphi^{1/p'}(X)) - h_p(X) = -p' \log \mathbb{E}(\varphi^{*1/p'}(X^*)) - h_p(X^*). \quad (29)$$

122 Note that from (6), this identity can be rewritten as

$$\frac{\mathbb{E}(\varphi^{*1/p'}(X^*))}{\|f^*\|_p} = \frac{\mathbb{E}(\varphi^{1/p'}(X))}{\|f\|_p} \quad (30)$$

123 **Proof.** Proceed to prove (30). Let f, f^* be the respective densities of X, X^* and recall that $X_p \sim f_p$ and
 124 $X_p^* \sim f_p^*$. By the transformation T the densities are related by

$$f_p^*(x^*) = f_p(T(x^*))|T'(x^*)| \quad (31)$$

$$\varphi^*(x^*) = \varphi(T(x^*))|T'(x^*)|. \quad (32)$$

125 where $|T'|$ denotes the Jacobian determinant of T . Using these relations and Definition 3,

$$\mathbb{E}(\varphi^{*1/p'}(X^*)) / \|f^*\|_p = \mathbb{E}(\varphi^{1/p'}(T(X^*))|T'(X^*)|^{1/p'}) / \|f^*\|_p \quad (33)$$

$$= \int \varphi^{1/p'}(T(x^*))|T'(x^*)|^{1/p'} f^*(x^*) dx^* / \|f^*\|_p \quad (34)$$

$$= \int \varphi^{1/p'}(T(x^*))|T'(x^*)|^{1/p'} f_p^*(x^*)^{1/p} dx^* \quad (35)$$

$$= \int \varphi^{1/p'}(T(x^*)) f_p(T(x^*))^{1/p} |T'(x^*)| dx^* \quad (36)$$

$$= \int \varphi^{1/p'}(x) f_p(x)^{1/p} dx \quad (37)$$

$$= \mathbb{E}(\varphi^{1/p'}(X)) / \|f\|_p \quad \square \quad (38)$$

126 **Remark 4.** The fact that φ is a density was not used in the proof of Theorem 2. Therefore (29) holds more
 127 generally for any function φ satisfying (32).

128 4. First Version of the Rényi EPI

129 For two independent random variables X and Y , the Shannon entropy power inequality can be
 130 expressed as follows [2,3]: For any $0 < \lambda < 1$,

$$h(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) \geq \lambda h(X) + (1-\lambda)h(Y) \quad (39)$$

131 with equality if and only if X, Y are i.i.d. normal. That is, the difference $h(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) - \lambda h(X) -$
 132 $(1-\lambda)h(Y)$ is minimum (zero) for i.i.d. normal X, Y . In this section, we study the natural generalization
 133 for Rényi entropies [3, Theorem 12], namely that the quantity

$$h_r(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) - \lambda h_p(X) - (1-\lambda)h_q(Y) \quad (40)$$

134 is minimum for i.i.d. normal X, Y . Here the triple (p, q, r) and its associated λ satisfy the following
 135 condition, which is used e.g., in Young's convolutional inequality.

136 **Definition 4** (Exponent Triple). An exponent triple $(p, q, r)_\lambda$ has conjugates p', q', r' of the same sign and
 137 such that

$$\frac{1}{p'} + \frac{1}{q'} = \frac{1}{r'}. \quad (41)$$

138 The corresponding coefficient $\lambda \in (0, 1)$ is defined by

$$\lambda = \frac{r'}{p'} = 1 - \frac{r'}{q'} \quad (42)$$

139 In other words, the exponents p, q, r are such that

$$\frac{1}{p} + \frac{1}{q} = 1 + \frac{1}{r} \quad (43)$$

140 and fulfill one the following two conditons:

$$141 \quad \begin{array}{c} \hline p, q, r > 1 \quad \parallel \quad 0 < p, q, r < 1 \\ p', q', r' > 1 \quad \parallel \quad p', q', r' < 0 \\ \hline r' < p', r' < q' \quad \parallel \quad |r'| < |p'|, |r'| < |q'| \\ r > p, r > q \quad \parallel \quad r < p, r < q \\ \hline \end{array}$$

142 The key argument used in this section is the following. If $X \sim f$ and $Y \sim g$, then for the escort
143 variables, $X_p \sim f_p$ and $Y_q \sim g_q$. By normal transport (Lemma 1), one can write

$$X_p = T(X_p^*) \quad (44)$$

$$Y_q = U(Y_q^*) \quad (45)$$

144 for two diffeomorphisms T and U , where X^*, Y^* are, say, i.i.d. standard normal $\mathcal{N}(0, \mathbf{I})$. (It follows
145 that $X_p^* \sim \mathcal{N}(0, \mathbf{I}/p)$ and $Y_q^* \sim \mathcal{N}(0, \mathbf{I}/q)$.) We then have the following straightforward extension of
146 Theorem 2:

147 **Lemma 6** (Transformational Invariance for Two Independent Variables). *For a two-dimensional $\varphi(x, y)$,*

$$\begin{aligned} -r' \log \mathbb{E}(\varphi^{1/r'}(X, Y)) - \lambda h_p(X) - (1 - \lambda) h_q(Y) \\ = -r' \log \mathbb{E}(\varphi^{*1/r'}(X^*, Y^*)) - \lambda h_p(X^*) - (1 - \lambda) h_q(Y^*) \end{aligned} \quad (46)$$

148 where

$$\varphi^*(x^*, y^*) = \varphi(T(x^*), U(y^*)) |T'(x^*)|^\lambda |U'(y^*)|^{1-\lambda}. \quad (47)$$

149 **Proof.** From (6) and the definition of λ , (57) can be rewritten as

$$\frac{\mathbb{E}(\varphi^{*1/r'}(X^*, Y^*))}{\|f^*\|_p \|g^*\|_q} = \frac{\mathbb{E}(\varphi^{1/r'}(X, Y))}{\|f\|_p \|g\|_q} \quad (48)$$

150 By the transformations T and U the densities of the escort variables are related by $f_p^*(x^*) =$
151 $f_p(T(x^*)) |T'(x^*)|$ and $g_q^*(y^*) = g_q(U(y^*)) |U'(y^*)|$. Now by the same calculation as in the proof
152 of Theorem 2,

$$\frac{\mathbb{E}(\varphi^{*1/r'}(X^*, Y^*))}{\|f^*\|_p \|g^*\|_q} = \frac{\mathbb{E}(\varphi^{1/r'}(T(X^*), U(Y^*)) |T'(X^*)|^{1/p'} |U'(Y^*)|^{1/q'})}{\|f^*\|_p \|g^*\|_q} \quad (49)$$

$$= \frac{\int \varphi^{1/r'}(T(x^*), U(y^*)) |T'(x^*)|^{1/p'} |U'(y^*)|^{1/q'} f^*(x^*) g^*(y^*) dx^* dy^*}{\|f^*\|_p \|g^*\|_q} \quad (50)$$

$$= \int \varphi^{1/r'}(T(x^*), U(y^*)) |T'(x^*)|^{1/p'} |U'(y^*)|^{1/q'} f_p^*(x^*)^{1/p} g_q^*(y^*)^{1/q} dx^* dy^* \quad (51)$$

$$= \int \varphi^{1/r'}(T(x^*), U(y^*)) |T'(x^*)| |U'(y^*)| f_p(T(x^*))^{1/p} g_q(U(y^*))^{1/q} dx^* dy^* \quad (52)$$

$$= \int \varphi^{1/r'}(x, y) f_p(x)^{1/p} g_q(y)^{1/q} dx dy \quad (53)$$

$$= \frac{\mathbb{E}(\varphi^{1/r'}(X, Y))}{\|f\|_p \|g\|_q} \quad \square \quad (54)$$

153 **Lemma 7.** Let φ be the density of $\sqrt{\lambda}X + \sqrt{1-\lambda}Y$. Then

$$\begin{aligned} & h_r(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) - \lambda h_p(X) - (1-\lambda)h_q(Y) \\ & \geq -r' \log \mathbb{E} \left\{ (\varphi_r(\sqrt{\lambda}T(X^*) + \sqrt{1-\lambda}U(Y^*)) \cdot |\lambda T'(X^*) + (1-\lambda)U'(Y^*)|)^{1/r'} \right\} \\ & \quad - \lambda h_p(X^*) - (1-\lambda)h_q(Y^*). \end{aligned} \quad (55)$$

154 **Proof.** By the equality case of Theorem 1 (see (23)), one has

$$h_r(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) = -r' \log \mathbb{E} \varphi_r^{1/r'}(\sqrt{\lambda}X + \sqrt{1-\lambda}Y). \quad (56)$$

155 Now by Lemma 6 applied to $\varphi(x, y) = \varphi_r(\sqrt{\lambda}x + \sqrt{1-\lambda}y)$, we have $\varphi^*(x, y) = \varphi_r(\sqrt{\lambda}T(x^*) +$
156 $\sqrt{1-\lambda}U(y^*))|T'(x^*)|^\lambda|U'(y^*)|^{1-\lambda}$, and, therefore,

$$\begin{aligned} & h_r(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) - \lambda h_p(X) - (1-\lambda)h_q(Y) \\ & = -r' \log \mathbb{E} \left\{ (\varphi_r(\sqrt{\lambda}T(X^*) + \sqrt{1-\lambda}U(Y^*)) \cdot |T'(X^*)|^\lambda|U'(Y^*)|^{1-\lambda})^{1/r'} \right\} \\ & \quad - \lambda h_p(X^*) - (1-\lambda)h_q(Y^*). \end{aligned} \quad (57)$$

157 Since from Lemma 1, T and U can be chosen such that T' and U' are (lower) triangular with positive
158 diagonal elements, it follows easily from the arithmetic-geometric mean inequality that

$$|T'(X^*)|^\lambda|U'(Y^*)|^{1-\lambda} \leq |\lambda T'(X^*) + (1-\lambda)U'(Y^*)|. \quad (58)$$

159 The result follows at once (for either positive or negative r'). \square

160 We can now use the normal rotation Lemma 2 to conclude by proving the following

161 **Theorem 3 (Rényi EPI [3]).** For independent X, Y and exponent triple $(p, q, r)_\lambda$,

$$h_r(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) - \lambda h_p(X) - (1-\lambda)h_q(Y) \geq \frac{n}{2} r' \left(\frac{\log r}{r} - \frac{\log p}{p} - \frac{\log q}{q} \right) \quad (59)$$

162 with equality if and only if X, Y are i.i.d. normal.

163 **Proof.** If X, Y are i.i.d. normal, then $\sqrt{\lambda}X + \sqrt{1-\lambda}Y$ is also identically distributed as X and Y , and
164 from Lemma 4, it is immediate to check that equality holds (irrespective of their covariances). Therefore,
165 inequality (59) is equivalent to

$$h_r(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) - \lambda h_p(X) - (1-\lambda)h_q(Y) \geq h_r(\sqrt{\lambda}X^* + \sqrt{1-\lambda}Y^*) - \lambda h_p(X^*) - (1-\lambda)h_q(Y^*) \quad (60)$$

166 where X^*, Y^* are, say, i.i.d. standard normal $\mathcal{N}(0, \mathbf{I})$.

167 To prove (60), consider the normal rotation of Lemma 2 and write X^*, Y^* in terms of \tilde{X}, \tilde{Y} using (2)
168 in the first term of the r.h.s. of (55) (Lemma 7). One obtains:

$$\begin{aligned} & h_r(\sqrt{\lambda}X + \sqrt{1-\lambda}Y) - \lambda h_p(X) - (1-\lambda)h_q(Y) \\ & \geq -r' \log \mathbb{E} \left\{ (\psi(\tilde{X}|\tilde{Y}))^{1/r'} \right\} - \lambda h_p(X^*) - (1-\lambda)h_q(Y^*). \end{aligned} \quad (61)$$

169 where

$$\begin{aligned} \psi(\tilde{x}|\tilde{y}) &= \varphi_r(\sqrt{\lambda}T(\sqrt{\lambda}\tilde{x} - \sqrt{1-\lambda}\tilde{y}) + \sqrt{1-\lambda}U(\sqrt{1-\lambda}\tilde{x} + \sqrt{\lambda}\tilde{y})) \\ &\quad \times |\lambda T'(\sqrt{\lambda}\tilde{x} - \sqrt{1-\lambda}\tilde{y}) + (1-\lambda)U'(\sqrt{1-\lambda}\tilde{x} + \sqrt{\lambda}\tilde{y})| \end{aligned} \quad (62)$$

170 Making the change of variable $z = \sqrt{\lambda}T(\sqrt{\lambda}\tilde{x} - \sqrt{1-\lambda}\tilde{y}) + \sqrt{1-\lambda}U(\sqrt{1-\lambda}\tilde{x} + \sqrt{\lambda}\tilde{y})$, one obtains

$$\int \psi(\tilde{x}|\tilde{y}) d\tilde{x} = \int \varphi_r(z) dz = 1, \quad (63)$$

171 since φ_r is a density. Hence $\psi(\tilde{x}|\tilde{y})$ is also a density in \tilde{x} for fixed \tilde{y} . Now since by Lemma 2, \tilde{X} and \tilde{Y}
172 are independent, by the conditional information inequality (25) of Corollary 1, one has

$$-r' \log \mathbb{E}\left\{(\psi(\tilde{X}|\tilde{Y}))^{1/r'}\right\} \geq h_r(\tilde{X}) = h_r(\sqrt{\lambda}X^* + \sqrt{1-\lambda}Y^*). \quad (64)$$

173 Combining with (61) yields the announced inequality (60).

174 It remains to settle the equality case in (60). From the above proof, equality holds in (60) if and
175 only if both (58) and (64) are equalities. Equality in (58) holds if and only if for all $i = 1, 2, \dots, n$,

$$\frac{\partial T_i}{\partial x_i}(X^*) = \frac{\partial U_i}{\partial y_i}(Y^*) \text{ a.s.} \quad (65)$$

176 Since X^* and Y^* are independent normal variables, this implies that $\frac{\partial T}{\partial x_i}$ and $\frac{\partial U}{\partial y_i}$ are constant and equal.

177 In particular the Jacobian $|\lambda T'(\sqrt{\lambda}\tilde{x} - \sqrt{1-\lambda}\tilde{y}) + (1-\lambda)U'(\sqrt{1-\lambda}\tilde{x} + \sqrt{\lambda}\tilde{y})|$ in (62) is constant.

178 From Corollary 1 equality in (64) holds if and only if $\psi(\tilde{x}|\tilde{y})$ does not depend on \tilde{y} , which implies
179 that $\sqrt{\lambda}T(\sqrt{\lambda}\tilde{x} - \sqrt{1-\lambda}\tilde{y}) + \sqrt{1-\lambda}U(\sqrt{1-\lambda}\tilde{x} + \sqrt{\lambda}\tilde{y})$ does not depend on the value of \tilde{y} . Taking
180 derivatives with respect to y_j for all $j = 1, 2, \dots, n$,

$$-\sqrt{\lambda}\sqrt{1-\lambda}\frac{\partial T_i}{\partial x_j}(\sqrt{\lambda}\tilde{X} - \sqrt{1-\lambda}\tilde{Y}) + \sqrt{\lambda}\sqrt{1-\lambda}\frac{\partial U_i}{\partial x_j}(\sqrt{1-\lambda}\tilde{X} + \sqrt{\lambda}\tilde{Y}) = 0 \quad (66)$$

181 which implies

$$\frac{\partial T_i}{\partial x_j}(X^*) = \frac{\partial U_i}{\partial y_j}(Y^*) \text{ a.s.} \quad (67)$$

182 for all $i, j = 1, 2, \dots, n$. Therefore, T and U are linear transformations, equal up to an additive constant
183 (equal to 0 since all variables are assumed of zero mean). It follows that $X_p = T(X_p^*)$ and $Y_q = U(Y_q^*)$
184 are normal with respective distributions $X_p \sim \mathcal{N}(0, \mathbf{K}/p)$ and $Y_q \sim \mathcal{N}(0, \mathbf{K}/q)$. Hence X and Y are
185 i.i.d. normal $\mathcal{N}(0, \mathbf{K})$. \square

186 A straightforward generalization to several independent variables is the following

187 **Corollary 2** (Rényi EPI for Several Variables). Let r_1, r_2, \dots, r_m, r be exponents those conjugates
188 $r'_1, r'_2, \dots, r'_m, r'$ are of the same sign and satisfy

$$\sum_{i=1}^m \frac{1}{r'_i} = \frac{1}{r'} \quad (68)$$

189 and let $\lambda_1, \lambda_2, \dots, \lambda_m$ be defined by

$$\lambda_i = \frac{r'}{r'_i} \quad (i = 1, 2, \dots, m). \quad (69)$$

190 Then for independent X_1, X_2, \dots, X_m ,

$$h_r\left(\sum_{i=1}^m \sqrt{\lambda_i} X_i\right) - \sum_{i=1}^m \lambda_i h_{r_i}(X_i) \geq \frac{n}{2} r' \left(\frac{\log r}{r} - \sum_{i=1}^m \frac{\log r_i}{r_i} \right) \quad (70)$$

191 with equality if and only if the X_i are i.i.d. normal.

192 **Proof.** By induction on m : The result for $m = 2$ is Theorem 3. Suppose the result satisfied at order $m - 1$
 193 and let $Y_m = \sum_{i=1}^{m-1} \sqrt{\lambda_i} X_i / \sqrt{1 - \lambda_m}$ and s_m be such that $\frac{1}{s_m} = \sum_{i=1}^{m-1} \frac{1}{r_i}$. Notice that $\frac{1}{r'} = \frac{1}{r_m} + \frac{1}{s_m} =$
 194 $\frac{\lambda_m}{r'} + \frac{1}{s_m}$, hence $r' = (1 - \lambda_m) s_m'$. By Theorem 3, $h_r\left(\sum_{i=1}^m \sqrt{\lambda_i} X_i\right) = h_r(\sqrt{\lambda_m} X_m + \sqrt{1 - \lambda_m} Y_m) \geq$
 195 $\lambda_m h_{r_m}(X_m) + (1 - \lambda_m) h_{s_m}(Y_m) + \frac{n}{2} r' \left(\frac{\log r}{r} - \frac{\log r_m}{r_m} - \frac{\log s_m}{s_m} \right)$ with equality if and only if X_m, Y_m are i.i.d.
 196 normal. Now by the induction hypothesis, $h_{s_m}(Y_m) \geq \sum_{i=1}^{m-1} \frac{\lambda_i}{1 - \lambda_m} h_{r_i}(X_i) + \frac{n}{2} s_m' \left(\frac{\log s_m}{s_m} - \sum_{i=1}^{m-1} \frac{\log r_i}{r_i} \right)$
 197 with equality if and only if the X_i ($i = 1, 2, \dots, m - 1$)—and hence Y_m —are i.i.d. normal. The result at
 198 order m follows by combining the two inequalities since $(1 - \lambda_m) s_m' = r'$. \square

199 5. Recent Versions of the Rényi EPI

200 **Definition 5** (Rényi Entropy Power [5]). *The Rényi entropy power of order r is defined by*

$$N_r(X) = e^{2h_r(X)/n}. \quad (71)$$

201 Up to a multiplicative constant, $N_r(X)$ is the (average) power of a white normal variable having
 202 the same Rényi entropy as X —hence the name “entropy power”. In fact, if $X^* \sim \mathcal{N}(0, \sigma^2)$ has the
 203 same Rényi entropy $h_r(X^*) = h_r(X)$, then by Lemma 4,

$$\sigma^2 = \frac{e^{2h_r(X)/n}}{2\pi r^{r'/r}}. \quad (72)$$

204 The Rényi entropy power enjoys the same scaling property as for the usual power: By Lemma 3, for
 205 any $a \in \mathbb{R}$,

$$N_r(aX) = a^2 N_r(X). \quad (73)$$

206 For independent X_1, X_2, \dots, X_m , Rényi entropy power inequalities take the either the form [5,6]

$$N_r\left(\sum_{i=1}^m X_i\right) \geq c \sum_{i=1}^m N_r(X_i) \quad (74)$$

207 for some positive constant c , or the form [7–9]

$$N_r^\alpha\left(\sum_{i=1}^m X_i\right) \geq \sum_{i=1}^m N_r^\alpha(X_i) \quad (75)$$

208 for some positive exponent α . The constants c and α may depend on the order r , the number m of
 209 variables and the dimension n . What is desired is:

- 210 • a maximum possible value of c in (74) since the inequality is automatically satisfied for all
 211 positive constants $c' < c$.
- 212 • a minimum possible value of α in (75) since the inequality is automatically satisfied for all positive
 213 exponents $\alpha' > \alpha$; in fact, since (75) is homogeneous by scaling the variables $X_i \mapsto aX_i$ as in (73),
 214 one may suppose without loss of generality that the r.h.s. of (75) is = 1; then $N_p(X_i) < 1$ for all i
 215 and $\sum_{i=1}^m N_r^{\alpha'}(X_i) \leq \sum_{i=1}^m N_r^\alpha(X_i)$.

216 The following useful characterization, which generalizes [8, Lemma 2.1], makes the link between the
 217 various versions (70), (74), (75) of the Rényi entropy power inequality.

218 **Lemma 8.** For independent X_1, X_2, \dots, X_m , the Rényi EPI in the general form

$$N_r^\alpha \left(\sum_{i=1}^m X_i \right) \geq c \sum_{i=1}^m N_r^\alpha (X_i) \quad (76)$$

219 for some constant $c > 0$ and exponent $\alpha > 0$ is equivalent to the following inequality

$$h_r \left(\sum_{i=1}^m \sqrt{\lambda_i} X_i \right) - \sum_{i=1}^m \lambda_i h_r (X_i) \geq \frac{n}{2} \left(\frac{\log c}{\alpha} + \left(\frac{1}{\alpha} - 1 \right) H(\lambda) \right) \quad (77)$$

220 for any positive $\lambda_1, \lambda_2, \dots, \lambda_m$ such that $\sum_{i=1}^m \lambda_i = 1$, where $H(\lambda) > 0$ denotes the discrete entropy

$$H(\lambda) = \sum_{i=1}^m \lambda_i \log \frac{1}{\lambda_i} > 0. \quad (78)$$

221 **Proof.** Suppose (76) holds. Then

$$h_r \left(\sum_{i=1}^m \sqrt{\lambda_i} X_i \right) = \frac{n}{2\alpha} \log N_r^\alpha \left(\sum_{i=1}^m \sqrt{\lambda_i} X_i \right) \quad (79)$$

$$\geq \frac{n}{2\alpha} \log \sum_{i=1}^m N_r^\alpha (\sqrt{\lambda_i} X_i) + \frac{n}{2\alpha} \log c \quad (80)$$

$$= \frac{n}{2\alpha} \log \sum_{i=1}^m \lambda_i^\alpha N_r^\alpha (X_i) + \frac{n}{2\alpha} \log c \quad (81)$$

$$\geq \frac{n}{2\alpha} \sum_{i=1}^m \lambda_i \log (\lambda_i^{\alpha-1} N_r^\alpha (X_i)) + \frac{n}{2\alpha} \log c \quad (82)$$

$$= \sum_{i=1}^m \lambda_i h_r (X_i) + \frac{n(\alpha-1)}{2\alpha} \sum_{i=1}^m \lambda_i \log \lambda_i + \frac{n}{2\alpha} \log c \quad (83)$$

222 where the scaling property (73) is used in (81) and the concavity of the logarithm is used in (82).

223 Conversely, suppose that (77) is satisfied for all $\lambda_i > 0$ such that $\sum_{i=1}^m \lambda_i = 1$. Set $\lambda_i =$

224 $N_r^\alpha (X_i) / \sum_{i=1}^m N_r^\alpha (X_i)$. Then

$$N_r^\alpha \left(\sum_{i=1}^m X_i \right) = \exp \frac{2\alpha}{n} h_r \left(\sum_{i=1}^m \sqrt{\lambda_i} \frac{X_i}{\sqrt{\lambda_i}} \right) \quad (84)$$

$$\geq \exp \frac{2\alpha}{n} \sum_{i=1}^m \lambda_i h_r \left(\frac{X_i}{\sqrt{\lambda_i}} \right) \cdot c \cdot \exp(1-\alpha) \sum_{i=1}^m \lambda_i \log \frac{1}{\lambda_i} \quad (85)$$

$$= c \prod_{i=1}^m \left(N_r^\alpha \left(\frac{X_i}{\sqrt{\lambda_i}} \right) \lambda_i^{\alpha-1} \right)^{\lambda_i} \quad (86)$$

$$= c \prod_{i=1}^m \left(N_r^\alpha (X_i) \lambda_i^{-1} \right)^{\lambda_i} \quad (87)$$

$$= c \left(\sum_{i=1}^m N_r^\alpha (X_i) \right)^{\sum_{i=1}^m \lambda_i} \quad (88)$$

$$= c \sum_{i=1}^m N_r^\alpha (X_i). \quad \square \quad (89)$$

225 5.1. Rényi Entropy Power Inequalities for Orders > 1

226 From Lemma 8 and Corollary 2 it is easy to recover known Rényi EPIs and obtain new ones for
 227 orders $r > 1$. In fact, if $r > 1$ then $r' > 0$ and all r'_i are positive and $> r'$. Therefore all r_i are $< r$ and by
 228 monotonicity (Lemma 5),

$$h_{r_i}(X_i) \geq h_r(X_i) \quad (i = 1, 2, \dots, m). \quad (90)$$

229 Plugging this into (70) one obtains

$$h_r\left(\sum_{i=1}^m \sqrt{\lambda_i} X_i\right) - \sum_{i=1}^m \lambda_i h_r(X_i) \geq \frac{n}{2} r' \left(\frac{\log r}{r} - \sum_{i=1}^m \frac{\log r_i}{r_i} \right) \quad (91)$$

230 where $\lambda_i = r'/r'_i$ for $i = 1, 2, \dots, m$. For future reference define¹

$$A(\lambda) = |r'| \left(\frac{\log r}{r} - \sum_{i=1}^m \frac{\log r_i}{r_i} \right) \quad (92)$$

$$= |r'| \left(\frac{\log r}{r} + \sum_{i=1}^m \left(1 - \frac{\lambda_i}{r'}\right) \log\left(1 - \frac{\lambda_i}{r'}\right) \right) \quad (93)$$

231 This function is strictly convex in $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_m)$ because $x \mapsto (1 - x/r') \log(1 - x/r')$ is strictly
 232 convex. Note that $A(\lambda)$ vanishes in the limiting cases where λ tends to one of the vectors $(1, 0, \dots, 0)$,
 233 $(0, 1, 0, \dots, 0), \dots, (0, 0, \dots, 0, 1)$ and since every λ is a convex combination of these vectors and $A(\lambda)$
 234 is strictly convex, one has $A(\lambda) < 0$.

235 **Theorem 4** (Ram and Sason [6]). *The Rényi EPI (74) holds for $r > 1$ and $c = r^{r'/r} (1 - \frac{1}{mr'})^{mr'-1}$.*

236 **Proof.** By Lemma 8 for $\alpha = 1$ we only need to check that the r.h.s. of (91) is greater than $\frac{n}{2} \log c$ for any
 237 choice of the λ_i 's, that is, for any choice of exponents r_i such that $\sum_{i=1}^m \frac{1}{r_i} = \frac{1}{r}$. Thus (74) will hold for
 238 $\log c = \min_{\lambda} A(\lambda)$. Now by the log-sum inequality [4, Theorem 2.7.1],

$$\sum_{i=1}^m \frac{1}{r_i} \log \frac{1}{r_i} \geq \left(\sum_{i=1}^m \frac{1}{r_i} \right) \log \frac{\sum_{i=1}^m \frac{1}{r_i}}{m} = (m - 1/r') \log \frac{m - 1/r'}{m} \quad (94)$$

239 with equality if and only if all r_i are equal, that is, the λ_i are equal to $1/m$. Thus $\min_{\lambda} A(\lambda) =$
 240 $r' \left(\frac{\log r}{r} + (m - 1/r') \log \frac{m - 1/r'}{m} \right)$ which yields $c = r^{r'/r} (1 - \frac{1}{mr'})^{mr'-1}$. \square

241 An alternate proof is to argue that $A(\lambda)$ is convex and symmetrical in $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_m)$ and is,
 242 therefore, minimized when all λ_i are equal.

243 **Remark 5.** *The above constant c is certainly not optimal since equality in (90) holds if and only if the X_i are*
 244 *uniformly distributed (Lemma 5) while equality in (70) holds if and only if the X_i are identically normally*
 245 *distributed (Corollary 2). Ram and Sason [6] tightened (74) further using optimization techniques, resulting in*
 246 *a constant that depends on the relative values of the entropy powers themselves.*

247 **Remark 6.** *It can be noted that $\log c = r' \frac{\log r}{r} + (mr' - 1) \log(1 - \frac{1}{mr'}) < 0$ decreases as m increases; in*
 248 *fact $\frac{\partial \log c}{\partial m} = r' \log(1 - \frac{1}{mr'}) + \frac{mr'}{r'^2 m^2} < r'(-\frac{1}{mr'}) + \frac{1}{m} = 0$. Thus, a universal constant independent of m is*
 249 *obtained by taking*

$$c = \inf_m r^{r'/r} (1 - \frac{1}{mr'})^{mr'-1} = r^{r'/r} \lim_{m \rightarrow \infty} (1 - \frac{1}{mr'})^{mr'-1} = \frac{r^{r'/r}}{e} \quad (95)$$

250 *This was the constant established by Bobkov and Chistyakov [5].*

251 **Theorem 5.** *The Rényi EPI (75) holds for $r > 1$ and $\alpha = \left(1 + \frac{r' \frac{\log r}{r} + (mr' - 1) \log(1 - \frac{1}{mr'})}{\log m}\right)^{-1}$.*

¹ The absolute value $|r'|$ is needed in the next subsection where r' will be negative.

Proof. By Lemma 8 for $c = 1$ we only need to check that the r.h.s. of (91) is greater than $\frac{n}{2}(1/\alpha - 1)H(\lambda)$ for any choice of the λ_i 's, that is, for any choice of exponents r_i such that $\sum_{i=1}^m \frac{1}{r_i} = \frac{1}{r}$. Thus (75) will hold for $\frac{1}{\alpha} - 1 = \min_{\lambda} \frac{A(\lambda)}{H(\lambda)}$. By the proof of the preceding theorem, the numerator is minimized when all λ_i are equal and this also maximizes the entropy $= \log m$ in the denominator. However one cannot conclude yet since the minimum in the numerator is negative.

Suppose that the minimum is *not* attained when all λ_i are equal. Since $A(\lambda)/H(\lambda)$ is symmetric in λ , if the minimum is attained in λ , then it is also attained at any permutation λ' of λ . Without loss of generality we may assume that $\lambda_1 \neq \lambda_2$ and take $\lambda' = (\lambda_2, \lambda_1, \lambda_3, \dots, \lambda_m)$. Now for every $\mu \in [0, 1]$ consider

$$\varphi(\mu) = \frac{A(\mu\lambda + (1-\mu)\lambda')}{H(\mu\lambda + (1-\mu)\lambda')} \quad (96)$$

so that $\min_{\lambda} \frac{A(\lambda)}{H(\lambda)} = \varphi(0) = \varphi(1)$. Differentiating one finds

$$\varphi'(1) = \frac{B(\lambda)}{H(\lambda)} - \frac{A(\lambda)C(\lambda)}{H(\lambda)^2} \quad (97)$$

where $B(\lambda) = \frac{\partial A(\mu\lambda + (1-\mu)\lambda')}{\partial \mu} \Big|_{\mu=1} = \sum_{i=1}^m (\lambda'_i - \lambda_i) \log(e(1 - \lambda_i/r')) = (\lambda_2 - \lambda_1) \log \frac{1-\lambda_1/r'}{1-\lambda_2/r'} > 0$ and $C(\lambda) = \frac{\partial H(\mu\lambda + (1-\mu)\lambda')}{\partial \mu} \Big|_{\mu=1} = \sum_{i=1}^m (\lambda'_i - \lambda_i) \log(e\lambda_i) = (\lambda_2 - \lambda_1) \log \frac{\lambda_1}{\lambda_2} < 0$. Therefore, $\varphi'(1) > 0$ and similarly one finds $\varphi'(0) < 0$. It follows that there exists $0 < \mu < 1$ such that $\varphi(\mu) < \varphi(0) = \varphi(1)$, a contradiction.

Therefore, the minimum of $A(\lambda)/H(\lambda)$ is attained when all λ_i are equal to $1/m$. This gives $\frac{1}{\alpha} - 1 = \left(r' \frac{\log r}{r} + (mr' - 1) \log \left(1 - \frac{1}{mr'} \right) \right) / \log m$. \square

Remark 7. The case $m = 2$ yields $\alpha = \frac{r-1}{(r+1) \log_2(r+1) - r \log_2 r - 2}$ which was found by Li [8] who remarked that this value of α is strictly smaller (better) than the value $\alpha = \frac{r+1}{2}$ obtained by Bobkov and Marsiglietti [7].

For $m > 2$ the exponent of Theorem 5 is even smaller (better). In fact, by Remark 6, it is immediate to see that $\left(r' \frac{\log r}{r} + (mr' - 1) \log \left(1 - \frac{1}{mr'} \right) \right) / \log m$ is negative and decreases toward 0 as m increases. Therefore, the exponent α decreases (is improved) as m increases.

In all cases $\alpha > 1$ and tends to 1 as $m \rightarrow +\infty$. However, using the same method it is easy to obtain Rényi EPIs with exponent values $\alpha < 1$. This is given by the following Theorem.

Theorem 6. The Rényi EPI (76) holds for $r > 1$, $0 < \alpha < 1$ with $c = \left(m r^{r'/r} \left(1 - \frac{1}{mr'} \right)^{mr'-1} \right)^{\alpha} / m$.

Proof. By Lemma 8 we only need to check that the r.h.s. of (91) is greater than $\frac{n}{2}((\log c)/\alpha + (1/\alpha - 1)H(\lambda))$, that is, $A(\lambda) \geq (\log c)/\alpha + (1/\alpha - 1)H(\lambda)$ for any choice of the λ_i 's, that is, for any choice of exponents r_i such that $\sum_{i=1}^m \frac{1}{r_i} = \frac{1}{r}$. Thus for a given $0 < \alpha < 1$, (76) will hold for $\log c = \min_{\lambda} \alpha A(\lambda) - (1 - \alpha)H(\lambda)$. From the preceding proofs the minimum is attained when all λ_i are equal. This gives $\log c = \alpha \left(r' \frac{\log r}{r} + (mr' - 1) \log \left(1 - \frac{1}{mr'} \right) \right) - (1 - \alpha) \log m$. \square

5.2. Rényi Entropy Power Inequalities for Orders < 1 and Log-Concave Densities

If $r < 1$ then $r' < 0$ and all r'_i are negative and $< r'$. Therefore all r_i are $> r$ and by monotonicity (Lemma 5), the opposite inequality of (90) holds and the method of the preceding subsection fails. For log-concave densities, however, (90) can be replaced by a similar inequality in the right direction.

Definition 6 (Log-Concave Density). A density f is log-concave if $\log f$ is concave in its support, i.e., for all $0 < \mu < 1$,

$$f(x)^\mu f(y)^{1-\mu} \leq f(\mu x + (1-\mu)y). \quad (98)$$

287 **Lemma 9.** If X has a log-concave density, then $h_p(pX) - ph_p(X) = n \log p + (1-p)h_p(X)$ is concave in p .

288 As noted below in Remark 8, this is essentially a result obtained by Fradelizi, Madiman and
289 Wang [18]. The following alternate proof uses the transport properties seen in Section 4.

290 **Proof.** Define $r = \lambda p + (1-\lambda)q$ where $0 < \lambda < 1$. By Lemma 1 there exists two diffeomorphisms
291 T, U such that one can write $pX_p = T(X^*)$ and $qX_q = U(X^*)$. Then X^* has density

$$\frac{1}{p^n} f_p\left(\frac{T(x^*)}{p}\right) |T'(x^*)| = \frac{1}{q^n} f_q\left(\frac{U(x^*)}{q}\right) |U'(x^*)| = \frac{1}{p^\lambda q^{(1-\lambda)n}} f_p\left(\frac{T(x^*)}{p}\right)^\lambda f_q\left(\frac{U(x^*)}{q}\right)^{1-\lambda} |T'(x^*)|^\lambda |U'(x^*)|^{1-\lambda} \quad (99)$$

292 Now, by log-concavity (98) with $\mu = \lambda p/r$,

$$f_p\left(\frac{T(x^*)}{p}\right)^\lambda f_q\left(\frac{U(x^*)}{q}\right)^{1-\lambda} = \frac{1}{\|f\|_p^{\lambda p} \|f\|_q^{(1-\lambda)q}} f\left(\frac{T(x^*)}{p}\right)^{\lambda p} f\left(\frac{U(x^*)}{q}\right)^{(1-\lambda)q} \quad (100)$$

$$\leq \frac{1}{\|f\|_p^{\lambda p} \|f\|_q^{(1-\lambda)q}} f\left(\frac{\lambda T(x^*) + (1-\lambda)U(x^*)}{r}\right)^r \quad (101)$$

$$= \frac{\|f\|_r^r}{\|f\|_p^{\lambda p} \|f\|_q^{(1-\lambda)q}} f_r\left(\frac{\lambda T(x^*) + (1-\lambda)U(x^*)}{r}\right) \quad (102)$$

293 Using the arithmetic-geometric mean inequality (58) and integrating the density (99) over $x^* \in \mathbb{R}^n$,
294 one obtains

$$(p^\lambda q^{1-\lambda})^n \|f\|_p^{\lambda p} \|f\|_q^{(1-\lambda)q} \leq r^n \|f\|_r^r. \quad (103)$$

295 Taking the logarithm yields the announced concavity. \square

296 **Remark 8.** Since $n \log p + (1-p)h_p(X)$ is concave, one has $\frac{\partial^2}{\partial p^2} (n \log p + (1-p)h_p(X)) \leq 0$.
297 Differentiating (16) with respect to p it is easy to show that this is equivalent to the following "varentropy bound"
298 $\text{Var} \log f(X_p) \leq n/p^2$, that is, $\text{Var} \log f_p(X_p) \leq n$ which was obtained in [18].

299 Since $n \log p + (1-p)h_p(X)$ is concave and vanishes for $p = 1$, the slopes $\frac{n \log p + (1-p)h_p(X) - 0}{p-1}$ are
300 nonincreasing in p . In other words $h_p(X) + n \frac{\log p}{1-p}$ is nondecreasing. Therefore:

301 **Corollary 3** (Marsiglietti and Melbourne [9]). If $p < q$ then for any X with log-concave density, $h_p(X) +$
302 $n \frac{\log p}{1-p} \leq h_q(X) + n \frac{\log q}{1-q}$.

303 We can now use Lemma 8 and Corollary 2 to obtain Rényi EPIs for orders $r < 1$. Since all r_i are
304 $> r$, by Corollary 3,

$$h_{r_i}(X) + n \frac{\log r_i}{1-r_i} \geq h_r(X) + n \frac{\log r}{1-r} \quad (i = 1, 2, \dots, m). \quad (104)$$

305 Plugging this into (70) one obtains

$$h_r\left(\sum_{i=1}^m \sqrt{\lambda_i} X_i\right) - \sum_{i=1}^m \lambda_i h_r(X_i) \geq n \left(\frac{\log r}{1-r} - \sum_{i=1}^m \lambda_i \frac{\log r_i}{1-r_i}\right) + \frac{n}{2} r' \left(\frac{\log r}{r} - \sum_{i=1}^m \frac{\log r_i}{r_i}\right) \quad (105)$$

$$= \frac{n}{2} r' \left(\sum_{i=1}^m \frac{\log r_i}{r_i} - \frac{\log r}{r}\right) \quad (106)$$

306 where we have used that $\lambda_i = r'/r'_i$ for $i = 1, 2, \dots, m$. Notice that the r.h.s. of (106) for $r < 1$ ($r' < 0$) is
 307 the opposite of that of (91) for $r > 1$ ($r' > 0$). Note, however, that since r' is now negative, the r.h.s.
 308 is exactly $= \frac{n}{2}A(\lambda)$ which is still convex and negative. For this reason, the proofs of the following
 309 theorems for $r < 1$ are such repeats of the theorems obtained previously for $r > 1$.

310 **Theorem 7.** *The Rényi EPI (74) for log-concave densities holds for $c = r^{-r'/r}(1 - \frac{1}{mr'})^{1-mr'}$ and $r < 1$.*

311 **Proof.** Identical to that of Theorem 4 except for the change $|r'| = -r'$ in the expression of $A(\lambda)$. \square

312 **Theorem 8.** *The Rényi EPI (75) for log-concave densities holds for $\alpha = \left(1 - \frac{r' \log r + (mr' - 1) \log(1 - \frac{1}{mr'})}{\log m}\right)^{-1}$
 313 and $r < 1$.*

314 **Proof.** Identical to that of Theorem 5 except for the change $|r'| = -r'$ in the expression of $A(\lambda)$. \square

315 **Remark 9.** *The case $m = 2$ yields $\alpha = \frac{1-r}{(r+1) \log_2(r+1) - r \log_2 r - 2r}$ which was found by Marsiglietti and
 316 Melbourne [9]. Again the exponent of the Theorem is improved (strictly smaller) for $m > 2$. In fact α decreases
 317 and tends to one as $m \rightarrow +\infty$.*

318 **Theorem 9.** *The Rényi EPI (76) for log-concave densities holds for $c = \left(mr^{-r'/r}(1 - \frac{1}{mr'})^{1-mr'}\right)^\alpha / m$ where
 319 $r < 1$ and $0 < \alpha < 1$.*

320 **Proof.** Identical to that of Theorem 6 except for the change $|r'| = -r'$ in the expression of $A(\lambda)$. \square

321 6. Conclusion

322 This article provides a comprehensive framework to derive all known Rényi entropy power
 323 inequalities (with much shorter proofs), and prove new ones. The framework is based on a transport
 324 argument from normal densities and a change of variable by rotation. Only basic properties of Rényi
 325 entropies are used in the proofs.

326 In particular, the α -modification of the EPI is generalized to more than two independent variables
 327 for Rényi entropy orders > 1 as well as for orders < 1 . Also, the Rényi EPI with multiplicative
 328 constant c is extended to Rényi entropy orders < 1 , and a more general formulation with both
 329 exponent α and constant c is obtained for all orders.

330 As a perspective, the methods developed in this paper can perhaps be generalized to obtain
 331 reverse Rényi entropy power inequalities (see e.g., the discussion in [8]).

332

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