

Article

Not peer-reviewed version

A Different Way to Count, Add, and Multiply

[Gideon Samid](#) *

Posted Date: 3 March 2025

doi: [10.20944/preprints202503.0082.v1](https://doi.org/10.20944/preprints202503.0082.v1)

Keywords: natural numbers; integers; counting; arithmetic operations



Preprints.org is a free multidisciplinary platform providing preprint service that is dedicated to making early versions of research outputs permanently available and citable. Preprints posted at Preprints.org appear in Web of Science, Crossref, Google Scholar, Scilit, Europe PMC.

Copyright: This open access article is published under a Creative Commons CC BY 4.0 license, which permit the free download, distribution, and reuse, provided that the author and preprint are cited in any reuse.

Disclaimer/Publisher's Note: The statements, opinions, and data contained in all publications are solely those of the individual author(s) and contributor(s) and not of MDPI and/or the editor(s). MDPI and/or the editor(s) disclaim responsibility for any injury to people or property resulting from any ideas, methods, instructions, or products referred to in the content.

Article

A Different Way to Count, Add, and Multiply

Gideon Samid *

Electrical, Computer and System Engineering, Computer and Data Sciences, Case Western Reserve University, Cleveland, OH

* Correspondence: Gideon.Samid@CASE.edu

Abstract: We write large numbers by grouping them in exponentially increasing groups: 1, 2, 4, 8, ..., 1, 10, 100, 1000, ..., or any other arbitrary power base. Is there another way? Let's explore incremental, rather than exponential grouping, but pile these groups on top of one another. Group 1 holds one number 1, Group 2 holds the next two numbers: 2, 3, group 3 holds the next three numbers: 4, 5, 6, ... etc. Every natural number will find its place in a group in which it will have its in-group count. E.g. number 5 is count 2 in group 3. Applying iteratively, we have before us a natural way to express integers. History of math has taught us that advanced representation leads to profound insight (e.g. algebra v. arithmetic). It is worthy, therefore, to explore what this grouping, let's call it *numerization* can offer us.

Keywords: natural numbers; integers; counting; arithmetic operations

1. Introduction

The Roman numeric system carried human computations throughout the Roman empire and well into the Middle Ages. It was limited in its expression abilities by the list of symbols used for large numbers. A dramatic revolution in computation, and hence in commerce, and civilization happened when power expressions came forth wherein any number x could be written as:

$$x = \sum a_n b^n \text{ for } n=0,1,2,\dots \text{ where } 0 \leq a_n \leq b-1$$

This exponential method advanced counting by allowing for a limited number of symbols (b) to systematically express any number, however large.

While for human counting the most popular value for the base b is $b=10$, many lament that $b=12$ was not chosen by our ancestors, having more divisors than 10. Today $b=2$, 8 , 16 are common and useful bases.

Innovation Science [8] calls for a revisiting established premises, searching for a useful modification. The essential advantage of the so called, "positional system" of counting is the grouping of numbers to very larger groups 10, 100, 1000, ... then counting these groups.

Grouping can be done on a fixed basis through power raising, but not necessarily so. In the existing system the power base, b , is an arbitrary choice. Is there a natural choice?

Yes, there is: incremental counting. Instead of 0, 1, 2, 4, 8, 16. where counting starts from zero: (0), (0, 1), (0, 1, 2, 3), (0, 1, 2, 3, 4, 5, 6, 7), (0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15), ...

We can count successively:

(0), (1), (2, 3), (4, 5, 6), (7, 8, 9, 10), (11, 12, 13, 14, 15), ...

counting these groups with our numeral symbols:

[0] = (0), [1] = (1), [2] = (2, 3), [3] = (4, 5, 6), [4] = (7, 8, 9, 10), [5] = (11, 12, 13, 14, 15)

Accordingly number 5 is expressed as number count 2 in group 3, and 14 is number count 4 in group 5, etc.

The grouping here increases incrementally not exponentially, but they add-up rather than start from 0. And most importantly this grouping, (let's call it numeration) does not rely on an arbitrarily selected base b -- it is *natural grouping*.



It is of great interest to explore more ways to write down the list of natural numbers. Given the power of AI to find hidden pattern, this *numerization* representation is worth the study.

2. Methodology

Natural number, n is defined as something that is not 1, not 2, not 3,...not $(n-1)$: $n \neq i$. for $i=1,2,\dots,(n-1)$.

These natural numbers are merely a succession of distinctions without any further attribution. In other words, natural numbers are highly abstracted, they only express mutual distinction and order of creation. But they can be further abstracted as *entities that qualify as numbers* but which are not specified as to being one number or another in particular.

So abstracted a natural number is designated by the Hebrew letter "Shin", \aleph , the first letter of the word 'sky' in Hebrew ("Shamaeem", שמיים), which is the first created entity in the Biblical account. Shin represents a number with unspecified designation.

Written as 1,2,3,... the natural numbers are mutually distinct, and no two of them are equivalent. However, listed as. $\aleph, \aleph, \aleph, \dots$ the list comprises elements of equivalence -- all being natural numbers. Equivalence and distinction are in the eyes of the beholder (the mathematician).

Given the unspecified list: $\aleph, \aleph, \aleph, \dots$ the mathematician wishes to express it as a series of natural numbers: 1, 2, 3, ... To that aim the mathematician defines one element as the first \aleph , and calls it "1". In order to define a "2" in the \aleph, \aleph, \dots list the mathematician will need to stack together more than one \aleph , say 2 \aleph s. Now stack 2 that contains two \aleph elements is distinct from the first stack that contains only one \aleph . This distinction is evident to any observer who may be blind towards the number designation of the \aleph entities, and is only aware of the presence of a numeric entity. Stack 2 comprising two numbers is clearly distinct from stack 1 comprising a single number.

And in order to identify a stack designated as "3" the "blind mathematician" will need to stack the next 3 numbers. Thereby stack 3 will be distinct from both stack 2 and stack 1. And so on, number n is comprised of n listed \aleph s, and is thereby distinct from the stacks numbered 1,2,...,(n-1).

The first numerization 1,2,3... originated from one rising out of none, zero, and is called the *zero-numerization* (N0). When the list of natural numbers (N0) is looked upon by an observer that only observes the numbers as equivalent entities (\aleph) where each \aleph has its own location on the list of natural numbers (N0) then, the observer performs a numerization of the N0 list thereby creating *numerization one*, N1: 1, 2, 3, as described above.

When a non-blind numbers observer who sees the specified designation of the \aleph s is examining. N1, she sees the following correspondence:

$$1 \text{ N1} = 1 \text{ N0}$$

$$2 \text{ N1} = 2, 3 \text{ N0}$$

$$3 \text{ N1} = 4, 5, 6. \text{ N0}$$

.....

$$n \text{ N1} = 0.5n(n-1)+1, 0.5n(n-1)+2, 0.5n(n-1)+3, \dots, 0.5n(n-1)+n$$

Every number m in N0 fits into a stack marked as some natural number n in N1. And in that stack, m is the i number in order, where. $1 \leq i \leq n$. m and (n,i) are a bijection. They satisfy the relations:

$$m = 0.5n(n-1)+i; 1 \leq i \leq n$$

Any number m in N0 can be written as a two-numbers tuple (n,i) where n the corresponding stack number in N1, and i is its order in that stack.

We introduce the underscore notation: n_i , where the natural number left of the underscore designates the stack number in N1, and the natural number right of the underscore designates the position, the order of the represented N0 number in that stack. We write:

$$m = n_i$$

where

$$m = 0.5n(n-1)+i; 1 \leq i \leq n$$

We now can map any integer m to a two component tuple n_i . We say that n_i is the N1 representation of the N0 number m .

Let's now discuss an underscore tuple in the form p_q . If $1 \leq q \leq p$ then p and q can be mapped to a number w in $N0$:

$$w = 0.5p(p-1) + q$$

However, if $q > p$ we will replace q with q' using modular arithmetic:

$$q' = q \bmod' p$$

The function MOD' is equivalent to MOD except that $q'=0$ is replaced with $q'=p$.

Thereby any underscore-connected tuple of integers p_q is mapped to a particular $N0$ integer, w . And as we have seen earlier, any integer w can be mapped into an underscore tuple p_q expressed in $N1$.

All together we have established a two tiers numerization which are mutually mappable.

We use the notation n_v to designate an $N0$ number that is in a stack designated as n in $N1$, but is not specified as to which number it is inside stack n . Since number (stack) n in $N1$ is comprising n numbers in $N0$, then v in n_v may be any of n possibilities:

$$n_v = n_i \text{ for } i=1,2,3,\dots,n$$

We also say that the $N1$ stack number of an $N0$ number is its "N1 approximation".

N1 Equivalence: Any two numbers, p , and q in $N0$ which share a stack number in $N1$ are regarded as $N1$ equivalent.

Let $p = n'_i$, and $q = n''_j$. If $n' = n''$ then p and q are $N1$ equivalent even if $i \neq j$. We can write then:

$$0 \leq |i-j| \leq n'-1$$

Negative Numerization: The numerization of $N1$ versus $N0$ is regarded as "positive" and it creates approximation and equivalence among numbers in $N0$. One can then view $N0$ as an approximation of a "negative" numerization $N-1$, such that certain numbers in $N-1$ share the same number in $N0$.

We write $N0: n_v$ to indicate a number in $N-1$ that belongs to stack n in $N0$.

Iteration. Any list of natural numbers $1,2,3,\dots$ may be 'numerized' itself, resulting in a new list of numbers: $1,2,3,\dots$

We say then that numerization of a list of numbers may be iterative, and both ways -- positive and negative. Thereby we define a consecutive series of positive numerizations:

$$N0, N1, N2, \dots, Nr$$

and also negative numerizations:

$$N0, N-1, N-2, N-3, \dots, N-r$$

A number x in $N0$ will be expressed in Nr as follows (standard numerization):

$$\begin{aligned} x &= x'_2 _ x_1 \\ x &= (x'_3 _ x_2) _ x_1 \\ x &= ((x'_4 _ x_3) _ x_2) _ x_1 \\ x &= (((\dots(x_r _ x_{r-1}) _ \dots) _ x_1) _ \dots) _ x_1 \end{aligned}$$

We can omit the parenthesis and agree to fold back the Nr expression from left to right (standard numerization). So we write:

$$x = x_r _ x_{r-1} _ x_{r-2} _ \dots _ x_1$$

where x_1, x_2, \dots, x_r are all natural numbers. And so are x'_2, x'_3, \dots, x'_r

A number x in $N0$ will be expressed in $N-r$ as:

$$x = x _ v_r _ v_{r-1} _ v_{r-2} _ \dots _ v_1$$

where v_i for $i=1,2,\dots,r$ are all unspecified numbers (v).

This will be regarded as natural (standard) expansion. Other possible expansions will be discussed ahead.

Example: Let $x=1,000,000$. ($N0$). Writing x in $N1$ comes to: 1414_1009 .

To represent x in $N2$ one will numerize 1414 :

$$1414 = 53_36$$

So $1,000,000$ $N0 = 53_36_1009$. ($N2$)

And further, to represent x in $N3$, one needs to numerize $53 = 10_8$

So we can write:

$$1,000,000 N0 = 10_8_36_1009 N3$$

And further: $10 = 4_4$, so:

$$1,000,000 N0 = 4_4_8_36_1009 N4$$

r-Level Equivalence Let p and q be two numbers in $N0$, written in Nr as follows:

$$p = p_r_p_{r-1}_p_{r-2}_....._p_r$$

and

$$q = q_r_q_{r-1}_q_{r-2}_....._q_r$$

Let it be that $p_i \neq q_i$ for $i=1,2,...,(r-1)$ while $p_r = q_r$

We will then state that p and q are Nr equivalent.

It is easy to see that however far apart p and q may be there is always an integer r such that p and q are Nr equivalent.

Nr Distance: Let the numbers p and q above comply with: $p_i \neq q_i$ for $i=1,2,...,(r-1)$ and $p_r = q_r$, then we say that the Nr distance between p and q , $D_r(p,q) = |p_r - q_r|$

If $D_r(p,q) = 0$ then p and q are Nr equivalent.

Abstracting the Positional Numeral System: When the Indians and Arabians during the 5th to 7th century introduced and applied the positional numeral system -- mathematics and its dependent fields have been catapulted into new horizons. The numerization ladder presented herein may be viewed in some respect as an abstraction of this old method. Instead of using an arbitrary basis for representing large count with a tuple of small counts, we here use a 'natural way' for doing so, hoping for it to open up roadways to new numerical knowledge.

Here below is a graphic representation of four rounds of numerization:(Figure 1)

Numerization			
N0	N1	N2	N3
1	1	1	1
2	2	2	2
3			
4	3		
5			
6			
7	4	3	
8			
9			
10			
11	5		
12			
13			
14			
15			
16	6		
17			
18			
19			
20			
21			

Figure 1. 4 rounds of Numerization.

3. From Distinction to Equivalence

Any two arbitrary numbers x , and y where $x \neq y$ are distinct. We explore two ways to extract equivalence between them.

1. Approximation, de-specification.
2. De-Approximation, specification

The first method calls for approximating both x and y to such degree that a number a is cast as an approximation of both x and y . The fewer the approximation steps the closer x and y are. Approximation amounts to de-specification, ignoring and removing specificity from both x and y .

The second method is based on the fact that both x and y may be viewed as different approximations of a pre-approximated number p . p was approximated one way from p to x , $p \rightarrow x$, and approximated another way from p to y : $p \rightarrow y$.

Distinction and equivalence are the essential building blocks for mathematical construction.

4. Approximation & Equivalence

Positive numerization leads to approximation and expanded equivalence, while negative numerization leads to specification and established distinction.

Given a natural number x . By writing it in $N1: y_z$, one identifies y as the first numerization grade approximation of x . There are x numbers which share same approximation. Say, with the reduced distinction expressed by $N1$ approximation, one creates an equivalence between two numbers x_1 and x_2 , written as $x_1 = y_{z_1}$, and $x_2 = y_{z_2}$.

Positive numerization, then, may be regarded as approximation, creating equivalence.

Any arbitrary number $x_i = x_{i+1}x'_i$, obeys: $x_{i+1} \leq x_i$ where the equality only applies for $x_i = 1$. Accordingly any arbitrary number $x > 1$ will eventually be approximated to: $2_11_11\dots$ Hence any two numbers x and y will become equivalent after t rounds of positive numerization.

Example: $x = 247837$, $y = 8976435$

We write:

$$247837 = 704_381 = 38_1_381 = 9_2_1_381 = 4_3_2_1_381 = 3_1_3_2_1_381$$

$$8976435 = 4237_2469 = 92_51_2469 = 14_1_51_2469 = 5_4_1_51_2469 = 3_2_4_1_51_2469$$

It takes 5 numerization rounds to achieve equality. We write:

$$247837 N5 = 8976435 N5$$

But if we compare y to $z = 6784113$ then we write:

$$6784113 = 3684_27 = 86_29_27 = 13_8_29_27 = 5_3_8_29_27$$

$$8976435 = 4237_2469 = 92_51_2469 = 14_1_51_2469 = 5_4_1_51_2469$$

Recording a proximity of degree 4: $8976435 N4 = 6784113 N4$

Given two arbitrary numbers x and y expressed in Numerization Nr :

$$x_r x_{r-1} \dots x_1 y_r y_{r-1} \dots y_1$$

where $x_r = y_r$, we then regard $|x_{r-1} - y_{r-1}|$ as the 1st approximate distinction between x and y . and in general, we regard:

$$|(x_{r-1} x_{r-2} \dots x_{r-i}) - (y_{r-1} y_{r-2} \dots y_{r-i})|$$

As the i -th approximate distinction between x and y .

Approximation establishes equivalence by becoming blind to distinctive details.

5. Specification & Distinction

Numerization implies specification and distinction upon adding a right side number to the underscore tuple.

We have seen that given two arbitrary numbers x and y expressed in Numerization Nr :

$$x_r x_{r-1} \dots x_1 y_r y_{r-1} \dots y_1$$

where $x_r = y_r$, we regard

$$|(x_{r-1} x_{r-2} \dots x_{r-i}) - (y_{r-1} y_{r-2} \dots y_{r-i})|$$

As the i -th approximate distinction between x and y .

We may apply this distinction over negative numerization much the same. So given two indistinct numbers $x=y$, we can build a distinction between them as x_i and y_j . where $i \neq j$. and $1 \leq i, j \leq x=y$. We say then that x_i and y_j are first degree distinct.

If $i=j$, then there is no distinction between x_i and y_j . However we can write x_{i-k} and y_{j-l} where $k \neq l$. and $1 \leq k, l \leq x_i = y_j$. We say then that x_{i-k} and y_{j-l} are 2nd degree distinct.

And similarly we can establish fine distinctions between otherwise indistinct numbers.

6. Equivocation

The number $y = x_{\mathcal{W}}$ is interpreted as a number that can be $0.5x(x-1)+1, 0.5x(x-1)+2, \dots, 0.5x(x-1)+x$. This is an extension of the notion of *qubit*. Equivocation can be extended:

$$z = x_{\mathcal{W}}_{\mathcal{W}} = y_{\mathcal{W}},$$

so z is equivocated over t numbers:

$$t = \sum_i (0.5x(x-1)+i) = 0.5x^2(x-1) + 0.5x(x+1) = 0.5x(x^2+1)$$

With equivocation growing exponentially for: $w = x_{\mathcal{W}_1}_{\mathcal{W}_2}_{\dots}_{\mathcal{W}_r}$

The equivocated numbers are consecutive and range from $x+1$ to $x+t$. Such equivocation numerization creates a range of arbitrary height for every x . Given two arbitrary numbers x , and y such that $y > x$, then x can be equivocated through t_x rounds and y can be equivocated through t_y rounds leading for the two equivocation ranges to overlap. There will be n arbitrary numbers z_1, z_2, \dots, z_n where the x -equivocation range and the y -equivocation range overlap, where n may be made as large as desired, by increasing the values of t_x and t_y .

Example: let $x=3$ for $t=2$ the equivocated numbers are: 7, 8, 9, 10, ... 21. Let $y=4$. for $t_y=1$ the equivocated numbers are: 7, 8, 9, 10. $n=4$ numbers overlap.

Computers which can handle computations with such equivocated entities may benefit from this framing of natural numbers.

7. Arithmetic

Nominal arithmetic can be readily extended to numerized expressions. Given two arbitrary numbers x and y expressed in Numerization Nr:

$$x_r x_{r-1} \dots x_1 y_r y_{r-1} \dots y_1$$

we will define Nr operations (addition, subtraction, multiplication, division, power raising) involving two Nr numbers resulting in an Nr number.

For $y_i = 0$ for $i=1,2,\dots,r$. ($0_0 \dots 0$) we write $x = x + y$, and for $y_i = 1$ for $i=1,2,\dots,r$ ($1_1 1 \dots 1$) we have $x = y^*x = x/y$.

Nr numerization addition of x and y : $z = x + y$. Nr:

$$z_i = x_i + y_i \text{ for } i=1 \text{ to } r$$

7.1. Addition/Subtraction

Given two arbitrary numbers x and y expressed in Numerization Nr:

$$x_r x_{r-1} \dots x_1 y_r y_{r-1} \dots y_1$$

we will define Nr numerization addition of x and y : $z = x + y$. Nr:

We first normalize both x , and y . Normalization ensures that for $x_i x_{i+1}$, we have: $x_i \leq x^*_{i+1}$. where:

$$x^*_{i+1} = (x_r x_{r-1} \dots x_{i+2} x_{i+1})$$

Namely x^*_{i+1} represents the N0 expression of the tuple x_r, \dots, x_{i+1} We set:

$$z_i = x_i + y_i \text{ for } i=1 \text{ to } r$$

Note that:

$$x + y \text{ N0} < x + y \text{ N1} < x + y \text{ N2} \dots$$

Example: $5768 + 9823 = 15591$

For N1 we write: $5768 = 197_7$; $9823 = 140_93$. Hence: $z_0 = (197+140)(7+93) = 337_100 = 56716$.

For N2 we write: $5768 = 20_7_71$. $9823 = 17_4_93$. Hence. $z_0 = (20+17)(7+4)(71+93) = 37_11_164 = 677_164 = 228990$

For $r=0,1,2$ respectively we have $5768 + 9823 = 15591, 56716, 228990$

Subtraction: Subtraction is defined opposite to addition: Given two arbitrary numbers x and y expressed in Numerization Nr: we first normalize x and y as defined above, (with the MOD' function) yielding the normalized tuples:

$$x_r x_{r-1} \dots x_1 y_r y_{r-1} \dots y_1$$

we will define Nr numerization subtraction of x and y : $z = x - y$. Nr:

$$z_i = x_i - y_i \text{for } i=1 \text{ to } r$$

This definition is straight forward as long as $x_i > y_i$ for $i=1,2,\dots,r$

Otherwise we need more mathematical definition.

Zero: Introducing the notion of zero. (0). Zero was described as the state in which "1" rose. It stands mentally to represent 'nothing'. Nothing cannot be properly represented because any representation will amount to something. Nothing is a limiting state of having as minimum as possible of anything. We will use the symbol zero to represent this limiting nothing.

Accordingly:

$$\begin{aligned} 0 &= 0_0_0_\dots \\ 0_n &= 0_0_0_\dots_0_n = n \\ n_0 &= n_0_0_\dots_0 = n \end{aligned}$$

for $n \neq 0$.

Thus: $x - x$ Nr = $0_0_\dots_0 = 0$. (a tuple of r underscores).

Negative Natural Numbers: We build a symmetric twin for the list of natural numbers: 1, 2, 3, ... marked as -1, -2, -3, ... and placed in a symmetric fashion with the symbol zero in between:

$$-n, -(n-1), \dots, -3, -2, -1 \ 0 \ 1, 2, 3, \dots, (n-1), n$$

The symmetric twin (negative natural numbers) is operated on the same as the positive integers (ignoring the minus sign), but when done, the minus sign is attached.

We denote the numerization of the negative numbers as, $N-1, N-2, \dots$

By definition x Nr - x N-r = 0.

$$x_i - x_{-i} = 0 \text{ for } i=1,2,\dots,r$$

Subtraction is carried out from left to right as follows: if $\text{sign} = x_r - y_r > 0$ then $z = x - y$ is a positive integer. If $\text{sign} < 0$ then z is a negative integer. If $\text{sign} = 0$ then x_r and y_r are ignored and x and y are treated as $N(r-1)$ numerizations:

$$x_{r-1}x_{r-2}_\dots\ x_1 \ y_{r-1}y_{r-2}_\dots\ y_1$$

Normalization should be followed. including negative number. Example: $3_5 = 3_2, 3_8 = 3_1$

As defined x_r and y_r determine if z is positive integer or negative integer.

Example:

$$x = 459 = 30_24 = 8_2_23$$

$$y = 244 = 22_13 = 7_1_13$$

Both x and y are normalized since $(8_2) > 23$, and $(7_1) > 13$. so $z = 1_1_10$ N3 = $0_1_10 = 1_10 \rightarrow$ normalizing = $1_1 = 1$

Now we try with $y = 1500 = 55_15 = 10_10_15$

$\text{sign} = x_3 - y_3 = 8 - 10 = -2$, which is < 0 therefore $z < 0$

We now subtract $(8_2_23) - (10_10_15)$. We concluded $z_3 = -2$. $z_2 = 2 - 10 = -8$, normalizing over $z_3 = -2$ we have $x_2 = -2$. $z_1 = 24 - 15 = 9$, normalizing over $(-2_2) = 3$ we get $z_1 = 1$. Hence: $z = -2_2_1$.

7.2. Multiplication/Division

Given two arbitrary numbers x and y expressed in Numerization Nr:

$$x_{r-1}x_{r-2}_\dots\ x_1 \ y_{r-1}y_{r-2}_\dots\ y_1$$

we will define Nr numerization multiplication of x and y : $z = x * y$ Nr:

$$z_i = x_i * y_i \text{ for } i=1,2,\dots,r$$

And normalization of the resultant z tuple.

Example:

$$x = 459 = 30_24 = 8_2_23$$

$$y = 244 = 22_13 = 7_1_13$$

Both x and y are normalized since $(8_2) > 23$, and $(7_1) > 13$.

$$z = (8*7)_(2*1)_(23*13) \ N3 = 56_2_299 = 1542_299 = 1188410$$

larger than $x * y$ N0 = $459 * 244 = 111996$

Division: For $z \text{ Nr}$ which is the multiplication of $x \text{ Nr}$ over $y \text{ Nr}$, we define division as a multiplication-reverse:

$$y = z/x \text{ N3}; x = z/y \text{ N3}$$

For a general tuple $z \text{ Nr}$ and $x \text{ Nr}$ the division y/x is carried on as follows:

Case 1 all $y_i = z_i/x_i$ for $i=1,2,\dots,r$ are integers. In that case

$$y = y_r _ y_{r-1} _ \dots _ y_1$$

Case 2. For $i = r, (r-1), (r-2), \dots, (i+1)$ $y_i = z_i/x_i$ are all integers while z_i is not an integer for $i=r$

In this case there are two division options: (i) approximation, (ii) specification.

In approximation we write:

$$y = y_r _ y_{r-1} _ \dots _ y_i _ 0 _ 0 _ \dots _ 0$$

In specification we write:

$$*z_i = z_r _ z_{r-1} _ \dots _ z_i _ \text{w}$$

$$*x_i = x_r _ x_{r-1} _ \dots _ x_i _ \text{w}$$

And we seek to replace the two w 'place holders' with z_i and x_i such that $y_i = z_i/x_i$ is an integer. If no such two specification numbers can be used to replace the w s then, we extend another w and write:

$$*z_i = z_r _ z_{r-1} _ \dots _ z_i _ \text{w} _ \text{w}$$

$$*x_i = x_r _ x_{r-1} _ \dots _ x_i _ \text{w} _ \text{w}$$

and try all possible w options to find $y_i = z_i/x_i$ as an integer. If no such four specification numbers can be used to replace the w s then, we extend another w and keep the extension until y_i is an integer.

And then continue with z_j and x_j to set y_j and on for $j = (i-1), (i-2), \dots, 1$

Example: $z = 12_4_18$, $x = 3_4_3$:

$$y \text{ N2} = (12/3) _ (4/4) _ (18/3) = 4_1_6 = 7_6 = 27 \text{ N0}$$

while for $z = 12_4_18 = 70_18 = 2433 \text{ N0}$. And $x = 3_4_3 = 7_3 = 24 \text{ N0}$ we write:

$$y \text{ N0} = (z \text{ N0}) / (x \text{ N0}) = 2433/24 = 101.375$$

But for $x \text{ N3} = 3_4_5$ we have:

$$y \text{ N3} = (12/3) _ (4/4) _ (18/5)$$

Since $18/5$ is not an integer we may choose the approximation method:

$$y \text{ N3 -- Approximation} = (12/3) _ (4/4) = 4_1 = 7$$

Or we may use specification method:

$$z = 12_4_18 _ \text{w} _ x = 3_4_3 _ \text{w}$$

We write:

$$z = 2433 _ \text{w} _ x = 24 _ \text{w}$$

We write:

$$z = 2433 _ 109 = 2958637$$

$$x = 24 _ 1 = 277$$

$$y = z/x = 2958637/277 = 10681$$

7.3. Power Raising

Given arbitrary x in Nr taken in normalized form:

$$xr _ xr-1 _ \dots _ x1$$

It can be raised by a power $p \text{ Nr}$ taken in normalized form:

$$pr _ pr-1 _ \dots _ p1$$

$y = x^p \text{ Nr}$ where: $y_i = (x_i)^{p_i}$.

Example: $x \text{ N2} = 23_4_11 = 257_11 = 32907$

$$p \text{ N2} = 2_2_1 = 3_1 \text{ N2} = 4$$

$$y = x^p = (232) _ (42) _ (111) = 529_16_11 = 139672_11 = 9754063967$$

while $y = x^p \text{ N0} = 32907^4 = 1172608842465681201$

Root Extraction: Given that $y = x^p$, we can say that $x = y^{1/p}$. In general root extraction will be handled just like division, ensuring an integer based on one of the two formerly discussed methods: approximation or specification.

8. Numerized Arithmetic

Variables and coefficients used in algebraic expressions can each be written in any non-negative numerization, Nr. Addition, subtraction, multiplication, division, power raising, root extraction are all well-defined for the various numerized expressions of integers. So one could manipulate symbols, calculate and process integers in any arbitrary numerized expression. This includes solving algebraic arithmetic questions.

Furthermore one could, at any point, switch numerized representation from some value N_i to some value N_j where $i \neq j$, and then to N_k where $k \neq i, k \neq j$. An algebraic operation O_1 can be done with numerization r_1 , then one may switch to numerization r_2 , then one may continue to process with operation O_2 , followed by a switch to numerization level r_3 , and so on. The result will depend on the shift pattern between numerization levels.

Example: $y = f(x) = 3x^2 + 2x + 4$

Compute $y = f(x)$ N_0 for $x=2$:

$$y = 3*2^2 + 2*2 + 4 = 20$$

Compute $y=f(x)$ N_1 for $x=2$

$$\begin{aligned} y N_1 &= (2_2)*(2_1)^{(2_1)} + (2_1)*(2_1) + (3_1) \\ &= (2_2)*(4_1) + (4_1) + (3_1) = (8_2) + (4_1) + (3_1) = (15_4) = 109 \end{aligned}$$

Multi Level Arithmetic An arbitrary number x N_r may be readily mapped to x N_p for any $p \neq r$. A given set of n variables V_1, V_2, \dots, V_n N_r may be operated on in N_r (O_r) to generate a set of m variables U_1, U_2, \dots, U_m N_r . These m variables may then migrate to N_p and be operated on with operation O_p to generate a set of q variables W_1, W_2, \dots, W_q N_p . And so on, operation after operation where each successive operation operates on a different numerization level.

Such numerization migration defines a *multi level arithmetic*. It defines an operational wealth for which the common arithmetic is a collapsed version.

Much as arithmetic abstracts itself to algebra so does multi-level arithmetic abstracts itself to multi-level algebra.

Numerized Functional Relationships

Let X and Y be two arbitrary integers in N_r . $x = x_r _ x_{r-1} _ \dots _ x_1$ $Y = Y_r _ Y_{r-1} _ \dots _ Y_1$. We may identify a variety of functional relationships between X and Y . For example: MAX, MIN, MOD.

Illustration: Let $X = 12_4_7$ and $Y=8_14_3$. We write:

$$Z = X \text{ MAX } Y = (12_4_7) \text{ MAX } (8_14_3) = 12_14_7$$

$$\text{We also write: } Z = X \text{ MOD } Y = (12_4_7) \text{ MOD } (8_14_3) = (4_4_1)$$

9. Attributes

Numeric attributes common with N_0 can be extended to N_r . We discuss, primes, composites, multiplications, fractions, as well as more esoteric attributes like uniform numbers and perfect numbers.

9.1. Primes

An arbitrary number x N_r where at least for some $i=1,2,\dots,r$. x_r is prime is regarded as "Base prime". since there is no y N_r number for which $y_i > 1$ for $i=1,2,\dots,r$ such that x/y is an integer.

A base prime N_r may be a composite number for N_p for $p < r$. Example: $x = 7_5_3$ N_2 (prime) = 26_3 N_1 (prime) = 328 N_0 (composite).

An arbitrary number x N_r where for all $i=1,2,\dots,r$. x_r is prime is regarded as "Full Prime". since there is no y N_r number for which $y_i > 1$ for which any x_i/y_i is an integer for any $i=1,2,\dots,r$

9.2. Composites

An arbitrary number x Nr where for $i=1,2,\dots,r$. x_i is composite is regarded as "Nr Composite" since there are one or more y Nr numbers for which $y_i > 1$ for $i=1,2,\dots,r$. such that x/y is an integer Nr.

Example: Let $x = 6_9_15$. We have:

for $y = 2_3_5$. $x/y = (6/2)(9/3)(15/5) = 3_3_3$.

9.3. Uniform Numbers

A number x Nr is considered uniform for x' if and only if $x'_i = x_i$ for $i=1,2,\dots,r$.

The series of uniform numbers is: $0_0_0, 1_1_1, 2_2_2, 3_3_3, 4_4_4, \dots, N_0$
which are: 0, 1, 5, 18, 49, N0

9.4. Fractions

Writing two Nr numbers, x, y , in fraction form, defines thereby a fraction, f_{xy}

$x_r x_{r-1} \dots x_1 y_r y_{r-1} \dots y_1$

$f_{xy} = (x_r x_{r-1} \dots x_1) / (y_r y_{r-1} \dots y_1)$

10. Applications

Applications involving approximation, specification, and obfuscations are all of interest in the numerized realm. Same for applications where one digs for unseen patterns. Hence numerization appears useful for probability calculus, quantum computing, AI inference. The field is attractive for cryptographic primitives.

10.1. Error Assessment

Given x Nr: $x_r x_{r-1} \dots x_1$

Let there be an error in some x_i , which is written as x'_i . where. $e = x'_i - x_i$.

The higher the value of i , the greater the impact of the error on x N0.

Example: let $x = 5_3_3 = 13_3 = 81$

let $e=1$ and $i=3$ hence $x_e = 6_3_3 = 18_3 = 156$

An error in x . N0 of 92% while for $5_3_4 = 5_3_1 = 18_1 = 79$ An error of $81/79 = 2\%$

10.2. Numerization Modes

In addition to the standard numerization mode in which each round the leftmost number is numerized, we can apply other modes. In particular: (i) double sided numerization, (ii) centric numerization, and (iii) full line numerization. In the first mode, in each numerization round both the rightmost and the leftmost number are numerized. In the second numerization the two most centric numbers in the tuple are numerized each round. In the third mode each round all numbers are numerized. Once the mode is known, these modes allow for reversal -- restoring the pre-numerization number.

Example: Nominal Numerizatio

$7892323817 = 125637_58751 = 501_387_58751 = 32_5_387_58751 = 8_4_5_387_58751 = 4_2_4_5_387_58751$

Double side numerization:

$7892323817 = 125637_58751 = 501_387_343_98 = 32_5_387_343_14_7$

Centric numerization:

$7892323817 = 125637_58751 = 501_387_343_98 = 501_28_9_26_18_98 = 501_28_4_3_7_5_18_98$

Full Line numerization

$7892323817 = 125637_58751 = 501_387_343_98 = 32_5_28_9_26_18_14_7$

10.2.1. Variations

Numerization may be explored through a series of variations. We mention here shift numerization, split numerization, circular numerization. Shift numerization is the situation where the counting shifts from 1,2,3,... to some number n , $n+1$, $n+2$, for $n > 1$. Split numerization refers to selecting an arbitrary point as 'zero' and applying positive numerization to its right and negative numerization to its left. Circular numerization describes numerization over a modular ring of integers.

10.2.2. Binary Strings Numerization

Numerization can be used to write binary strings differently, for any prospective purpose.

Consider the following binary string:

$x = 1100110011101100011011110010100111011010101011110$

Its numeric value is: $x = 20303025270$ which can be Numerized in the full front method:

$x = 20303025270 = 201509_187484 = 635_214_612_518$

Onward:

$x = 36_5_21_4_35_17_32_22 = 8_8_3_2_6_6_3_1_8_7_6_2_8_4_7_1$

This string can be expressed using binary alphabet by writing every one of the 16 numbers above through a unary alphabet flipping between the bits identity. So we start with 8 zeros followed by 8 "ones" followed by 3 zeros, etc:

$x = 0000000011111110001100000011111000000110000000111100000001$

The above string is content identical to the original x string

Alternatively, given that the numerization proceeded such that 8 is the largest number, one can express this numerization by a series of 3 consecutive bits, per each numerization number. Since 0 is excluded from the numerization, we express the series, $8_8_3_2\ldots$ as 000 000 011 010.. which computes to:

$x = 0000000110101101001100100011110010000010111001$

The above string also is of the same content as the original string x (bijection)

10.3. Estimation

When x represents a measurement, or a calculation of a well-defined value. There may be a difference between the 'true' value one tries to measure or calculate, x^{\wedge} , and x . But when both x and x^{\wedge} are numerized then x_i is more likely to coincide with x_i^{\wedge} , the higher the value of i .

Example. Let $x^{\wedge} = 50$ and $x = 46$, we write $x^{\wedge} = 50 = 10_5$, and $x = 46 = 10_1$ We see that

$$x_2 = x^{\wedge}_2 = 10$$

We therefore can reduce 10_5 to $10_{\mathfrak{v}}$ to indicate that x_1 in the measurement or calculation are of too low validity, and should not qualify as data. We can then replace the \mathfrak{v} with the top of the x_1 range, the bottom, or the middle, per an arbitrary decision. (Note: nominally the middle range is best, it has the shortest distance from the unknown true value. For even count select top or bottom). And then we use this choice in rewriting x . So instead of $x=46$ we write $x = 10_6 = 51$.

To express even greater doubt about the accuracy of the measurement or calculation we may go to N2:

$$46 = x_3_x_2_x_1 = 10_1 = 4_4_1$$

And write $x = 4_{\mathfrak{v}}_{\mathfrak{v}}$. Similarly. $x^{\wedge} = 4_4_5$

Now writing $x = 4_4_10 = 55$

A more common approach uses the negative numerization to chart a range. For numerization level N-1 we write $x = x_2_{\mathfrak{v}}$, and replace \mathfrak{v} once with 1 and once with x_2 to define a range.

In the example above it figures as a range from 10_1 to 10_10 , namely: 46-55, which includes x'

$$10_1 \leq x' \leq 10_10 \text{ Namely: } 46 \leq x' \leq 55$$

If we are less confident about our measurement we will go down to N-2: $x = x_2_{\mathfrak{v}}_{\mathfrak{v}}$, and again define the range between the lowest possible values for \mathfrak{v} and the highest.

Writing the measurement x as $x = 46 = 10_1 = 4_4_1 = 4_{\text{v}}_{\text{v}}$, we identify the bottom of the range: $4_1_1 = 22$, and the top of the range: $4_4_{10} = 55$, writing $22 \leq x' \leq 55$
The same may apply to N-3, N-4, etc.

10.4. Non-Linear Addition, Multiplication.

In quite a few situations the linear addition $z = x + y$ is non-reflective of reality. At times z should be larger than the linear addition, and at times smaller, same for multiplication. When adding risk for example, combined risk is larger than linear add-on, in adding probabilities, a range of overlapping cases will result in a lower probability than linear add on. Numerization arithmetic can readily be used to express such non-linear situations.

Given $z = x + y$. N0, one may map, x , and y to N1: $x_1_{x_2}$, $y_1_{y_2}$, and add: $z = x + y$ N1. and then map z N1 to $(z$ N1) N0, where $(z$ N1) N0 $>$ z N0. This non-linearity may be extended at will by choosing a higher numerization level i :

$$(z \text{ N}_i) \text{ N}_0 > (z \text{ N}_j) \text{ N}_0 \text{ for } i > j$$

Example: $z = x + y$; $x = 54$, $y = 85$. $z = 139$.

We write: $54 = 10_9$, $85 = 13_7$ and so $z = (10+13)_{(9+7)} = 23_{16} = 269$

Indeed $269 > 139$.

Going further to N2: $54 = 4_4_9$; $85 = 5_3_7$ and so $z = (4+5)_{(4+3)}_{(9+7)} = 9_7_{16} = 43_{16} = 919$

So the N0 mapping of z N2, z N1, z N0 is 919, 269, 139

For situations where the nonlinearity is in the opposite direction, to suppress the outcome, one would treat x and y as the x_2 and y_2 values respectively of negative numerization expansion N-1: $x' = x_2_{x_1}$, and N-1 $y' = y_2_{y_1}$ where $x_1 = y_1 = \text{v}$. Given that $1 \leq x_1 \leq x_2$, one can choose any value within this range, nominally -- the top of the range, so $x_1 = x_2$, and similarly $y_1 = y_2$.

And to increase the suppression one might map x and y in N0 to $x_2_{\text{v}}_{\text{v}}$ and $y_2_{\text{v}}_{\text{v}}$, and so on. In general:

$$(z \text{ N}_i) \text{ N}_j < z \text{ N}_j \text{ for } i, j = 1, 2, \dots \text{ and } i > j$$

Example: $z = x + y$; $x = 54$, $y = 85$. $z = 139$.

We write. x N0 = 54_{v} , y N0 = 85_{v} , becoming x N0 = 54_{54} , and y N0 = 85_{85}

Mapping $(x$ N0) N-1 = $54_{54} = 1485$. And mapping $(y$ N0) N-1 = $85_{85} = 3655$

And hence z N-1 = $(x$ N0) N-1 + $(y$ N0) N-1 = $1485 + 3655 = 5140$

Further. $(z$ N-1) N0 = $5140_{N-1} = 101_{90}$

which becomes 101_{v} , converting to 101 which is smaller than $x+y = 139$

Going for N-2, x , and y are written as: $x_2_{\text{v}}_{\text{v}}$, and $y_2_{\text{v}}_{\text{v}}$, where we replace the v with the top of the range: x N0 = 54_{54}_{1485} and y N0 = 85_{85}_{3655} , and the corresponding numbers at N-2 are:

$$54_{54}_{1485} = 1485_{1485} = 1103355 = x$$
 N-2

$$85_{85}_{3655} = 3655_{3655} = 6681340 = y$$
 N-2

And so:

$$z$$
 N-2 = $1103355 + 6681340 = 7784695$

for which the corresponding N0 is: $89_{30}_{1210} = 89_{\text{v}}_{\text{v}}$.

So in summary $x+y$ plain (N0), N-1, N-2 values are: 139, 101, 89

Similarly for multiplication both expansion and contraction option are numerization enabled.

10.5. Re-Anchoring

Numerization may be re-anchored anywhere on the original list of natural numbers. Anchor A means that A+1 is the integer deemed as 1 for numerization process, A+2 is integer 2, A+i integer i, and A is represented by the 0 indication in numerized vocabulary.

Everything that takes place in pre-anchoring numerization has a corresponding event in anchored representation.

There is a clear mapping between any numerized tuple before anchoring and after anchoring as well as between two different anchored designations.

Example: $x = 111$, Pre-anchoring $x N1 = 15_6$. Let $A = 48$, namely $x N0 (A=48) = 63$ $N0 (A=48)$ so: $x N1 (A=48) = 11_8$, and similarly:

Pre A 15_7 = 11_9 (A=48).
 Pre A 15_8 = 11_10 (A=48)
 Pre A. 15_9 = 11_11 (A = 48)
 Pre A. 15_10 = 12_1 (A=48)

10.6. Pattern Recognition

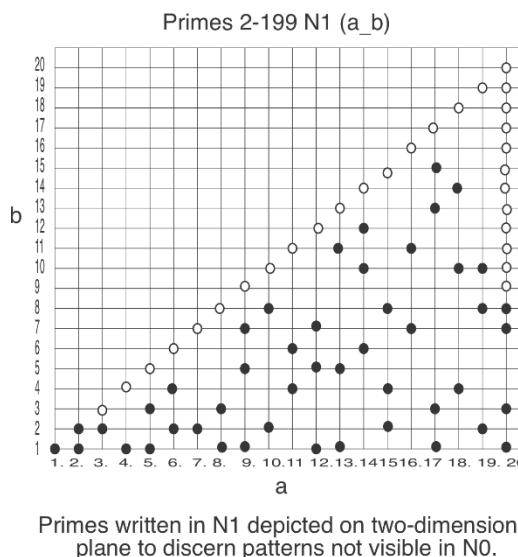
Let there be n measurements. x_1, x_2, \dots, x_n , of n items, each with respect to a particular attribute of the items. And let there be m measurements of the same attribute of different m items y_1, y_2, \dots, y_m . The n first items belong to a category X and the m latter items belong to a category Y. One tries to find a discriminating pattern between the two categories so that given another measurement z , one will be able to state with some confidence whether the z measured item belongs to category X or category Y.

This cluster analysis is normally done over the nominal measured values, all in $N0$. However, mapping these $n+m+1$ measurements each to Nr will offer the pattern analyzer a much richer domain in which to discern pattern. The Nr expression of the measurements will depict the $n+m+1$ cases as $n+m+1$ points on an r -dimensional space exposing clusters not visible with $N0$.

This enhanced pattern recognition is widely applicable. For fractional measurement, a multiplication by all denominators will yield natural numbers that are readily numerized.

Example, for measurement. 1.44, 0.28, 1.11, multiplying $\times 100$ yields: 144, 28, 111.

The numerization born pattern recognition may also be used to find profound mathematical attributes like prime numbers.



11. Analysis

This numerization construction is faithful to the unmatched clarity of natural numbers, a succession of distinction -- the foundation of mathematics. It stops short of the common extension to irrational numbers, continuity, infinity, motion. This extension is well handled by mathematical formalism, but at the cost of the conceptual clarity that is claimed by the natural numbers.

Historically the battle between the discrete and the continuous was raging for centuries, either notion regards the other as an approximation. A no lesser authority than David Hilbert asserted:

"Our First Naive impression of Nature and matter is that of continuity. Be it a piece of metal or a volume of liquid, we invariably conceive it as divisible into infinity."

It stands to reason to explore mathematics from either direction. This thesis does so from the side of natural numbers, with numerization offers them a new wealth of expression to handle situations more naturally handled through continuity, irrationality, infinity.

The main interest in numerization is its aspect of novelty regarding the fundamental entities of mathematics: the series of natural numbers. They harbor secret, patterns, relationships which are still hidden from the common knowledge of mathematics. Any novel method that writes and manipulates these natural entities with minimum arbitrariness and maximum naturalness is of interest because of its potential to reveal numeric properties not yet known.

Numerization keeps its analysis in the realm of natural numbers, avoiding the 'continuity trap' that leads to irrational numbers. It introduces two opposing complimentary concepts: approximation and specifications with which it meets the known limitations of integers normally solved with irrational numbers, and qubits.

The numerization concepts presented herein appears qualified for further investigation by the community at large, and to their attention this piece is addressed.

Reference

1. Dantzig, Tobias "Number -- The Language of Science". 4th Edition, The Free Press, 1954
2. Waismann, F., Introduction to Mathematical Thinking, Harper Torchbooks, New York, 1959.
3. Benacerraf, P., "What Numbers Could Not Be," Phil. Rev., vol. LXXIV, 1965, pp. 47–73.10.2307/2183530
4. Sich Jeffrey, " Counting and the Natural Numbers" Philosophy of Science , Volume 37 , Issue 3 , September 1970 , pp. 405 - 416. DOI: <https://doi.org/10.1086/288316>
5. Helmut Hasse, "Number Theory". Feb 2022 · Walter de Gruyter GmbH & Co KG
6. Trygve Nagell " Introduction to Number Theory" AMS Chelsea Publishing: July 2021
7. G. Samid " Negotiating Darwin's Barrier: Evolution Limits Our View of Reality, AI Breaks Through". <https://www.preprints.org/manuscript/202409.1505/v1>
8. Innovation Solution Protocol, www.InnovationSP.net.
9. G. Samid, "Metaverse Oriented Geometry" https://www.researchgate.net/publication/360084946_Metaverse_Oriented_Geometry.
10. G. Samid "Artificial Intelligence Assisted Innovation" 2023 <https://www.intechopen.com/chapters/75159>

Disclaimer/Publisher's Note: The statements, opinions and data contained in all publications are solely those of the individual author(s) and contributor(s) and not of MDPI and/or the editor(s). MDPI and/or the editor(s) disclaim responsibility for any injury to people or property resulting from any ideas, methods, instructions or products referred to in the content.