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Article

Contra-HyperSoft Set and Contra-SuperHyperSoft Set

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Abstract

A Soft Set is a parameterized family of subsets of a universe, where each parameter selects elements relevant under that condition. A ContraSoft Set is a parameterized soft set in which each parameter's values are linked with a contradiction degree, and a threshold mechanism is applied to retain only those values that do not exceed a specified level of contradiction with respect to a chosen reference. In this paper, we explore two new concepts, namely the *Contra-HyperSoft Set* and the *Contra-SuperHyperSoft Set*, extending the framework of contradiction-aware modeling.

Keywords: soft set; contradiction; ContraSoft set; Contra-HyperSoft set; Contra-SuperHyperSoft set

1. Preliminaries

We collect the basic terminology and notation used in what follows. The definitions in this paper are assumed to be finite.

1.1. Soft Set

A Soft Set is a parameterized family of subsets selecting universe elements relevant to each parameter, supporting flexible decision modeling [1–3]. The definitions of the Soft Set are provided below.

Definition 1 (Soft Set). [1] Let U be a universal set and E a set of parameters. A soft set over U is defined as an ordered pair (F, E) , where F is a mapping from E to the power set $\mathcal{P}(U)$:

$$F : E \rightarrow \mathcal{P}(U).$$

For each parameter $e \in E$, $F(e) \subseteq U$ represents the set of e -approximate elements in U , with (F, E) forming a parameterized family of subsets of U .

Example 1 (Soft Set — Hotel Filtering with Parameterized Conditions). **Universe and parameters.** Let the universe of candidate hotels be

$$U = \{h_1, h_2, h_3, h_4, h_5, h_6, h_7, h_8\}.$$

Let the parameter set be

$$E = \{\text{near_station}, \text{free_breakfast}, \text{onsen}, \text{under}\text{¥}12000, \text{twin_room}\}.$$

Define a soft set (F, E) with $F : E \rightarrow \mathcal{P}(U)$ by

$$\begin{aligned} F(\text{near_station}) &= \{h_1, h_2, h_5, h_7\}, \\ F(\text{free_breakfast}) &= \{h_2, h_4, h_5, h_6\}, \\ F(\text{onsen}) &= \{h_3, h_5, h_8\}, \\ F(\text{under}¥12000) &= \{h_1, h_4, h_5, h_7, h_8\}, \\ F(\text{twin_room}) &= \{h_2, h_4, h_6, h_7\}. \end{aligned}$$

Concrete queries and explicit computations.

(i) *near_station & under¥12000*: $F(\text{near_station}) \cap F(\text{under}¥12000) = \{h_1, h_2, h_5, h_7\} \cap \{h_1, h_4, h_5, h_7, h_8\} = \{h_1, h_5, h_7\}, |\cdot| = 3.$

(ii) *(free_breakfast ∪ onsen) & twin_room*: $(F(\text{free_breakfast}) \cup F(\text{onsen})) \cap F(\text{twin_room}) = (\{h_2, h_4, h_5, h_6\} \cup \{h_3, h_5, h_8\}) \cap \{h_2, h_4, h_6, h_7\} = \{h_2, h_4, h_6\}, |\cdot| = 3.$

(iii) *onsen & twin_room*: $F(\text{onsen}) \cap F(\text{twin_room}) = \{h_3, h_5, h_8\} \cap \{h_2, h_4, h_6, h_7\} = \emptyset.$

These results illustrate how a soft set supports multi-criterion filtering by standard set operations with exact outputs.

1.2. ContraSoft Set

A ContraSoft Set is a parameterized soft set where each parameter's values are associated with a contradiction degree, and thresholding is used to aggregate only those values that are not too contradictory with respect to a chosen reference. This allows soft-set modeling to filter or weight information based on contradiction, rather than uncertainty.

Definition 2 (Contradiction on attribute values). *Let V be a nonempty finite set of attribute values. A contradiction function on V is a map*

$$c : V \times V \longrightarrow [0, 1]$$

such that

$$c(v, v) = 0 \quad (\text{reflexivity}), \quad c(v, w) = c(w, v) \quad (\text{symmetry}).$$

The quantity $c(v, w)$ measures the degree of contradiction between v and w (larger means more contradictory).

Example 2 (Contradiction on attribute values — temperature preference). *Let $V = \{\text{cold, mild, hot}\}$. Define the symmetric contradiction $c : V \times V \rightarrow [0, 1]$ (with $c(v, v) = 0$) by*

		cold	mild	hot
c	cold	0	0.4	0.9
	mild	0.4	0	0.5
	hot	0.9	0.5	0

so, e.g., $c(\text{cold, hot}) = 0.9$ expresses a strong contradiction, while $c(\text{cold, mild}) = 0.4$ is moderate.

Definition 3 (ContraSoft structure). *Let U be a nonempty universe and E a nonempty set of parameters. For each $e \in E$ fix:*

- a nonempty finite value set V_e ;
- a contradiction function $c_e : V_e \times V_e \rightarrow [0, 1]$ (Definition 2);
- a designated reference value $v_e^* \in V_e$.

Write $V := \bigsqcup_{e \in E} (\{e\} \times V_e)$ for the disjoint union of all parameter–value pairs.

Example 3 (ContraSoft structure — hotels by noise and price). Let the universe be $U = \{h_1, h_2, h_3, h_4\}$ and parameters $E = \{\text{noise, price}\}$. For each $e \in E$ fix a finite value-set V_e , a contradiction $c_e : V_e \times V_e \rightarrow [0, 1]$, and a reference value $v_e^* \in V_e$:

		quiet	moderate	loud	
$V_{\text{noise}} = \{\text{quiet, moderate, loud}\}$, c_{noise}	quiet	0	0.3	0.8	$v_{\text{noise}}^* = \text{quiet}$.
	moderate	0.3	0	0.4	
	loud	0.8	0.4	0	
		cheap	mid	expensive	
$V_{\text{price}} = \{\text{cheap, mid, expensive}\}$, c_{price}	cheap	0	0.2	0.7	$v_{\text{price}}^* = \text{mid}$.
	mid	0.2	0	0.3	
	expensive	0.7	0.3	0	

These choices realize Definition (ContraSoft structure) by specifying value domains, their contradiction degrees, and per-parameter references.

Definition 4 (ContraSoft Set). Let U be a finite universe of objects and E a finite set of parameters. A ContraSoft Set is a quadruple

$$\text{CS} := (U, E, F, c),$$

where

- $F : E \rightarrow \mathcal{P}(U)$ is the (crisp) soft mapping; $F(e) \subseteq U$ is the set of objects accepted (or classified as positive) under parameter e ;
- $c : E \times E \rightarrow [0, 1]$ is a contradiction degree on parameters, symmetric and reflexive on the diagonal:

$$c(e, e) = 0, \quad c(e, f) = c(f, e) \quad (\forall e, f \in E).$$

For $x \in U$ and $e \in E$, the atomic lemma “ x is accepted by e ” is represented by

$$A(x, e) : \quad x \in F(e),$$

with truth value **T** if $x \in F(e)$ and **F** otherwise.

Example 4 (ContraSoft Set — Noise-Aware Hotel Selection with Contradiction Thresholding). **Universe, parameters, and soft mapping.** Let the same universe U be as above. Consider parameters

$$E = \{\text{quiet, nightlife, coworking, scenic}\}.$$

Define $F : E \rightarrow \mathcal{P}(U)$ by

$$\begin{aligned} F(\text{quiet}) &= \{h_1, h_3, h_5, h_8\}, & F(\text{nightlife}) &= \{h_2, h_4, h_6\}, \\ F(\text{coworking}) &= \{h_2, h_5, h_6, h_7\}, & F(\text{scenic}) &= \{h_3, h_5, h_7, h_8\}. \end{aligned}$$

Contradiction degrees on parameters. Let $c : E \times E \rightarrow [0, 1]$ be symmetric with

$c(\cdot, \cdot)$	quiet	nightlife	coworking	scenic
quiet	0	0.9	0.4	0.2
nightlife	0.9	0	0.3	0.5
coworking	0.4	0.3	0	0.4
scenic	0.2	0.5	0.4	0

(diagonal 0, larger values mean more contradictory).

Reference and thresholded aggregation. Fix the reference parameter $e^* = \text{quiet}$ and threshold $\tau = 0.4$. Define the accepted envelope

$$S^{(\tau)}(e^*) := \bigcup_{e \in E: c(e, e^*) \leq \tau} F(e).$$

Eligible parameters are those within the contradiction radius:

$$c(\text{quiet}, \text{quiet}) = 0, \quad c(\text{coworking}, \text{quiet}) = 0.4, \quad c(\text{scenic}, \text{quiet}) = 0.2, \quad c(\text{nightlife}, \text{quiet}) = 0.9 > \tau.$$

Hence

$$\begin{aligned} S^{(\tau)}(\text{quiet}) &= F(\text{quiet}) \cup F(\text{coworking}) \cup F(\text{scenic}) \\ &= \{h_1, h_3, h_5, h_8\} \cup \{h_2, h_5, h_6, h_7\} \cup \{h_3, h_5, h_7, h_8\} \\ &= \{h_1, h_2, h_3, h_5, h_6, h_7, h_8\}, \quad |S^{(\tau)}| = 7. \end{aligned}$$

Tighter threshold for comparison. With $\tau' = 0.2$, only quiet and scenic are admitted:

$$S^{(\tau')}(\text{quiet}) = F(\text{quiet}) \cup F(\text{scenic}) = \{h_1, h_3, h_5, h_7, h_8\}, \quad |S^{(\tau')}| = 5.$$

Thus $S^{(\tau)}$ is monotone in τ , and the contradiction metric controls how widely we aggregate across potentially conflicting parameters.

1.3. HyperSoft Set and SuperHyperSoft Set

HyperSoft Set maps each multi-attribute tuple from a Cartesian product to a subset of the universe consistent with those values [4–8]. SuperHyperSoft Set maps tuples of subsets from power-set domains to universe subsets, generalizing HyperSoft; singletons in each coordinate recover HyperSoft [9,10].

Definition 5 (HyperSoft Set). [4] Let U be a finite universe and let $\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_m$ be m attribute value domains. Consider the Cartesian product

$$\mathcal{C} = \mathcal{A}_1 \times \mathcal{A}_2 \times \dots \times \mathcal{A}_m,$$

so that each parameter $\gamma = (\gamma_1, \gamma_2, \dots, \gamma_m) \in \mathcal{C}$ chooses a single value $\gamma_i \in \mathcal{A}_i$ for every attribute. A HyperSoft Set over U is a pair (G, \mathcal{C}) where

$$G : \mathcal{C} \longrightarrow \mathcal{P}(U)$$

assigns to each multi-attribute parameter γ a subset $G(\gamma) \subseteq U$. Equivalently,

$$(G, \mathcal{C}) = \{(\gamma, G(\gamma)) : \gamma \in \mathcal{C}\}.$$

Example 5 (HyperSoft Set — Multi-Attribute Restaurant Finder). **Universe and attributes.** Let the universe of candidate restaurants be

$$U = \{r_1, r_2, r_3, r_4, r_5, r_6, r_7, r_8, r_9\}.$$

Let the attribute domains be

$$\mathcal{A}_1 = \{\text{jpn, ita, ind}\}, \quad \mathcal{A}_2 = \{\text{low, mid, high}\}, \quad \mathcal{A}_3 = \{\text{omn, veg, vgn}\}.$$

The parameter space is the Cartesian product $\mathcal{C} = \mathcal{A}_1 \times \mathcal{A}_2 \times \mathcal{A}_3$. A HyperSoft Set is a mapping $G : \mathcal{C} \rightarrow \mathcal{P}(U)$ that assigns a subset of restaurants to each single-valued tuple $(\text{cuisine, price, diet})$.

Specification (nonempty images).

$$\begin{aligned} G(\text{jpn, mid, veg}) &= \{r_2, r_5\}, & G(\text{jpn, low, omn}) &= \{r_1, r_3\}, \\ G(\text{ita, mid, veg}) &= \{r_4\}, & G(\text{ind, low, vgn}) &= \{r_6, r_8\}, \\ G(\text{ita, high, omn}) &= \{r_7, r_9\}, \end{aligned}$$

and $G(\gamma) = \emptyset$ for all other $\gamma \in \mathcal{C}$.

Concrete queries with exact set calculations.

- (i) *Exactly* (jpn, mid, veg) : $G(\text{jpn, mid, veg}) = \{r_2, r_5\}$, $|\cdot| = 2$.
- (ii) *Union of two precise asks*: $G(\text{jpn, low, omn}) \cup G(\text{ita, high, omn})$
 $= \{r_1, r_3\} \cup \{r_7, r_9\} = \{r_1, r_3, r_7, r_9\}$, $|\cdot| = 4$.
- (iii) *Disjointness of incompatible tuples*: $G(\text{jpn, low, omn}) \cap G(\text{ind, low, vgn}) = \{r_1, r_3\} \cap \{r_6, r_8\} = \emptyset$.

The HyperSoft Set captures single-value choices per attribute; each tuple pinpoints a crisp slice of U .

Definition 6 (SuperHyperSoft Set). [9,11] Let U be a finite universe. Let a_1, a_2, \dots, a_n be distinct attributes with finite, pairwise disjoint value-sets A_1, A_2, \dots, A_n (i.e., $A_i \cap A_j = \emptyset$ for $i \neq j$). Write $\mathcal{P}(A_i)$ for the power set of A_i and form

$$\mathcal{C} = \mathcal{P}(A_1) \times \mathcal{P}(A_2) \times \dots \times \mathcal{P}(A_n).$$

A SuperHyperSoft Set over U is a pair (F, \mathcal{C}) with

$$F : \mathcal{C} \longrightarrow \mathcal{P}(U),$$

so that for each $\gamma = (\alpha_1, \alpha_2, \dots, \alpha_n) \in \mathcal{C}$ (where $\alpha_i \subseteq A_i$) we have a subset $F(\gamma) \subseteq U$. Formally,

$$(F, \mathcal{C}) = \{(\gamma, F(\gamma)) : \gamma \in \mathcal{C}, F(\gamma) \subseteq U\}.$$

Example 6 (SuperHyperSoft Set — Flexible Restaurant Finder with Set-Valued Coordinates). **Universe and attributes.** Use the same U and attribute value-sets $A_1 = \{\text{jpn, ita, ind}\}$, $A_2 = \{\text{low, mid, high}\}$, $A_3 = \{\text{omn, veg, vgn}\}$. In the SuperHyperSoft setting, the parameter space is

$$\mathcal{C} = \mathcal{P}(A_1) \times \mathcal{P}(A_2) \times \mathcal{P}(A_3),$$

so each coordinate is a subset of admissible values (a flexible filter).

Mapping (nonempty images). Define $F : \mathcal{C} \rightarrow \mathcal{P}(U)$ by

$$\begin{aligned} F(\{\text{jpn, ita}\}, \{\text{low, mid}\}, \{\text{veg}\}) &= \{r_2, r_4, r_5\}, \\ F(\{\text{ind}\}, \{\text{low, mid}\}, \{\text{vgn}\}) &= \{r_6, r_8\}, \\ F(\{\text{jpn, ind}\}, \{\text{low}\}, \{\text{omn, veg}\}) &= \{r_1, r_3, r_6\}, \\ F(\{\text{ita}\}, \{\text{high}\}, \{\text{omn}\}) &= \{r_7, r_9\}, \end{aligned}$$

and $F(\alpha) = \emptyset$ otherwise.

Reading the parameters. For example, $\alpha = (\{\text{jpn, ita}\}, \{\text{low, mid}\}, \{\text{veg}\})$ means: cuisine is Japanese or Italian, price is low or mid, diet is vegetarian. Then $F(\alpha) = \{r_2, r_4, r_5\}$ is the recommended subset.

Coherence with HyperSoft via singletons. If we restrict to singletons in each coordinate, SuperHyperSoft reduces to HyperSoft. Concretely,

$$F(\{\text{jpn}\}, \{\text{mid}\}, \{\text{veg}\}) = \{r_2, r_5\} = G(\text{jpn, mid, veg}),$$

so the singleton tuple reproduces the HyperSoft slice exactly. Moreover,

$$F(\{\text{jpn}\}, \{\text{mid}\}, \{\text{veg}\}) \subseteq F(\{\text{jpn, ita}\}, \{\text{low, mid}\}, \{\text{veg}\}) = \{r_2, r_4, r_5\},$$

exhibiting the intended flexible expansion when coordinates are broadened from single values to sets of values.

Cardinality checks.

$$|F(\{\text{jpn, ita}\}, \{\text{low, mid}\}, \{\text{veg}\})| = 3, \quad |F(\{\text{ind}\}, \{\text{low, mid}\}, \{\text{vgn}\})| = 2.$$

Thus SuperHyperSoft enables compact specification of multi-value preferences per attribute and directly returns the filtered subset of U .

2. Main Results

In this section, we present and analyze the principal outcomes of our study.

2.1. Contra-HyperSoft Set

Contra-HyperSoft Set augments HyperSoft with a tuple-wise contradiction metric, reference selector, and threshold, uniting parameter slices within the admissible radius.

Definition 7 (Coordinatewise contradiction). Let $\mathcal{A}_1, \dots, \mathcal{A}_m$ be nonempty finite sets. For each $i \in \{1, \dots, m\}$ a contradiction function is a map

$$c_i : \mathcal{A}_i \times \mathcal{A}_i \longrightarrow [0, 1] \quad \text{with} \quad c_i(a, a) = 0, \quad c_i(a, b) = c_i(b, a).$$

When needed for exact reductions, we assume the zero-separation property $c_i(a, b) = 0 \Rightarrow a = b$.

Definition 8 (Tuple-level contradiction). Let $\mathcal{C} := \mathcal{A}_1 \times \dots \times \mathcal{A}_m$ and write $\gamma = (\gamma_1, \dots, \gamma_m)$, $\delta = (\delta_1, \dots, \delta_m) \in \mathcal{C}$. Define the aggregated contradiction by

$$\Delta(\gamma, \delta) := \max_{1 \leq i \leq m} c_i(\gamma_i, \delta_i) \in [0, 1].$$

Then $\Delta(\gamma, \delta) = \Delta(\delta, \gamma)$ and $\Delta(\gamma, \gamma) = 0$. If each c_i is zero-separating, then $\Delta(\gamma, \delta) = 0 \iff \gamma = \delta$.

Definition 9 (Reference selector). A reference selector is a map $\rho : \mathcal{C} \rightarrow \mathcal{C}$. Two canonical choices are

$$(\text{self-centered}) \quad \rho(\gamma) = \gamma, \quad (\text{fixed-reference}) \quad \rho(\gamma) \equiv r \text{ for a fixed } r \in \mathcal{C}.$$

Definition 10 (Contra-HyperSoft Set (CHS)). Let U be a finite universe and let $G : \mathcal{C} \rightarrow \mathcal{P}(U)$ be a HyperSoft mapping. Fix contradiction kernels $\{c_i\}_{i=1}^m$, a reference selector ρ , and a threshold $\tau \in [0, 1]$. The associated Contra-HyperSoft Set is the tuple

$$\text{CHS} := (U, \{\mathcal{A}_i, c_i\}_{i=1}^m, G, \rho, \tau),$$

together with the filtered mapping

$$G_\rho^{(\tau)} : \mathcal{C} \longrightarrow \mathcal{P}(U), \quad G_\rho^{(\tau)}(\gamma) := \bigcup_{\delta \in \mathcal{C} : \Delta(\delta, \rho(\gamma)) \leq \tau} G(\delta).$$

Example 7 (Contra-HyperSoft Set — Candidate Shortlisting under Conflicting Signals (self-centered selector)). **Universe and attributes.** Let the candidate pool be $U = \{c_1, c_2, c_3, c_4, c_5, c_6, c_7, c_8\}$. Consider three single-valued attribute domains:

$$\mathcal{A}_1 = \{\text{junior, mid, senior}\}, \quad \mathcal{A}_2 = \{\text{onsite, hybrid, remote}\}, \quad \mathcal{A}_3 = \{\text{backend, frontend, data}\}.$$

The parameter space is $\mathcal{C} = \mathcal{A}_1 \times \mathcal{A}_2 \times \mathcal{A}_3$.

Coordinatewise contradictions. All c_i are symmetric with 0 on the diagonal.

$$c_1 = \begin{array}{c|ccc} & \text{junior} & \text{mid} & \text{senior} \\ \hline \text{junior} & 0 & 0.4 & 0.9 \\ \text{mid} & 0.4 & 0 & 0.4 \\ \text{senior} & 0.9 & 0.4 & 0 \end{array},$$

$$c_2 = \begin{array}{c|ccc} & \text{onsite} & \text{hybrid} & \text{remote} \\ \hline \text{onsite} & 0 & 0.3 & 0.8 \\ \text{hybrid} & 0.3 & 0 & 0.3 \\ \text{remote} & 0.8 & 0.3 & 0 \end{array},$$

$$c_3 = \begin{array}{c|ccc} & \text{backend} & \text{frontend} & \text{data} \\ \hline \text{backend} & 0 & 0.5 & 0.4 \\ \text{frontend} & 0.5 & 0 & 0.6 \\ \text{data} & 0.4 & 0.6 & 0 \end{array}.$$

Aggregate tuple-contradiction: $\Delta(\gamma, \delta) := \max\{c_1(\gamma_1, \delta_1), c_2(\gamma_2, \delta_2), c_3(\gamma_3, \delta_3)\}$.

HyperSoft mapping $G : \mathcal{C} \rightarrow \mathcal{P}(U)$ (nonempty images).

$$G(\text{junior, remote, frontend}) = \{c_1, c_3\}, \quad G(\text{mid, hybrid, backend}) = \{c_2, c_5\},$$

$$G(\text{senior, onsite, data}) = \{c_4\}, \quad G(\text{mid, remote, data}) = \{c_6\},$$

$$G(\text{senior, hybrid, backend}) = \{c_7, c_8\}.$$

CHS filter. Choose the self-centered selector $\rho(\gamma) = \gamma$ and threshold $\tau = 0.4$. Let $\gamma^* = (\text{mid, hybrid, backend})$. Compute $\Delta(\cdot, \gamma^*)$ on the above tuples:

$$\begin{array}{c|c} \delta & \Delta(\delta, \gamma^*) \\ \hline (\text{mid, hybrid, backend}) & \max(0, 0, 0) = 0 \ (\leq \tau) \\ (\text{senior, hybrid, backend}) & \max(0.4, 0, 0) = 0.4 \ (\leq \tau) \\ (\text{mid, remote, data}) & \max(0, 0.3, 0.4) = 0.4 \ (\leq \tau) \\ (\text{senior, onsite, data}) & \max(0.4, 0.3, 0.4) = 0.4 \ (\leq \tau) \\ (\text{junior, remote, frontend}) & \max(0.4, 0.3, 0.5) = 0.5 \ (> \tau) \end{array}$$

Hence

$$G_\rho^{(\tau)}(\gamma^*) = \{c_2, c_5\} \cup \{c_7, c_8\} \cup \{c_6\} \cup \{c_4\} = \{c_2, c_4, c_5, c_6, c_7, c_8\}, \quad |\cdot| = 6.$$

With a tighter threshold $\tau' = 0.3$, only the base slice survives: $G_\rho^{(\tau')}(\gamma^*) = \{c_2, c_5\}$, illustrating monotonicity in τ .

Example 8 (Contra-HyperSoft Set — Travel Package Selection (fixed reference)). **Universe and attributes.** Let $U = \{\text{pkg}_1, \dots, \text{pkg}_{10}\}$ be travel packages. Attributes:

$$\mathcal{A}_1 = \{\text{winter, spring, summer, autumn}\}, \quad \mathcal{A}_2 = \{\text{ski, beach, culture}\}, \quad \mathcal{A}_3 = \{\text{solo, couple, family}\}.$$

Contradiction matrices (symmetric, 0 on diagonal).

$$c_1 = \begin{array}{c|cccc} & \text{winter} & \text{spring} & \text{summer} & \text{autumn} \\ \hline \text{winter} & 0 & 0.3 & 0.8 & 0.5 \\ \text{spring} & 0.3 & 0 & 0.3 & 0.5 \\ \text{summer} & 0.8 & 0.3 & 0 & 0.3 \\ \text{autumn} & 0.5 & 0.5 & 0.3 & 0 \end{array},$$

$$c_2 = \begin{array}{c|ccc} & \text{ski} & \text{beach} & \text{culture} \\ \hline \text{ski} & 0 & 0.9 & 0.4 \\ \text{beach} & 0.9 & 0 & 0.4 \\ \text{culture} & 0.4 & 0.4 & 0 \end{array},$$

$$c_3 = \begin{array}{c|ccc} & \text{solo} & \text{couple} & \text{family} \\ \hline \text{solo} & 0 & 0.2 & 0.6 \\ \text{couple} & 0.2 & 0 & 0.3 \\ \text{family} & 0.6 & 0.3 & 0 \end{array}.$$

Aggregate $\Delta(\gamma, \delta) := \max\{c_1(\gamma_1, \delta_1), c_2(\gamma_2, \delta_2), c_3(\gamma_3, \delta_3)\}$.

HyperSoft mapping (nonempty images).

$$G(\text{winter, ski, family}) = \{\text{pkg}_1, \text{pkg}_2\}, \quad G(\text{summer, beach, couple}) = \{\text{pkg}_3, \text{pkg}_4\},$$

$$G(\text{spring, culture, solo}) = \{\text{pkg}_5\}, \quad G(\text{autumn, culture, family}) = \{\text{pkg}_6, \text{pkg}_7\},$$

$$G(\text{summer, culture, family}) = \{\text{pkg}_8\}, \quad G(\text{winter, beach, solo}) = \{\text{pkg}_9\}.$$

CHS filter (fixed reference). Choose the fixed reference $r = (\text{summer, beach, family})$ and threshold $\tau = 0.5$. Evaluate $\Delta(\cdot, r)$:

γ	$\Delta(\gamma, r)$
(winter, ski, family)	$\max(0.8, 0.9, 0) = 0.9 (> \tau)$
(summer, beach, couple)	$\max(0, 0, 0.3) = 0.3 (\leq \tau)$
(spring, culture, solo)	$\max(0.3, 0.4, 0.6) = 0.6 (> \tau)$
(autumn, culture, family)	$\max(0.3, 0.4, 0) = 0.4 (\leq \tau)$
(summer, culture, family)	$\max(0, 0.4, 0) = 0.4 (\leq \tau)$
(winter, beach, solo)	$\max(0.8, 0, 0.6) = 0.8 (> \tau)$

Thus the accepted tuples are the 2nd, 4th, and 5th. The CHS envelope at r is

$$G_{\rho \equiv r}^{(\tau)}(r) = \{\text{pkg}_3, \text{pkg}_4\} \cup \{\text{pkg}_6, \text{pkg}_7\} \cup \{\text{pkg}_8\} = \{\text{pkg}_3, \text{pkg}_4, \text{pkg}_6, \text{pkg}_7, \text{pkg}_8\},$$

with cardinality 5. If we tighten to $\tau' = 0.3$, only the 2nd tuple remains, so $G_{\rho \equiv r}^{(\tau')}(r) = \{\text{pkg}_3, \text{pkg}_4\}$, demonstrating the control afforded by the contradiction threshold.

Proposition 1 (Basic properties). For fixed $(U, \{\mathcal{A}_i, c_i\}, G, \rho)$ the family $\{G_{\rho}^{(\tau)}\}_{\tau \in [0,1]}$ is monotone in τ : if $0 \leq \tau_1 \leq \tau_2 \leq 1$ then $G_{\rho}^{(\tau_1)}(\gamma) \subseteq G_{\rho}^{(\tau_2)}(\gamma)$ for all $\gamma \in \mathcal{C}$. Moreover $G_{\rho}^{(1)}(\gamma) = \bigcup_{\delta \in \mathcal{C}} G(\delta)$ for all γ .

Proof. If $\tau_1 \leq \tau_2$ then $\{\delta : \Delta(\delta, \rho(\gamma)) \leq \tau_1\} \subseteq \{\delta : \Delta(\delta, \rho(\gamma)) \leq \tau_2\}$, hence the unions are nested. For $\tau = 1$ the constraint is vacuous since $\Delta \in [0, 1]$. \square

Theorem 1 (CHS generalizes the HyperSoft Set). Assume each c_i is zero-separating and take the self-centered selector $\rho(\gamma) = \gamma$. Then for $\tau = 0$ one has

$$G_{\rho}^{(0)}(\gamma) = G(\gamma) \quad (\forall \gamma \in \mathcal{C}).$$

Proof. By definition, $G_{\rho}^{(0)}(\gamma) = \bigcup_{\delta: \Delta(\delta, \gamma) \leq 0} G(\delta)$. Since $\Delta(\delta, \gamma) \geq 0$ always, the inequality forces $\Delta(\delta, \gamma) = 0$. Zero-separation gives $\delta = \gamma$, thus the union is $G(\gamma)$. \square

Definition 11 (Neighborhood-based ContraSoft on a single attribute). *Let V be a finite set with contradiction $c : V \times V \rightarrow [0, 1]$, and let $F : V \rightarrow \mathcal{P}(U)$. For $\tau \in [0, 1]$, the neighborhood-based ContraSoft transform is*

$$F^{(\tau)}(v) := \bigcup_{w \in V: c(w,v) \leq \tau} F(w) \quad (v \in V).$$

Fixing $v^* \in V$ yields the fixed-reference variant $F^{(\tau;v^*)}(v) := \bigcup_{w: c(w,v^*) \leq \tau} F(w)$.

Example 9 (Neighborhood-based ContraSoft — Destination Selection by Climate Preference). *Universe and attribute.* Let the universe of candidate destinations be

$$U = \{u_1 = \text{Reykjavik}, u_2 = \text{Zurich}, u_3 = \text{Lisbon}, u_4 = \text{Dubai}, u_5 = \text{Helsinki}, u_6 = \text{Vancouver}\}.$$

Consider a single attribute “preferred climate” with value set

$$V = \{\text{cold, mild, warm, hot}\}.$$

Contradiction on V . Let $c : V \times V \rightarrow [0, 1]$ be symmetric with $c(v, v) = 0$:

		cold	mild	warm	hot
c =	cold	0	0.3	0.6	0.9
	mild	0.3	0	0.3	0.7
	warm	0.6	0.3	0	0.3
	hot	0.9	0.7	0.3	0

Baseline soft mapping. Define $F : V \rightarrow \mathcal{P}(U)$ by

$$F(\text{cold}) = \{u_1, u_5\}, \quad F(\text{mild}) = \{u_2, u_6\}, \quad F(\text{warm}) = \{u_3\}, \quad F(\text{hot}) = \{u_4\}.$$

Neighborhood-based ContraSoft transform. For threshold $\tau \in [0, 1]$ and center $v \in V$,

$$F^{(\tau)}(v) = \bigcup_{w \in V: c(w,v) \leq \tau} F(w).$$

Case 1 (moderate neighborhood). Let $v = \text{mild}$ and $\tau = 0.35$. Eligible neighbors satisfy $c(w, \text{mild}) \leq 0.35$:

$$c(\text{cold, mild}) = 0.3 (\checkmark), \quad c(\text{mild, mild}) = 0 (\checkmark), \quad c(\text{warm, mild}) = 0.3 (\checkmark), \quad c(\text{hot, mild}) = 0.7 (\times).$$

Therefore

$$\begin{aligned} F^{(0.35)}(\text{mild}) &= F(\text{cold}) \cup F(\text{mild}) \cup F(\text{warm}) \\ &= \{u_1, u_5\} \cup \{u_2, u_6\} \cup \{u_3\} = \{u_1, u_2, u_3, u_5, u_6\}, \quad |\cdot| = 5. \end{aligned}$$

Case 2 (tight neighborhood). Let $\tau = 0.20$. Only $w = \text{mild}$ satisfies $c(w, \text{mild}) \leq 0.20$, hence

$$F^{(0.20)}(\text{mild}) = F(\text{mild}) = \{u_2, u_6\}, \quad |\cdot| = 2.$$

These computations show how increasing τ expands the accepted neighborhood in V and unions the corresponding destination sets in U .

Theorem 2 (CHS generalizes ContraSoft). *Suppose $m = 1$, so $\mathcal{C} = \mathcal{A}_1 =: V$, and let $c_1 = c$. Identify $G : V \rightarrow \mathcal{P}(U)$ with F . Then:*

(a) With the self-centered selector $\rho(v) = v$, one has

$$G_{\rho}^{(\tau)}(v) = \bigcup_{w: c(w,v) \leq \tau} G(w) = F^{(\tau)}(v), \quad \forall v \in V.$$

(b) With the fixed-reference selector $\rho(v) \equiv v^*$, one has

$$G_{\rho}^{(\tau)}(v) = \bigcup_{w: c(w,v^*) \leq \tau} G(w) = F^{(\tau,v^*)}(v), \quad \forall v \in V.$$

Hence, for $m = 1$ the CHS construction recovers both standard *ContraSoft* variants.

Proof. When $m = 1$, $\Delta(w, \rho(v)) = c(w, \rho(v))$. Substituting $\rho(v) = v$ gives (a); substituting $\rho(v) \equiv v^*$ gives (b). The set-theoretic unions agree by definition in both cases. \square

2.2. *Contra-SuperHyperSoft* Set

Contra-SuperHyperSoft Set extends to set-valued coordinates, using lifted subset contradictions and aggregate radius; selector-threshold filtering unions nearby *SuperHyperSoft* slices effectively.

Definition 12 (Base and lifted contradictions). Let A_1, \dots, A_m be nonempty finite sets of attribute values and let $c_i : A_i \times A_i \rightarrow [0, 1]$ be contradiction functions (symmetric and reflexive: $c_i(a, a) = 0 = c_i(a, a)$, $c_i(a, b) = c_i(b, a)$). Assume the zero-separation property $c_i(a, b) = 0 \Rightarrow a = b$. For subsets $S, T \subseteq A_i$ define the lifted contradiction

$$\widehat{c}_i(S, T) := \begin{cases} \max \left\{ \max_{a \in S} \min_{b \in T} c_i(a, b), \max_{b \in T} \min_{a \in S} c_i(a, b) \right\}, & S, T \neq \emptyset, \\ 0, & S = T = \emptyset, \\ 1, & \text{otherwise.} \end{cases}$$

(With finiteness, the max / min are attained.)

Lemma 1 (Symmetry, reflexivity, and zero-separation on subsets). For each i and $S, T \subseteq A_i$:

- (a) $\widehat{c}_i(S, T) = \widehat{c}_i(T, S)$ and $\widehat{c}_i(S, S) = 0$.
- (b) If c_i is zero-separating, then $\widehat{c}_i(S, T) = 0$ implies $S = T$.

Proof. (a) Symmetry follows by exchanging the two max terms; reflexivity is immediate. (b) If $\widehat{c}_i(S, T) = 0$ with $S, T \neq \emptyset$, then $\max_{a \in S} \min_{b \in T} c_i(a, b) = 0$ and $\max_{b \in T} \min_{a \in S} c_i(a, b) = 0$. Thus, for each $a \in S$ there is $b \in T$ with $c_i(a, b) = 0$, hence $a = b$ by zero-separation, so $S \subseteq T$. The second equality gives $T \subseteq S$. The cases with empties are by definition. \square

Definition 13 (Product parameter space and aggregate contradiction). Let $\mathcal{C} := \mathcal{P}(A_1) \times \dots \times \mathcal{P}(A_m)$ and write $\alpha = (\alpha_1, \dots, \alpha_m)$, $\beta = (\beta_1, \dots, \beta_m) \in \mathcal{C}$. Define the tuple-level contradiction by

$$\Delta(\alpha, \beta) := \max_{1 \leq i \leq m} \widehat{c}_i(\alpha_i, \beta_i) \in [0, 1].$$

Then $\Delta(\alpha, \beta) = \Delta(\beta, \alpha)$ and $\Delta(\alpha, \alpha) = 0$; if each c_i is zero-separating, then by Lemma 1 we have $\Delta(\alpha, \beta) = 0 \iff \alpha = \beta$.

Definition 14 (Reference selector). A reference selector is any map $\rho : \mathcal{C} \rightarrow \mathcal{C}$. Two common choices are the self-centered selector $\rho(\alpha) = \alpha$ and a fixed-reference selector $\rho(\alpha) \equiv r$ for a fixed $r \in \mathcal{C}$.

Definition 15 (Contra-SuperHyperSoft Set (CSHS)). *Let U be a finite universe and let $F : \mathcal{C} \rightarrow \mathcal{P}(U)$ be a SuperHyperSoft mapping. Fix contradiction kernels $\{c_i\}$, their lifts $\{\hat{c}_i\}$, an aggregate Δ , a selector ρ , and a threshold $\tau \in [0, 1]$. The associated Contra-SuperHyperSoft Set is the tuple*

$$\text{CSHS} := (U, \{A_i, c_i\}_{i=1}^m, F, \rho, \tau),$$

together with the filtered mapping

$$F_\rho^{(\tau)} : \mathcal{C} \longrightarrow \mathcal{P}(U), \quad F_\rho^{(\tau)}(\alpha) := \bigcup_{\beta \in \mathcal{C} : \Delta(\beta, \rho(\alpha)) \leq \tau} F(\beta).$$

Example 10 (CSHS in E-commerce Fraud Review (self-centered selector)). *Setup.* Let the universe of orders be $U = \{o_1, o_2, o_3, o_4, o_5, o_6, o_7, o_8\}$. Take two attribute domains:

$$A_1 = \{\text{card, crypto}\}, \quad A_2 = \{\text{verified, partial, missing}\}.$$

Base contradictions $c_1, c_2 : [\cdot] \rightarrow [0, 1]$ (symmetric, 0 on the diagonal):

$$c_1 = \begin{array}{c|cc} & \text{card} & \text{crypto} \end{array}, \quad c_2 = \begin{array}{c|ccc} & \text{verified} & \text{partial} & \text{missing} \end{array}.$$

$$\begin{array}{c|cc} \text{card} & 0 & 0.8 \\ \hline \text{crypto} & 0.8 & 0 \end{array} \quad \begin{array}{c|ccc} \text{verified} & 0 & 0.3 & 0.9 \\ \text{partial} & 0.3 & 0 & 0.6 \\ \text{missing} & 0.9 & 0.6 & 0 \end{array}.$$

Lift \hat{c}_i to subsets by Definition (lifted contradiction) and aggregate

$$\Delta((S_1, S_2), (T_1, T_2)) := \max\{\hat{c}_1(S_1, T_1), \hat{c}_2(S_2, T_2)\}.$$

The SuperHyperSoft mapping $F : \mathcal{P}(A_1) \times \mathcal{P}(A_2) \rightarrow \mathcal{P}(U)$ is specified by

$$\begin{aligned} F(\{\text{crypto}\}, \{\text{missing}\}) &= \{o_1, o_2\}, & F(\{\text{crypto}\}, \{\text{partial}\}) &= \{o_3\}, \\ F(\{\text{card}\}, \{\text{missing}\}) &= \{o_4\}, & F(\{\text{card}\}, \{\text{partial}\}) &= \{o_5, o_6\}, \\ F(\{\text{card, crypto}\}, \{\text{partial, missing}\}) &= \{o_7\}, & \text{all other pairs map to } \emptyset. \end{aligned}$$

Choose the self-centered selector $\rho(\alpha) = \alpha$ and threshold $\tau = 0.6$.

Filtering at $\alpha_0 = (\{\text{crypto}\}, \{\text{missing}\})$.

To avoid overfull lines, we list the computations in an aligned display:

$$\begin{aligned} \Delta((\{\text{crypto}\}, \{\text{missing}\}), \alpha_0) &= \max(0, 0) = 0 \leq \tau, \\ \Delta((\{\text{crypto}\}, \{\text{partial}\}), \alpha_0) &= \max(0, c_2(\text{partial, missing})) = \max(0, 0.6) = 0.6 \leq \tau, \\ \Delta((\{\text{card}\}, \{\text{missing}\}), \alpha_0) &= \max(0.8, 0) = 0.8 > \tau, \\ \Delta((\{\text{card}\}, \{\text{partial}\}), \alpha_0) &= \max(0.8, 0.6) = 0.8 > \tau, \\ \Delta((\{\text{card, crypto}\}, \{\text{partial, missing}\}), \alpha_0) &= \max(\hat{c}_1(\{\text{card, crypto}\}, \{\text{crypto}\}), \hat{c}_2(\{\text{partial, missing}\}, \{\text{missing}\})) \\ &= \max(0.8, 0.6) = 0.8 > \tau. \end{aligned}$$

Hence

$$F_\rho^{(\tau)}(\alpha_0) = \{o_1, o_2\} \cup \{o_3\} = \{o_1, o_2, o_3\}, \quad |F_\rho^{(\tau)}(\alpha_0)| = 3.$$

Filtering at $\alpha_1 = (\{\text{card}\}, \{\text{partial}\})$.

β	$\Delta(\beta, \alpha_1)$
$(\{\text{card}\}, \{\text{partial}\})$	$0 \leq \tau$
$(\{\text{card}\}, \{\text{missing}\})$	$\max(0, 0.6) = 0.6 \leq \tau$
$(\{\text{crypto}\}, \{\text{partial}\})$	$\max(0.8, 0) = 0.8 > \tau$
$(\{\text{crypto}\}, \{\text{missing}\})$	$\max(0.8, 0.6) = 0.8 > \tau$
$(\{\text{card}, \text{crypto}\}, \{\text{partial, missing}\})$	$\max(0.8, 0.6) = 0.8 > \tau$

Therefore

$$F_\rho^{(\tau)}(\alpha_1) = \{o_5, o_6\} \cup \{o_4\} = \{o_4, o_5, o_6\}, \quad |F_\rho^{(\tau)}(\alpha_1)| = 3.$$

This illustrates how the CSHS envelope aggregates nearby subset-parameters under the contradiction metric.

Example 11 (CSHS in Cloud Deployment Recommendation (self-centered selector)). **Setup.** Let the universe of candidate nodes be $U = \{n_1, n_2, n_3, n_4, n_5, n_6, n_7\}$. Attributes:

$$A_1 = \{\text{us-east, us-west, eu}\}, \quad A_2 = \{\text{cpu, gpu, memory}\}.$$

Base contradictions (0 on the diagonal, symmetric):

$$c_1 = \begin{array}{c|ccc} & \text{us-east} & \text{us-west} & \text{eu} \\ \hline \text{us-east} & 0 & 0.4 & 0.7 \\ \text{us-west} & 0.4 & 0 & 0.8 \\ \text{eu} & 0.7 & 0.8 & 0 \end{array}, \quad c_2 = \begin{array}{c|ccc} & \text{cpu} & \text{gpu} & \text{memory} \\ \hline \text{cpu} & 0 & 0.5 & 0.3 \\ \text{gpu} & 0.5 & 0 & 0.7 \\ \text{memory} & 0.3 & 0.7 & 0 \end{array}.$$

Lift to subsets by \hat{c}_i and aggregate by $\Delta((S_1, S_2), (T_1, T_2)) = \max\{\hat{c}_1(S_1, T_1), \hat{c}_2(S_2, T_2)\}$.

SuperHyperSoft mapping F (nonempty images shown):

$$\begin{aligned} F(\{\text{us-east}\}, \{\text{cpu}\}) &= \{n_1, n_2\}, \quad F(\{\text{us-east}\}, \{\text{gpu}\}) = \{n_3\}, \\ F(\{\text{us-west}\}, \{\text{cpu}\}) &= \{n_4\}, \quad F(\{\text{eu}\}, \{\text{cpu}\}) = \{n_5\}, \\ F(\{\text{us-east, us-west}\}, \{\text{cpu}\}) &= \{n_6\}, \\ F(\{\text{us-east}\}, \{\text{cpu, gpu}\}) &= \{n_7\}. \end{aligned}$$

Choose the self-centered selector $\rho(\alpha) = \alpha$ and threshold $\tau = 0.5$.

Filtering at $\alpha^* = (\{\text{us-east}\}, \{\text{cpu}\})$. For each β with $F(\beta) \neq \emptyset$, compute $\Delta(\beta, \alpha^*)$:

β	$\Delta(\beta, \alpha^*)$
$(\{\text{us-east}\}, \{\text{cpu}\})$	$\max(0, 0) = 0 \leq \tau$
$(\{\text{us-east}\}, \{\text{gpu}\})$	$\max(0, c_2(\text{gpu, cpu}) = 0.5) = 0.5 \leq \tau$
$(\{\text{us-west}\}, \{\text{cpu}\})$	$\max(c_1(\text{us-west, us-east}) = 0.4, 0) = 0.4 \leq \tau$
$(\{\text{eu}\}, \{\text{cpu}\})$	$\max(0.7, 0) = 0.7 > \tau$
$(\{\text{us-east, us-west}\}, \{\text{cpu}\})$	$\max(\hat{c}_1(\{\text{us-east, us-west}\}, \{\text{us-east}\}) = 0.4, 0) = 0.4 \leq \tau$
$(\{\text{us-east}\}, \{\text{cpu, gpu}\})$	$\max(0, \hat{c}_2(\{\text{cpu, gpu}\}, \{\text{cpu}\}) = 0.5) = 0.5 \leq \tau$

Hence

$$F_\rho^{(\tau)}(\alpha^*) = \{n_1, n_2\} \cup \{n_3\} \cup \{n_4\} \cup \{n_6\} \cup \{n_7\} = \{n_1, n_2, n_3, n_4, n_6, n_7\},$$

with $|F_\rho^{(\tau)}(\alpha^*)| = 6$. Nodes requiring a European region are excluded by the contradiction bound.

Proposition 2 (Monotonicity in the threshold). If $0 \leq \tau_1 \leq \tau_2 \leq 1$, then $F_\rho^{(\tau_1)}(\alpha) \subseteq F_\rho^{(\tau_2)}(\alpha)$ for all $\alpha \in \mathcal{C}$. Moreover, $F_\rho^{(1)}(\alpha) = \bigcup_{\beta \in \mathcal{C}} F(\beta)$ for all α .

Proof. The index sets $\{\beta : \Delta(\beta, \rho(\alpha)) \leq \tau\}$ are nested as τ grows, hence so are the unions. For $\tau = 1$ the constraint is vacuous since $\Delta \in [0, 1]$. \square

Theorem 3 (CSHS generalizes SuperHyperSoft). *Assume zero-separation for each c_i and take the self-centered selector $\rho(\alpha) = \alpha$. Then for $\tau = 0$ we have*

$$F_\rho^{(0)}(\alpha) = F(\alpha) \quad (\forall \alpha \in \mathcal{C}).$$

Proof. By definition, $F_\rho^{(0)}(\alpha) = \bigcup_{\beta: \Delta(\beta, \alpha) \leq 0} F(\beta)$. Since $\Delta \geq 0$, the inequality forces $\Delta(\beta, \alpha) = 0$. By Definition 13 and Lemma 1, this occurs iff $\beta = \alpha$. Hence the union collapses to $F(\alpha)$. \square

Definition 16 (Singleton embedding of HyperSoft into SuperHyperSoft). *Let $\iota : A_1 \times \cdots \times A_m \rightarrow \mathcal{C}$ be*

$$\iota(\gamma_1, \dots, \gamma_m) := (\{\gamma_1\}, \dots, \{\gamma_m\}).$$

We say $F : \mathcal{C} \rightarrow \mathcal{P}(U)$ is singleton-supported for $G : A_1 \times \cdots \times A_m \rightarrow \mathcal{P}(U)$ if $F(\iota(\gamma)) = G(\gamma)$ for all γ and $F(\alpha) = \emptyset$ whenever some α_i is not a singleton.

Lemma 2 (Compatibility of contradictions on singletons). *For any $a, b \in A_i$ we have $\hat{c}_i(\{a\}, \{b\}) = c_i(a, b)$. Consequently, for $\gamma, \delta \in A_1 \times \cdots \times A_m$,*

$$\Delta(\iota(\gamma), \iota(\delta)) = \max_{1 \leq i \leq m} c_i(\gamma_i, \delta_i).$$

Proof. From Definition 12 with singletons, $\max\{\min c_i(a, b), \min c_i(b, a)\} = c_i(a, b)$ by symmetry, and the product case follows. \square

Theorem 4 (CSHS generalizes Contra-HyperSoft). *Let $G : A_1 \times \cdots \times A_m \rightarrow \mathcal{P}(U)$ be a HyperSoft set and F be a singleton-supported extension (Definition 16). Let ρ_H be a selector on $A_1 \times \cdots \times A_m$ and define ρ_S on \mathcal{C} by $\rho_S(\iota(\gamma)) := \iota(\rho_H(\gamma))$, with arbitrary values elsewhere. Then for every γ and $\tau \in [0, 1]$,*

$$F_{\rho_S}^{(\tau)}(\iota(\gamma)) = \bigcup_{\delta: \max_i c_i(\delta_i, \rho_H(\gamma)_i) \leq \tau} G(\delta) =: G_{\rho_H}^{(\tau)}(\gamma),$$

i.e. the CSHS filtered mapping restricted to singleton parameters coincides with the Contra-HyperSoft mapping for G built from $\{c_i\}$ and the same threshold.

Proof. By Definition 15,

$$F_{\rho_S}^{(\tau)}(\iota(\gamma)) = \bigcup_{\beta: \Delta(\beta, \rho_S(\iota(\gamma))) \leq \tau} F(\beta).$$

Since F is singleton-supported, only β of the form $\iota(\delta)$ contribute. By Lemma 2 and the definition of ρ_S , $\Delta(\iota(\delta), \rho_S(\iota(\gamma))) = \max_i c_i(\delta_i, \rho_H(\gamma)_i)$. Substitute and use $F(\iota(\delta)) = G(\delta)$ to obtain the stated equality. \square

Theorem 5 (CSHS generalizes ContraSoft). *Let $m = 1$ with base set $A_1 =: V$ and contradiction $c := c_1$. Let $F : \mathcal{P}(V) \rightarrow \mathcal{P}(U)$ be singleton-supported for some $F_0 : V \rightarrow \mathcal{P}(U)$ via $F(\{v\}) = F_0(v)$ and $F(S) = \emptyset$ if $|S| \neq 1$. Then, with the self-centered selector, for all $v \in V$ and $\tau \in [0, 1]$,*

$$F_{\text{id}}^{(\tau)}(\{v\}) = \bigcup_{w: c(w, v) \leq \tau} F_0(w),$$

which is exactly the neighborhood-based ContraSoft transform on (V, c) . With a fixed reference $v^ \in V$, the same construction yields the fixed-reference ContraSoft variant.*

Proof. Specialize Theorem 4 to $m = 1$ with $G \equiv F_0$ and observe that $\Delta(\{w\}, \{v\}) = \widehat{c}_1(\{w\}, \{v\}) = c(w, v)$. Singleton support reduces the union to singletons, yielding the stated form. \square

3. Conclusion

In this paper, we examined the concepts of the ContraSoft Set, the Contra-HyperSoft Set, and the Contra-SuperHyperSoft Set.

Future work will focus on extending these ideas by integrating them with richer frameworks, including Neutrosophic Sets[12,13], Plithogenic Sets[14–16], Rough Sets[17,18], and TreeSoft Sets[19,20]. Moreover, we anticipate the development of generalized structures that apply these contra-based approaches to Graphs[21,22], HyperGraphs[23–25], SuperHyperGraphs[26–28], and broader HyperStructures[29–31]. Such investigations are expected to open new directions in handling contradiction-aware representations across diverse domains.

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Use of Artificial Intelligence: I use generative AI and AI-assisted tools for tasks such as English grammar checking, and I do not employ them in any way that violates ethical standards.

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