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Not peer-reviewed version

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Posted Date: 11 October 2025

doi: 10.20944/preprints202510.0187.v1

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Article

Contra-HyperSoft Set and Contra-SuperHyperSoft Set

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Abstract

A Soft Set is a parameterized family of subsets of a universe, where each parameter selects elements relevant under that condition. A ContraSoft Set is a parameterized soft set in which each parameter's values are linked with a contradiction degree, and a threshold mechanism is applied to retain only those values that do not exceed a specified level of contradiction with respect to a chosen reference. In this paper, we explore two new concepts, namely the *Contra-HyperSoft Set* and the *Contra-SuperHyperSoft Set*, extending the framework of contradiction-aware modeling.

Keywords: soft set; contradiction; ContraSoft set; Contra-HyperSoft set; Contra-SuperHyperSoft set

1. Preliminaries

We collect the basic terminology and notation used in what follows. The definitions in this paper are assumed to be finite.

1.1. Soft Set

A Soft Set is a parameterized family of subsets selecting universe elements relevant to each parameter, supporting flexible decision modeling [1–3]. The definitions of the Soft Set are provided below.

Definition 1 (Soft Set). [1] Let U be a universal set and E a set of parameters. A soft set over U is defined as an ordered pair (F, E) , where F is a mapping from E to the power set $\mathcal{P}(U)$:

$$F : E \rightarrow \mathcal{P}(U).$$

For each parameter $e \in E$, $F(e) \subseteq U$ represents the set of e -approximate elements in U , with (F, E) forming a parameterized family of subsets of U .

Example 1 (Soft Set — Hotel Filtering with Parameterized Conditions). *Universe and parameters.* Let the universe of candidate hotels be

$$U = \{h_1, h_2, h_3, h_4, h_5, h_6, h_7, h_8\}.$$

Let the parameter set be

$$E = \{\text{near_station, free_breakfast, onsen, under}\yen12000, \text{twin_room}\}.$$

Define a soft set (F, E) with $F : E \rightarrow \mathcal{P}(U)$ by

$$\begin{aligned} F(\text{near_station}) &= \{h_1, h_2, h_5, h_7\}, \\ F(\text{free_breakfast}) &= \{h_2, h_4, h_5, h_6\}, \\ F(\text{onsen}) &= \{h_3, h_5, h_8\}, \\ F(\text{under¥12000}) &= \{h_1, h_4, h_5, h_7, h_8\}, \\ F(\text{twin_room}) &= \{h_2, h_4, h_6, h_7\}. \end{aligned}$$

Concrete queries and explicit computations.

$$\begin{aligned} \text{(i) near_station \& under¥12000: } & F(\text{near_station}) \cap F(\text{under¥12000}) = \{h_1, h_2, h_5, h_7\} \cap \{h_1, h_4, h_5, h_7, h_8\} \\ & = \{h_1, h_5, h_7\}, \quad |\cdot| = 3. \\ \text{(ii) (free_breakfast} \cup \text{onsen) \& twin_room: } & (F(\text{free_breakfast}) \cup F(\text{onsen})) \cap F(\text{twin_room}) \\ & = (\{h_2, h_4, h_5, h_6\} \cup \{h_3, h_5, h_8\}) \cap \{h_2, h_4, h_6, h_7\} \\ & = \{h_2, h_4, h_6\}, \quad |\cdot| = 3. \\ \text{(iii) onsen \& twin_room: } & F(\text{onsen}) \cap F(\text{twin_room}) = \{h_3, h_5, h_8\} \cap \{h_2, h_4, h_6, h_7\} = \emptyset. \end{aligned}$$

These results illustrate how a soft set supports multi-criterion filtering by standard set operations with exact outputs.

1.2. ContraSoft Set

A ContraSoft Set is a parameterized soft set where each parameter's values are associated with a contradiction degree, and thresholding is used to aggregate only those values that are not too contradictory with respect to a chosen reference. This allows soft-set modeling to filter or weight information based on contradiction, rather than uncertainty.

Definition 2 (Contradiction on attribute values). *Let V be a nonempty finite set of attribute values. A contradiction function on V is a map*

$$c : V \times V \longrightarrow [0, 1]$$

such that

$$c(v, v) = 0 \quad (\text{reflexivity}), \quad c(v, w) = c(w, v) \quad (\text{symmetry}).$$

The quantity $c(v, w)$ measures the degree of contradiction between v and w (larger means more contradictory).

Example 2 (Contradiction on attribute values — temperature preference). *Let $V = \{\text{cold}, \text{mild}, \text{hot}\}$. Define the symmetric contradiction $c : V \times V \rightarrow [0, 1]$ (with $c(v, v) = 0$) by*

	cold	mild	hot
cold	0	0.4	0.9
mild	0.4	0	0.5
hot	0.9	0.5	0

so, e.g., $c(\text{cold}, \text{hot}) = 0.9$ expresses a strong contradiction, while $c(\text{cold}, \text{mild}) = 0.4$ is moderate.

Definition 3 (ContraSoft structure). *Let U be a nonempty universe and E a nonempty set of parameters. For each $e \in E$ fix:*

- a nonempty finite value set V_e ;
- a contradiction function $c_e : V_e \times V_e \rightarrow [0, 1]$ (Definition 2);
- a designated reference value $v_e^* \in V_e$.

Write $V := \sqcup_{e \in E} (\{e\} \times V_e)$ for the disjoint union of all parameter–value pairs.

Example 3 (ContraSoft structure — hotels by noise and price). Let the universe be $U = \{h_1, h_2, h_3, h_4\}$ and parameters $E = \{\text{noise}, \text{price}\}$. For each $e \in E$ fix a finite value-set V_e , a contradiction $c_e : V_e \times V_e \rightarrow [0, 1]$, and a reference value $v_e^* \in V_e$:

$V_{\text{noise}} = \{\text{quiet}, \text{moderate}, \text{loud}\}, \quad c_{\text{noise}} =$		quiet	moderate	loud	$, \quad v_{\text{noise}}^* = \text{quiet}.$
	quiet	0	0.3	0.8	
	moderate	0.3	0	0.4	
	loud	0.8	0.4	0	
$V_{\text{price}} = \{\text{cheap}, \text{mid}, \text{expensive}\}, \quad c_{\text{price}} =$		cheap	mid	expensive	$, \quad v_{\text{price}}^* = \text{mid}.$
	cheap	0	0.2	0.7	
	mid	0.2	0	0.3	
	expensive	0.7	0.3	0	

These choices realize Definition (ContraSoft structure) by specifying value domains, their contradiction degrees, and per-parameter references.

Definition 4 (ContraSoft Set). Let U be a finite universe of objects and E a finite set of parameters. A ContraSoft Set is a quadruple

$$\text{CS} := (U, E, F, c),$$

where

- $F : E \rightarrow \mathcal{P}(U)$ is the (crisp) soft mapping; $F(e) \subseteq U$ is the set of objects accepted (or classified as positive) under parameter e ;
- $c : E \times E \rightarrow [0, 1]$ is a contradiction degree on parameters, symmetric and reflexive on the diagonal:

$$c(e, e) = 0, \quad c(e, f) = c(f, e) \quad (\forall e, f \in E).$$

For $x \in U$ and $e \in E$, the atomic lemma “ x is accepted by e ” is represented by

$$A(x, e) : \quad x \in F(e),$$

with truth value **T** if $x \in F(e)$ and **F** otherwise.

Example 4 (ContraSoft Set — Noise-Aware Hotel Selection with Contradiction Thresholding). **Universe, parameters, and soft mapping.** Let the same universe U be as above. Consider parameters

$$E = \{\text{quiet}, \text{nightlife}, \text{coworking}, \text{scenic}\}.$$

Define $F : E \rightarrow \mathcal{P}(U)$ by

$$\begin{aligned} F(\text{quiet}) &= \{h_1, h_3, h_5, h_8\}, & F(\text{nightlife}) &= \{h_2, h_4, h_6\}, \\ F(\text{coworking}) &= \{h_2, h_5, h_6, h_7\}, & F(\text{scenic}) &= \{h_3, h_5, h_7, h_8\}. \end{aligned}$$

Contradiction degrees on parameters. Let $c : E \times E \rightarrow [0, 1]$ be symmetric with

$c(\cdot, \cdot)$	quiet	nightlife	coworking	scenic
quiet	0	0.9	0.4	0.2
nightlife	0.9	0	0.3	0.5
coworking	0.4	0.3	0	0.4
scenic	0.2	0.5	0.4	0

(diagonal 0, larger values mean more contradictory).

Reference and thresholded aggregation. Fix the reference parameter $e^* = \text{quiet}$ and threshold $\tau = 0.4$. Define the accepted envelope

$$S^{(\tau)}(e^*) := \bigcup_{e \in E: c(e, e^*) \leq \tau} F(e).$$

Eligible parameters are those within the contradiction radius:

$$c(\text{quiet}, \text{quiet}) = 0, \quad c(\text{coworking}, \text{quiet}) = 0.4, \quad c(\text{scenic}, \text{quiet}) = 0.2, \quad c(\text{nightlife}, \text{quiet}) = 0.9 > \tau.$$

Hence

$$\begin{aligned} S^{(\tau)}(\text{quiet}) &= F(\text{quiet}) \cup F(\text{coworking}) \cup F(\text{scenic}) \\ &= \{h_1, h_3, h_5, h_8\} \cup \{h_2, h_5, h_6, h_7\} \cup \{h_3, h_5, h_7, h_8\} \\ &= \{h_1, h_2, h_3, h_5, h_6, h_7, h_8\}, \quad |S^{(\tau)}| = 7. \end{aligned}$$

Tighter threshold for comparison. With $\tau' = 0.2$, only quiet and scenic are admitted:

$$S^{(\tau')}(\text{quiet}) = F(\text{quiet}) \cup F(\text{scenic}) = \{h_1, h_3, h_5, h_7, h_8\}, \quad |S^{(\tau')}| = 5.$$

Thus $S^{(\tau)}$ is monotone in τ , and the contradiction metric controls how widely we aggregate across potentially conflicting parameters.

1.3. HyperSoft Set and SuperHyperSoft Set

HyperSoft Set maps each multi-attribute tuple from a Cartesian product to a subset of the universe consistent with those values [4–8]. SuperHyperSoft Set maps tuples of subsets from power-set domains to universe subsets, generalizing HyperSoft; singletons in each coordinate recover HyperSoft [9,10].

Definition 5 (HyperSoft Set). [4] Let U be a finite universe and let $\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_m$ be m attribute value domains. Consider the Cartesian product

$$\mathcal{C} = \mathcal{A}_1 \times \mathcal{A}_2 \times \dots \times \mathcal{A}_m,$$

so that each parameter $\gamma = (\gamma_1, \gamma_2, \dots, \gamma_m) \in \mathcal{C}$ chooses a single value $\gamma_i \in \mathcal{A}_i$ for every attribute. A HyperSoft Set over U is a pair (G, \mathcal{C}) where

$$G : \mathcal{C} \longrightarrow \mathcal{P}(U)$$

assigns to each multi-attribute parameter γ a subset $G(\gamma) \subseteq U$. Equivalently,

$$(G, \mathcal{C}) = \{ (\gamma, G(\gamma)) : \gamma \in \mathcal{C} \}.$$

Example 5 (HyperSoft Set — Multi-Attribute Restaurant Finder). **Universe and attributes.** Let the universe of candidate restaurants be

$$U = \{r_1, r_2, r_3, r_4, r_5, r_6, r_7, r_8, r_9\}.$$

Let the attribute domains be

$$\mathcal{A}_1 = \{\text{jpn}, \text{ita}, \text{ind}\}, \quad \mathcal{A}_2 = \{\text{low}, \text{mid}, \text{high}\}, \quad \mathcal{A}_3 = \{\text{omn}, \text{veg}, \text{vgn}\}.$$

The parameter space is the Cartesian product $\mathcal{C} = \mathcal{A}_1 \times \mathcal{A}_2 \times \mathcal{A}_3$. A HyperSoft Set is a mapping $G : \mathcal{C} \rightarrow \mathcal{P}(U)$ that assigns a subset of restaurants to each single-valued tuple (cuisine, price, diet).

Specification (nonempty images).

$$\begin{aligned} G(\text{jpn}, \text{mid}, \text{veg}) &= \{r_2, r_5\}, & G(\text{jpn}, \text{low}, \text{omn}) &= \{r_1, r_3\}, \\ G(\text{ita}, \text{mid}, \text{veg}) &= \{r_4\}, & G(\text{ind}, \text{low}, \text{vgn}) &= \{r_6, r_8\}, \\ G(\text{ita}, \text{high}, \text{omn}) &= \{r_7, r_9\}, \end{aligned}$$

and $G(\gamma) = \emptyset$ for all other $\gamma \in \mathcal{C}$.

Concrete queries with exact set calculations.

- (i) Exactly (jpn, mid, veg): $G(\text{jpn}, \text{mid}, \text{veg}) = \{r_2, r_5\}, |\cdot| = 2$.
- (ii) Union of two precise asks: $G(\text{jpn}, \text{low}, \text{omn}) \cup G(\text{ita}, \text{high}, \text{omn})$
 $= \{r_1, r_3\} \cup \{r_7, r_9\} = \{r_1, r_3, r_7, r_9\}, |\cdot| = 4$.
- (iii) Disjointness of incompatible tuples: $G(\text{jpn}, \text{low}, \text{omn}) \cap G(\text{ind}, \text{low}, \text{vgn}) = \{r_1, r_3\} \cap \{r_6, r_8\} = \emptyset$.

The HyperSoft Set captures single-value choices per attribute; each tuple pinpoints a crisp slice of U .

Definition 6 (SuperHyperSoft Set). [9,11] Let U be a finite universe. Let a_1, a_2, \dots, a_n be distinct attributes with finite, pairwise disjoint value-sets A_1, A_2, \dots, A_n (i.e., $A_i \cap A_j = \emptyset$ for $i \neq j$). Write $\mathcal{P}(A_i)$ for the power set of A_i and form

$$\mathcal{C} = \mathcal{P}(A_1) \times \mathcal{P}(A_2) \times \dots \times \mathcal{P}(A_n).$$

A SuperHyperSoft Set over U is a pair (F, \mathcal{C}) with

$$F : \mathcal{C} \longrightarrow \mathcal{P}(U),$$

so that for each $\gamma = (\alpha_1, \alpha_2, \dots, \alpha_n) \in \mathcal{C}$ (where $\alpha_i \subseteq A_i$) we have a subset $F(\gamma) \subseteq U$. Formally,

$$(F, \mathcal{C}) = \{(\gamma, F(\gamma)) : \gamma \in \mathcal{C}, F(\gamma) \subseteq U\}.$$

Example 6 (SuperHyperSoft Set — Flexible Restaurant Finder with Set-Valued Coordinates). *Universe and attributes.* Use the same U and attribute value-sets $A_1 = \{\text{jpn}, \text{ita}, \text{ind}\}$, $A_2 = \{\text{low}, \text{mid}, \text{high}\}$, $A_3 = \{\text{omn}, \text{veg}, \text{vgn}\}$. In the SuperHyperSoft setting, the parameter space is

$$\mathcal{C} = \mathcal{P}(A_1) \times \mathcal{P}(A_2) \times \mathcal{P}(A_3),$$

so each coordinate is a subset of admissible values (a flexible filter).

Mapping (nonempty images). Define $F : \mathcal{C} \rightarrow \mathcal{P}(U)$ by

$$\begin{aligned} F(\{\text{jpn}, \text{ita}\}, \{\text{low}, \text{mid}\}, \{\text{veg}\}) &= \{r_2, r_4, r_5\}, \\ F(\{\text{ind}\}, \{\text{low}, \text{mid}\}, \{\text{vgn}\}) &= \{r_6, r_8\}, \\ F(\{\text{jpn}, \text{ind}\}, \{\text{low}\}, \{\text{omn}, \text{veg}\}) &= \{r_1, r_3, r_6\}, \\ F(\{\text{ita}\}, \{\text{high}\}, \{\text{omn}\}) &= \{r_7, r_9\}, \end{aligned}$$

and $F(\alpha) = \emptyset$ otherwise.

Reading the parameters. For example, $\alpha = (\{\text{jpn}, \text{ita}\}, \{\text{low}, \text{mid}\}, \{\text{veg}\})$ means: cuisine is Japanese or Italian, price is low or mid, diet is vegetarian. Then $F(\alpha) = \{r_2, r_4, r_5\}$ is the recommended subset.

Coherence with HyperSoft via singletons. If we restrict to singletons in each coordinate, SuperHyperSoft reduces to HyperSoft. Concretely,

$$F(\{\text{jpn}\}, \{\text{mid}\}, \{\text{veg}\}) = \{r_2, r_5\} = G(\text{jpn}, \text{mid}, \text{veg}),$$

so the singleton tuple reproduces the HyperSoft slice exactly. Moreover,

$$F(\{\text{jpn}\}, \{\text{mid}\}, \{\text{veg}\}) \subseteq F(\{\text{jpn}, \text{ita}\}, \{\text{low}, \text{mid}\}, \{\text{veg}\}) = \{r_2, r_4, r_5\},$$

exhibiting the intended flexible expansion when coordinates are broadened from single values to sets of values.

Cardinality checks.

$$|F(\{\text{jpn}, \text{ita}\}, \{\text{low}, \text{mid}\}, \{\text{veg}\})| = 3, \quad |F(\{\text{ind}\}, \{\text{low}, \text{mid}\}, \{\text{vgn}\})| = 2.$$

Thus SuperHyperSoft enables compact specification of multi-value preferences per attribute and directly returns the filtered subset of U .

2. Main Results

In this section, we present and analyze the principal outcomes of our study.

2.1. Contra-HyperSoft Set

Contra-HyperSoft Set augments HyperSoft with a tuple-wise contradiction metric, reference selector, and threshold, uniting parameter slices within the admissible radius.

Definition 7 (Coordinatewise contradiction). Let $\mathcal{A}_1, \dots, \mathcal{A}_m$ be nonempty finite sets. For each $i \in \{1, \dots, m\}$ a contradiction function is a map

$$c_i : \mathcal{A}_i \times \mathcal{A}_i \longrightarrow [0, 1] \quad \text{with} \quad c_i(a, a) = 0, \quad c_i(a, b) = c_i(b, a).$$

When needed for exact reductions, we assume the zero-separation property $c_i(a, b) = 0 \Rightarrow a = b$.

Definition 8 (Tuple-level contradiction). Let $\mathcal{C} := \mathcal{A}_1 \times \dots \times \mathcal{A}_m$ and write $\gamma = (\gamma_1, \dots, \gamma_m)$, $\delta = (\delta_1, \dots, \delta_m) \in \mathcal{C}$. Define the aggregated contradiction by

$$\Delta(\gamma, \delta) := \max_{1 \leq i \leq m} c_i(\gamma_i, \delta_i) \in [0, 1].$$

Then $\Delta(\gamma, \delta) = \Delta(\delta, \gamma)$ and $\Delta(\gamma, \gamma) = 0$. If each c_i is zero-separating, then $\Delta(\gamma, \delta) = 0 \iff \gamma = \delta$.

Definition 9 (Reference selector). A reference selector is a map $\rho : \mathcal{C} \rightarrow \mathcal{C}$. Two canonical choices are

$$(\text{self-centered}) \quad \rho(\gamma) = \gamma, \quad (\text{fixed-reference}) \quad \rho(\gamma) \equiv r \text{ for a fixed } r \in \mathcal{C}.$$

Definition 10 (Contra-HyperSoft Set (CHS)). Let U be a finite universe and let $G : \mathcal{C} \rightarrow \mathcal{P}(U)$ be a HyperSoft mapping. Fix contradiction kernels $\{c_i\}_{i=1}^m$, a reference selector ρ , and a threshold $\tau \in [0, 1]$. The associated Contra-HyperSoft Set is the tuple

$$\text{CHS} := (U, \{\mathcal{A}_i, c_i\}_{i=1}^m, G, \rho, \tau),$$

together with the filtered mapping

$$G_\rho^{(\tau)} : \mathcal{C} \longrightarrow \mathcal{P}(U), \quad G_\rho^{(\tau)}(\gamma) := \bigcup_{\delta \in \mathcal{C} : \Delta(\delta, \rho(\gamma)) \leq \tau} G(\delta).$$

Example 7 (Contra-HyperSoft Set — Candidate Shortlisting under Conflicting Signals (self-centered selector)). **Universe and attributes.** Let the candidate pool be $U = \{c_1, c_2, c_3, c_4, c_5, c_6, c_7, c_8\}$. Consider three single-valued attribute domains:

$$\mathcal{A}_1 = \{\text{junior}, \text{mid}, \text{senior}\}, \quad \mathcal{A}_2 = \{\text{onsite}, \text{hybrid}, \text{remote}\}, \quad \mathcal{A}_3 = \{\text{backend}, \text{frontend}, \text{data}\}.$$

The parameter space is $\mathcal{C} = \mathcal{A}_1 \times \mathcal{A}_2 \times \mathcal{A}_3$.

Coordinatewise contradictions. All c_i are symmetric with 0 on the diagonal.

		junior	mid	senior	$c_1 =$
junior	junior	0	0.4	0.9	
	mid	0.4	0	0.4	
	senior	0.9	0.4	0	

		onsite	hybrid	remote	$c_2 =$
onsite	onsite	0	0.3	0.8	
	hybrid	0.3	0	0.3	
	remote	0.8	0.3	0	

		backend	frontend	data	$c_3 =$
backend	backend	0	0.5	0.4	
	frontend	0.5	0	0.6	
	data	0.4	0.6	0	

Aggregate tuple-contradiction: $\Delta(\gamma, \delta) := \max\{c_1(\gamma_1, \delta_1), c_2(\gamma_2, \delta_2), c_3(\gamma_3, \delta_3)\}$.

HyperSoft mapping $G : \mathcal{C} \rightarrow \mathcal{P}(U)$ (nonempty images).

$$\begin{aligned} G(\text{junior}, \text{remote}, \text{frontend}) &= \{c_1, c_3\}, & G(\text{mid}, \text{hybrid}, \text{backend}) &= \{c_2, c_5\}, \\ G(\text{senior}, \text{onsite}, \text{data}) &= \{c_4\}, & G(\text{mid}, \text{remote}, \text{data}) &= \{c_6\}, \\ G(\text{senior}, \text{hybrid}, \text{backend}) &= \{c_7, c_8\}. \end{aligned}$$

CHS filter. Choose the self-centered selector $\rho(\gamma) = \gamma$ and threshold $\tau = 0.4$. Let $\gamma^* = (\text{mid}, \text{hybrid}, \text{backend})$. Compute $\Delta(\cdot, \gamma^*)$ on the above tuples:

δ	$\Delta(\delta, \gamma^*)$
(mid, hybrid, backend)	$\max(0, 0, 0) = 0 \ (\leq \tau)$
(senior, hybrid, backend)	$\max(0.4, 0, 0) = 0.4 \ (\leq \tau)$
(mid, remote, data)	$\max(0, 0.3, 0.4) = 0.4 \ (\leq \tau)$
(senior, onsite, data)	$\max(0.4, 0.3, 0.4) = 0.4 \ (\leq \tau)$
(junior, remote, frontend)	$\max(0.4, 0.3, 0.5) = 0.5 \ (> \tau)$

Hence

$$G_\rho^{(\tau)}(\gamma^*) = \{c_2, c_5\} \cup \{c_7, c_8\} \cup \{c_6\} \cup \{c_4\} = \{c_2, c_4, c_5, c_6, c_7, c_8\}, \quad |\cdot| = 6.$$

With a tighter threshold $\tau' = 0.3$, only the base slice survives: $G_\rho^{(\tau')}(\gamma^*) = \{c_2, c_5\}$, illustrating monotonicity in τ .

Example 8 (Contra-HyperSoft Set — Travel Package Selection (fixed reference)). **Universe and attributes.** Let $U = \{\text{pkg}_1, \dots, \text{pkg}_{10}\}$ be travel packages. Attributes:

$$\mathcal{A}_1 = \{\text{winter}, \text{spring}, \text{summer}, \text{autumn}\}, \quad \mathcal{A}_2 = \{\text{ski}, \text{beach}, \text{culture}\}, \quad \mathcal{A}_3 = \{\text{solo}, \text{couple}, \text{family}\}.$$

Contradiction matrices (symmetric, 0 on diagonal).

		winter	spring	summer	autumn	$c_1 =$
winter	winter	0	0.3	0.8	0.5	
	spring	0.3	0	0.3	0.5	
	summer	0.8	0.3	0	0.3	
	autumn	0.5	0.5	0.3	0	

		ski	beach	culture	
$c_2 =$	ski	0	0.9	0.4	,
	beach	0.9	0	0.4	
	culture	0.4	0.4	0	
		solo	couple	family	
$c_3 =$	solo	0	0.2	0.6	.
	couple	0.2	0	0.3	
	family	0.6	0.3	0	

Aggregate $\Delta(\gamma, \delta) := \max\{c_1(\gamma_1, \delta_1), c_2(\gamma_2, \delta_2), c_3(\gamma_3, \delta_3)\}$.

HyperSoft mapping (nonempty images).

$$G(\text{winter}, \text{ski}, \text{family}) = \{\text{pkg}_1, \text{pkg}_2\}, \quad G(\text{summer}, \text{beach}, \text{couple}) = \{\text{pkg}_3, \text{pkg}_4\},$$

$$G(\text{spring}, \text{culture}, \text{solo}) = \{\text{pkg}_5\}, \quad G(\text{autumn}, \text{culture}, \text{family}) = \{\text{pkg}_6, \text{pkg}_7\},$$

$$G(\text{summer}, \text{culture}, \text{family}) = \{\text{pkg}_8\}, \quad G(\text{winter}, \text{beach}, \text{solo}) = \{\text{pkg}_9\}.$$

CHS filter (fixed reference). Choose the fixed reference $r = (\text{summer}, \text{beach}, \text{family})$ and threshold $\tau = 0.5$. Evaluate $\Delta(\cdot, r)$:

γ	$\Delta(\gamma, r)$
(winter, ski, family)	$\max(0.8, 0.9, 0) = 0.9 (> \tau)$
(summer, beach, couple)	$\max(0, 0, 0.3) = 0.3 (\leq \tau)$
(spring, culture, solo)	$\max(0.3, 0.4, 0.6) = 0.6 (> \tau)$
(autumn, culture, family)	$\max(0.3, 0.4, 0) = 0.4 (\leq \tau)$
(summer, culture, family)	$\max(0, 0.4, 0) = 0.4 (\leq \tau)$
(winter, beach, solo)	$\max(0.8, 0, 0.6) = 0.8 (> \tau)$

Thus the accepted tuples are the 2nd, 4th, and 5th. The CHS envelope at r is

$$G_{\rho \equiv r}^{(\tau)}(r) = \{\text{pkg}_3, \text{pkg}_4\} \cup \{\text{pkg}_6, \text{pkg}_7\} \cup \{\text{pkg}_8\} = \{\text{pkg}_3, \text{pkg}_4, \text{pkg}_6, \text{pkg}_7, \text{pkg}_8\},$$

with cardinality 5. If we tighten to $\tau' = 0.3$, only the 2nd tuple remains, so $G_{\rho \equiv r}^{(\tau')}(r) = \{\text{pkg}_3, \text{pkg}_4\}$, demonstrating the control afforded by the contradiction threshold.

Proposition 1 (Basic properties). For fixed $(U, \{\mathcal{A}_i, c_i\}, G, \rho)$ the family $\{G_{\rho}^{(\tau)}\}_{\tau \in [0,1]}$ is monotone in τ : if $0 \leq \tau_1 \leq \tau_2 \leq 1$ then $G_{\rho}^{(\tau_1)}(\gamma) \subseteq G_{\rho}^{(\tau_2)}(\gamma)$ for all $\gamma \in \mathcal{C}$. Moreover $G_{\rho}^{(1)}(\gamma) = \bigcup_{\delta \in \mathcal{C}} G(\delta)$ for all γ .

Proof. If $\tau_1 \leq \tau_2$ then $\{\delta : \Delta(\delta, \rho(\gamma)) \leq \tau_1\} \subseteq \{\delta : \Delta(\delta, \rho(\gamma)) \leq \tau_2\}$, hence the unions are nested. For $\tau = 1$ the constraint is vacuous since $\Delta \in [0, 1]$. \square

Theorem 1 (CHS generalizes the HyperSoft Set). Assume each c_i is zero-separating and take the self-centered selector $\rho(\gamma) = \gamma$. Then for $\tau = 0$ one has

$$G_{\rho}^{(0)}(\gamma) = G(\gamma) \quad (\forall \gamma \in \mathcal{C}).$$

Proof. By definition, $G_{\rho}^{(0)}(\gamma) = \bigcup_{\delta: \Delta(\delta, \gamma) \leq 0} G(\delta)$. Since $\Delta(\delta, \gamma) \geq 0$ always, the inequality forces $\Delta(\delta, \gamma) = 0$. Zero-separation gives $\delta = \gamma$, thus the union is $G(\gamma)$. \square

Definition 11 (Neighborhood-based ContraSoft on a single attribute). Let V be a finite set with contradiction $c : V \times V \rightarrow [0, 1]$, and let $F : V \rightarrow \mathcal{P}(U)$. For $\tau \in [0, 1]$, the neighborhood-based ContraSoft transform is

$$F^{(\tau)}(v) := \bigcup_{w \in V: c(w,v) \leq \tau} F(w) \quad (v \in V).$$

Fixing $v^* \in V$ yields the fixed-reference variant $F^{(\tau; v^*)}(v) := \bigcup_{w: c(w, v^*) \leq \tau} F(w)$.

Example 9 (Neighborhood-based ContraSoft — Destination Selection by Climate Preference). **Universe and attribute.** Let the universe of candidate destinations be

$$U = \{u_1 = \text{Reykjavik}, u_2 = \text{Zurich}, u_3 = \text{Lisbon}, u_4 = \text{Dubai}, u_5 = \text{Helsinki}, u_6 = \text{Vancouver}\}.$$

Consider a single attribute “preferred climate” with value set

$$V = \{\text{cold}, \text{mild}, \text{warm}, \text{hot}\}.$$

Contradiction on V . Let $c : V \times V \rightarrow [0, 1]$ be symmetric with $c(v, v) = 0$:

	cold	mild	warm	hot
cold	0	0.3	0.6	0.9
mild	0.3	0	0.3	0.7
warm	0.6	0.3	0	0.3
hot	0.9	0.7	0.3	0

Baseline soft mapping. Define $F : V \rightarrow \mathcal{P}(U)$ by

$$F(\text{cold}) = \{u_1, u_5\}, \quad F(\text{mild}) = \{u_2, u_6\}, \quad F(\text{warm}) = \{u_3\}, \quad F(\text{hot}) = \{u_4\}.$$

Neighborhood-based ContraSoft transform. For threshold $\tau \in [0, 1]$ and center $v \in V$,

$$F^{(\tau)}(v) = \bigcup_{w \in V: c(w,v) \leq \tau} F(w).$$

Case 1 (moderate neighborhood). Let $v = \text{mild}$ and $\tau = 0.35$. Eligible neighbors satisfy $c(w, \text{mild}) \leq 0.35$:

$$c(\text{cold}, \text{mild}) = 0.3 (\checkmark), \quad c(\text{mild}, \text{mild}) = 0 (\checkmark), \quad c(\text{warm}, \text{mild}) = 0.3 (\checkmark), \quad c(\text{hot}, \text{mild}) = 0.7 (\times).$$

Therefore

$$\begin{aligned} F^{(0.35)}(\text{mild}) &= F(\text{cold}) \cup F(\text{mild}) \cup F(\text{warm}) \\ &= \{u_1, u_5\} \cup \{u_2, u_6\} \cup \{u_3\} = \{u_1, u_2, u_3, u_5, u_6\}, \quad |\cdot| = 5. \end{aligned}$$

Case 2 (tight neighborhood). Let $\tau = 0.20$. Only $w = \text{mild}$ satisfies $c(w, \text{mild}) \leq 0.20$, hence

$$F^{(0.20)}(\text{mild}) = F(\text{mild}) = \{u_2, u_6\}, \quad |\cdot| = 2.$$

These computations show how increasing τ expands the accepted neighborhood in V and unions the corresponding destination sets in U .

Theorem 2 (CHS generalizes ContraSoft). Suppose $m = 1$, so $\mathcal{C} = \mathcal{A}_1 =: V$, and let $c_1 = c$. Identify $G : V \rightarrow \mathcal{P}(U)$ with F . Then:

(a) With the self-centered selector $\rho(v) = v$, one has

$$G_\rho^{(\tau)}(v) = \bigcup_{w: c(w,v) \leq \tau} G(w) = F^{(\tau)}(v), \quad \forall v \in V.$$

(b) With the fixed-reference selector $\rho(v) \equiv v^*$, one has

$$G_\rho^{(\tau)}(v) = \bigcup_{w: c(w,v^*) \leq \tau} G(w) = F^{(\tau;v^*)}(v), \quad \forall v \in V.$$

Hence, for $m = 1$ the CHS construction recovers both standard ContraSoft variants.

Proof. When $m = 1$, $\Delta(w, \rho(v)) = c(w, \rho(v))$. Substituting $\rho(v) = v$ gives (a); substituting $\rho(v) \equiv v^*$ gives (b). The set-theoretic unions agree by definition in both cases. \square

2.2. Contra-SuperHyperSoft Set

Contra-SuperHyperSoft Set extends to set-valued coordinates, using lifted subset contradictions and aggregate radius; selector-threshold filtering unions nearby SuperHyperSoft slices effectively.

Definition 12 (Base and lifted contradictions). Let A_1, \dots, A_m be nonempty finite sets of attribute values and let $c_i : A_i \times A_i \rightarrow [0, 1]$ be contradiction functions (symmetric and reflexive: $c_i(a, a) = 0 = c_i(a, a)$, $c_i(a, b) = c_i(b, a)$). Assume the zero-separation property $c_i(a, b) = 0 \Rightarrow a = b$. For subsets $S, T \subseteq A_i$ define the lifted contradiction

$$\hat{c}_i(S, T) := \begin{cases} \max \left\{ \max_{a \in S} \min_{b \in T} c_i(a, b), \max_{b \in T} \min_{a \in S} c_i(a, b) \right\}, & S, T \neq \emptyset, \\ 0, & S = T = \emptyset, \\ 1, & \text{otherwise.} \end{cases}$$

(With finiteness, the max / min are attained.)

Lemma 1 (Symmetry, reflexivity, and zero-separation on subsets). For each i and $S, T \subseteq A_i$:

- (a) $\hat{c}_i(S, T) = \hat{c}_i(T, S)$ and $\hat{c}_i(S, S) = 0$.
- (b) If c_i is zero-separating, then $\hat{c}_i(S, T) = 0$ implies $S = T$.

Proof. (a) Symmetry follows by exchanging the two max terms; reflexivity is immediate. (b) If $\hat{c}_i(S, T) = 0$ with $S, T \neq \emptyset$, then $\max_{a \in S} \min_{b \in T} c_i(a, b) = 0$ and $\max_{b \in T} \min_{a \in S} c_i(a, b) = 0$. Thus, for each $a \in S$ there is $b \in T$ with $c_i(a, b) = 0$, hence $a = b$ by zero-separation, so $S \subseteq T$. The second equality gives $T \subseteq S$. The cases with empties are by definition. \square

Definition 13 (Product parameter space and aggregate contradiction). Let $\mathcal{C} := \mathcal{P}(A_1) \times \dots \times \mathcal{P}(A_m)$ and write $\alpha = (\alpha_1, \dots, \alpha_m)$, $\beta = (\beta_1, \dots, \beta_m) \in \mathcal{C}$. Define the tuple-level contradiction by

$$\Delta(\alpha, \beta) := \max_{1 \leq i \leq m} \hat{c}_i(\alpha_i, \beta_i) \in [0, 1].$$

Then $\Delta(\alpha, \beta) = \Delta(\beta, \alpha)$ and $\Delta(\alpha, \alpha) = 0$; if each c_i is zero-separating, then by Lemma 1 we have $\Delta(\alpha, \beta) = 0 \iff \alpha = \beta$.

Definition 14 (Reference selector). A reference selector is any map $\rho : \mathcal{C} \rightarrow \mathcal{C}$. Two common choices are the self-centered selector $\rho(\alpha) = \alpha$ and a fixed-reference selector $\rho(\alpha) \equiv r$ for a fixed $r \in \mathcal{C}$.

Definition 15 (Contra-SuperHyperSoft Set (CSHS)). Let U be a finite universe and let $F : \mathcal{C} \rightarrow \mathcal{P}(U)$ be a SuperHyperSoft mapping. Fix contradiction kernels $\{c_i\}$, their lifts $\{\hat{c}_i\}$, an aggregate Δ , a selector ρ , and a threshold $\tau \in [0, 1]$. The associated Contra-SuperHyperSoft Set is the tuple

$$\text{CSHS} := (U, \{A_i, c_i\}_{i=1}^m, F, \rho, \tau),$$

together with the filtered mapping

$$F_\rho^{(\tau)} : \mathcal{C} \longrightarrow \mathcal{P}(U), \quad F_\rho^{(\tau)}(\alpha) := \bigcup_{\beta \in \mathcal{C} : \Delta(\beta, \rho(\alpha)) \leq \tau} F(\beta).$$

Example 10 (CSHS in E-commerce Fraud Review (self-centered selector)). **Setup.** Let the universe of orders be $U = \{o_1, o_2, o_3, o_4, o_5, o_6, o_7, o_8\}$. Take two attribute domains:

$$A_1 = \{\text{card}, \text{crypto}\}, \quad A_2 = \{\text{verified}, \text{partial}, \text{missing}\}.$$

Base contradictions $c_1, c_2 : [\cdot] \rightarrow [0, 1]$ (symmetric, 0 on the diagonal):

$$c_1 = \begin{array}{c|cc} & \text{card} & \text{crypto} \\ \hline \text{card} & 0 & 0.8 \\ \text{crypto} & 0.8 & 0 \end{array}, \quad c_2 = \begin{array}{c|ccc} & \text{verified} & \text{partial} & \text{missing} \\ \hline \text{verified} & 0 & 0.3 & 0.9 \\ \text{partial} & 0.3 & 0 & 0.6 \\ \text{missing} & 0.9 & 0.6 & 0 \end{array}.$$

Lift \hat{c}_i to subsets by Definition (lifted contradiction) and aggregate

$$\Delta((S_1, S_2), (T_1, T_2)) := \max\{\hat{c}_1(S_1, T_1), \hat{c}_2(S_2, T_2)\}.$$

The SuperHyperSoft mapping $F : \mathcal{P}(A_1) \times \mathcal{P}(A_2) \rightarrow \mathcal{P}(U)$ is specified by

$$\begin{aligned} F(\{\text{crypto}\}, \{\text{missing}\}) &= \{o_1, o_2\}, & F(\{\text{crypto}\}, \{\text{partial}\}) &= \{o_3\}, \\ F(\{\text{card}\}, \{\text{missing}\}) &= \{o_4\}, & F(\{\text{card}\}, \{\text{partial}\}) &= \{o_5, o_6\}, \\ F(\{\text{card}, \text{crypto}\}, \{\text{partial}, \text{missing}\}) &= \{o_7\}, & \text{all other pairs map to } \emptyset. \end{aligned}$$

Choose the self-centered selector $\rho(\alpha) = \alpha$ and threshold $\tau = 0.6$.

Filtering at $\alpha_0 = (\{\text{crypto}\}, \{\text{missing}\})$.

To avoid overfull lines, we list the computations in an aligned display:

$$\begin{aligned} \Delta((\{\text{crypto}\}, \{\text{missing}\}), \alpha_0) &= \max(0, 0) = 0 \leq \tau, \\ \Delta((\{\text{crypto}\}, \{\text{partial}\}), \alpha_0) &= \max(0, c_2(\text{partial}, \text{missing})) = \max(0, 0.6) = 0.6 \leq \tau, \\ \Delta((\{\text{card}\}, \{\text{missing}\}), \alpha_0) &= \max(0.8, 0) = 0.8 > \tau, \\ \Delta((\{\text{card}\}, \{\text{partial}\}), \alpha_0) &= \max(0.8, 0.6) = 0.8 > \tau, \\ \Delta((\{\text{card}, \text{crypto}\}, \{\text{partial}, \text{missing}\}), \alpha_0) &= \max(\hat{c}_1(\{\text{card}, \text{crypto}\}, \{\text{crypto}\}), \hat{c}_2(\{\text{partial}, \text{missing}\}, \{\text{missing}\})) \\ &= \max(0.8, 0.6) = 0.8 > \tau. \end{aligned}$$

Hence

$$F_\rho^{(\tau)}(\alpha_0) = \{o_1, o_2\} \cup \{o_3\} = \{o_1, o_2, o_3\}, \quad |F_\rho^{(\tau)}(\alpha_0)| = 3.$$

Filtering at $\alpha_1 = (\{\text{card}\}, \{\text{partial}\})$.

β	$\Delta(\beta, \alpha_1)$
$(\{\text{card}\}, \{\text{partial}\})$	$0 \leq \tau$
$(\{\text{card}\}, \{\text{missing}\})$	$\max(0, 0.6) = 0.6 \leq \tau$
$(\{\text{crypto}\}, \{\text{partial}\})$	$\max(0.8, 0) = 0.8 > \tau$
$(\{\text{crypto}\}, \{\text{missing}\})$	$\max(0.8, 0.6) = 0.8 > \tau$
$(\{\text{card}, \text{crypto}\}, \{\text{partial}, \text{missing}\})$	$\max(0.8, 0.6) = 0.8 > \tau$

Therefore

$$F_p^{(\tau)}(\alpha_1) = \{o_5, o_6\} \cup \{o_4\} = \{o_4, o_5, o_6\}, \quad |F_p^{(\tau)}(\alpha_1)| = 3.$$

This illustrates how the CSHS envelope aggregates nearby subset-parameters under the contradiction metric.

Example 11 (CSHS in Cloud Deployment Recommendation (self-centered selector)). *Setup.* Let the universe of candidate nodes be $U = \{n_1, n_2, n_3, n_4, n_5, n_6, n_7\}$. Attributes:

$$A_1 = \{\text{us-east}, \text{us-west}, \text{eu}\}, \quad A_2 = \{\text{cpu}, \text{gpu}, \text{memory}\}.$$

Base contradictions (0 on the diagonal, symmetric):

$c_1 =$		us-east	us-west	eu	$c_2 =$		cpu	gpu	memory
	us-east	0	0.4	0.7		cpu	0	0.5	0.3
	us-west	0.4	0	0.8		gpu	0.5	0	0.7
	eu	0.7	0.8	0		memory	0.3	0.7	0

Lift to subsets by \hat{c}_i and aggregate by $\Delta((S_1, S_2), (T_1, T_2)) = \max\{\hat{c}_1(S_1, T_1), \hat{c}_2(S_2, T_2)\}$.

SuperHyperSoft mapping F (nonempty images shown):

$$\begin{aligned} F(\{\text{us-east}\}, \{\text{cpu}\}) &= \{n_1, n_2\}, & F(\{\text{us-east}\}, \{\text{gpu}\}) &= \{n_3\}, \\ F(\{\text{us-west}\}, \{\text{cpu}\}) &= \{n_4\}, & F(\{\text{eu}\}, \{\text{cpu}\}) &= \{n_5\}, \\ F(\{\text{us-east}, \text{us-west}\}, \{\text{cpu}\}) &= \{n_6\}, \\ F(\{\text{us-east}\}, \{\text{cpu}, \text{gpu}\}) &= \{n_7\}. \end{aligned}$$

Choose the self-centered selector $\rho(\alpha) = \alpha$ and threshold $\tau = 0.5$.

Filtering at $\alpha^* = (\{\text{us-east}\}, \{\text{cpu}\})$. For each β with $F(\beta) \neq \emptyset$, compute $\Delta(\beta, \alpha^*)$:

β	$\Delta(\beta, \alpha^*)$
$(\{\text{us-east}\}, \{\text{cpu}\})$	$\max(0, 0) = 0 \leq \tau$
$(\{\text{us-east}\}, \{\text{gpu}\})$	$\max(0, c_2(\text{gpu}, \text{cpu}) = 0.5) = 0.5 \leq \tau$
$(\{\text{us-west}\}, \{\text{cpu}\})$	$\max(c_1(\text{us-west}, \text{us-east}) = 0.4, 0) = 0.4 \leq \tau$
$(\{\text{eu}\}, \{\text{cpu}\})$	$\max(0.7, 0) = 0.7 > \tau$
$(\{\text{us-east}, \text{us-west}\}, \{\text{cpu}\})$	$\max(\hat{c}_1(\{\text{us-east}, \text{us-west}\}, \{\text{us-east}\}) = 0.4, 0) = 0.4 \leq \tau$
$(\{\text{us-east}\}, \{\text{cpu}, \text{gpu}\})$	$\max(0, \hat{c}_2(\{\text{cpu}, \text{gpu}\}, \{\text{cpu}\}) = 0.5) = 0.5 \leq \tau$

Hence

$$F_p^{(\tau)}(\alpha^*) = \{n_1, n_2\} \cup \{n_3\} \cup \{n_4\} \cup \{n_6\} \cup \{n_7\} = \{n_1, n_2, n_3, n_4, n_6, n_7\},$$

with $|F_p^{(\tau)}(\alpha^*)| = 6$. Nodes requiring a European region are excluded by the contradiction bound.

Proposition 2 (Monotonicity in the threshold). If $0 \leq \tau_1 \leq \tau_2 \leq 1$, then $F_p^{(\tau_1)}(\alpha) \subseteq F_p^{(\tau_2)}(\alpha)$ for all $\alpha \in \mathcal{C}$. Moreover, $F_p^{(1)}(\alpha) = \bigcup_{\beta \in \mathcal{C}} F(\beta)$ for all α .

Proof. The index sets $\{\beta : \Delta(\beta, \rho(\alpha)) \leq \tau\}$ are nested as τ grows, hence so are the unions. For $\tau = 1$ the constraint is vacuous since $\Delta \in [0, 1]$. \square

Theorem 3 (CSHS generalizes SuperHyperSoft). *Assume zero-separation for each c_i and take the self-centered selector $\rho(\alpha) = \alpha$. Then for $\tau = 0$ we have*

$$F_{\rho}^{(0)}(\alpha) = F(\alpha) \quad (\forall \alpha \in \mathcal{C}).$$

Proof. By definition, $F_{\rho}^{(0)}(\alpha) = \bigcup_{\beta: \Delta(\beta, \alpha) \leq 0} F(\beta)$. Since $\Delta \geq 0$, the inequality forces $\Delta(\beta, \alpha) = 0$. By Definition 13 and Lemma 1, this occurs iff $\beta = \alpha$. Hence the union collapses to $F(\alpha)$. \square

Definition 16 (Singleton embedding of HyperSoft into SuperHyperSoft). *Let $\iota : A_1 \times \cdots \times A_m \rightarrow \mathcal{C}$ be*

$$\iota(\gamma_1, \dots, \gamma_m) := (\{\gamma_1\}, \dots, \{\gamma_m\}).$$

We say $F : \mathcal{C} \rightarrow \mathcal{P}(U)$ is singleton-supported for $G : A_1 \times \cdots \times A_m \rightarrow \mathcal{P}(U)$ if $F(\iota(\gamma)) = G(\gamma)$ for all γ and $F(\alpha) = \emptyset$ whenever some α_i is not a singleton.

Lemma 2 (Compatibility of contradictions on singletons). *For any $a, b \in A_i$ we have $\widehat{c}_i(\{a\}, \{b\}) = c_i(a, b)$. Consequently, for $\gamma, \delta \in A_1 \times \cdots \times A_m$,*

$$\Delta(\iota(\gamma), \iota(\delta)) = \max_{1 \leq i \leq m} c_i(\gamma_i, \delta_i).$$

Proof. From Definition 12 with singletons, $\max\{\min c_i(a, b), \min c_i(b, a)\} = c_i(a, b)$ by symmetry, and the product case follows. \square

Theorem 4 (CSHS generalizes Contra-HyperSoft). *Let $G : A_1 \times \cdots \times A_m \rightarrow \mathcal{P}(U)$ be a HyperSoft set and F be a singleton-supported extension (Definition 16). Let ρ_H be a selector on $A_1 \times \cdots \times A_m$ and define ρ_S on \mathcal{C} by $\rho_S(\iota(\gamma)) := \iota(\rho_H(\gamma))$, with arbitrary values elsewhere. Then for every γ and $\tau \in [0, 1]$,*

$$F_{\rho_S}^{(\tau)}(\iota(\gamma)) = \bigcup_{\delta: \max_i c_i(\delta_i, \rho_H(\gamma)_i) \leq \tau} G(\delta) =: G_{\rho_H}^{(\tau)}(\gamma),$$

i.e. the CSHS filtered mapping restricted to singleton parameters coincides with the Contra-HyperSoft mapping for G built from $\{c_i\}$ and the same threshold.

Proof. By Definition 15,

$$F_{\rho_S}^{(\tau)}(\iota(\gamma)) = \bigcup_{\beta: \Delta(\beta, \rho_S(\iota(\gamma))) \leq \tau} F(\beta).$$

Since F is singleton-supported, only β of the form $\iota(\delta)$ contribute. By Lemma 2 and the definition of ρ_S , $\Delta(\iota(\delta), \rho_S(\iota(\gamma))) = \max_i c_i(\delta_i, \rho_H(\gamma)_i)$. Substitute and use $F(\iota(\delta)) = G(\delta)$ to obtain the stated equality. \square

Theorem 5 (CSHS generalizes ContraSoft). *Let $m = 1$ with base set $A_1 =: V$ and contradiction $c := c_1$. Let $F : \mathcal{P}(V) \rightarrow \mathcal{P}(U)$ be singleton-supported for some $F_0 : V \rightarrow \mathcal{P}(U)$ via $F(\{v\}) = F_0(v)$ and $F(S) = \emptyset$ if $|S| \neq 1$. Then, with the self-centered selector, for all $v \in V$ and $\tau \in [0, 1]$,*

$$F_{\text{id}}^{(\tau)}(\{v\}) = \bigcup_{w: c(w, v) \leq \tau} F_0(w),$$

which is exactly the neighborhood-based ContraSoft transform on (V, c) . With a fixed reference $v^ \in V$, the same construction yields the fixed-reference ContraSoft variant.*

Proof. Specialize Theorem 4 to $m = 1$ with $G \equiv F_0$ and observe that $\Delta(\{w\}, \{v\}) = \hat{c}_1(\{w\}, \{v\}) = c(w, v)$. Singleton support reduces the union to singletons, yielding the stated form. \square

3. Conclusion

In this paper, we examined the concepts of the ContraSoft Set, the Contra-HyperSoft Set, and the Contra-SuperHyperSoft Set.

Future work will focus on extending these ideas by integrating them with richer frameworks, including Neutrosophic Sets[12,13], Plithogenic Sets[14–16], Rough Sets[17,18], and TreeSoft Sets[19,20]. Moreover, we anticipate the development of generalized structures that apply these contra-based approaches to Graphs[21,22], HyperGraphs[23–25], SuperHyperGraphs[26–28], and broader HyperStructures[29–31]. Such investigations are expected to open new directions in handling contradiction-aware representations across diverse domains.

Funding: This study did not receive any financial or external support from organizations or individuals.

Institutional Review Board Statement: As this research is entirely theoretical in nature and does not involve human participants or animal subjects, no ethical approval is required.

Data Availability Statement: This research is purely theoretical, involving no data collection or analysis. We encourage future researchers to pursue empirical investigations to further develop and validate the concepts introduced here.

Acknowledgments: We extend our sincere gratitude to everyone who provided insights, inspiration, and assistance throughout this research. We particularly thank our readers for their interest and acknowledge the authors of the cited works for laying the foundation that made our study possible. We also appreciate the support from individuals and institutions that provided the resources and infrastructure needed to produce and share this paper. Finally, we are grateful to all those who supported us in various ways during this project.

Use of Artificial Intelligence: I use generative AI and AI-assisted tools for tasks such as English grammar checking, and I do not employ them in any way that violates ethical standards.

Conflicts of Interest: The authors confirm that there are no conflicts of interest related to the research or its publication.

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