

Review

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Review

A Review of Population Count Algorithms in Binary Sequence : Through A Software Perspective

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Abstract: Counting of number of set bits as high in a binary sequence is a very important and a powerful domain question that has it's wide applications in cryptography, statistical machine learning, as well as information theory of data retrieval. With the increasing demand in software systems for tabulating the hamming weight/ Population Count / Bit count , unfortunately the classical algorithm of bitwise iteration remains widespread. The present study explores the various software-based algorithms used in the current objective like Brian-Kerningham, LookUp Table algorithm, etc. Upon performing the complexity comparison study of time as well as space complexities its concludable that Lookup table method is about 5 times more efficient than the traditional sequential search.

Keywords: algorithmic complexity; bit counting; cryptography; density count; document retrieval; efficient algorithm; hamming weight

1. Introduction

The rapid increase in computation devices following their architectural progression has witnessed the need for counting the number of Bit's High(that is 1) in a binary sequence. Initial discussions in early 1960's [1] revealed it's significance in varying fields of Computational Science and Mathematics like those in cryptography, database management, statistical machine learning. The following table reveals the various common terminologies associated with population count in binary sequence.

Table 1. Contextual Description useful in terms related to the population count.

S. No	Synonym	Contextual Description
1	Bit Count	Generic term for counting bits, often implies counting 1s
2	Population Count	Most widely used term in software/hardware literature
3	Hamming Weight	Term used in information theory and coding
4	1's Count	Number of 1s in the binary sequence
5	Set Bit Count	Bits that are 1
6	Bitwise Weight	Usud in Hardware Literatures , specially related to FPGA

Authors specifically E. El-Qawasmeh and Marilyn Mack have contributed significantly in this domain, with their papers having the following titles :

- Papers by Marilyn Mack :

(a) A bit-counting algorithm using the frequency division principle [2]

- Papers by E. El Qawasmeh :

(a) Beating the PopcountT[3]

- While both the researchers that is Marilyn Mack and E. El Qawasmeh contributed the following works together:

(a) Increasing the Efficiency of Bit-Counting[4]

(b) Organization of near matching in bit attribute matrix applied to associative access methods in information retrieval. [5]

On comparing with Hardware perspective papers of Daniel Lemire, titled the following (a) introduces

Advancement recent in this field for Faster Positional-Population Counts for AVX2, AVX-512, and ASIMD. [6]

(b) Faster Population Counts Using AVX2 Instructions [7]

These emphasizes a lot over the technicality and need to compute mathematically the Population-Count of given Binary Sequence.

2. Case Study

Algorithm 1 Population count using Naive Bitwise Count

```
def popcount_naive(NUM : int) -> int:
    COUNT = 0
    while NUM:
        COUNT += NUM & 1
        NUM = NUM » 1
    return COUNT
```

Algorithm 2 Population count using Brian Kernighan's Algorithm

```
def popcount_b_kernighan(NUM : int) -> int:
    COUNT = 0
    while NUM:
        NUM = NUM & (NUM - 1)
        COUNT += 1
    return COUNT
```

Algorithm 3 Population count using Lookup Table

```
bit_Table_256 = [0] * 256
```

```
for bit_i in range(256):
    bit_Table_256[bit_i] = (bit_i & 1) + bit_Table_256[bit_i » 1]
```

```
def popcount_LookupTable(NUM : int) -> int :
    return bit_Table_256[ NUM & 0xff ] + bit_Table_256[ (NUM » 8) & 0xff ] +
           bit_Table_256[ (NUM » 16) & 0xff ] + bit_Table_256[ (NUM » 24) & 0xff ]
```

Algorithm 4 Population count using CSA - Carry Save Adder

from typing import List

```
def csa(a: int, b: int, c: int) -> (int, int):
    total = a ^ b ^ c
    carry = (a & b) | (b & c) | (a & c)
    return carry, total
```

```
def popcount(x: int) -> int:
```

```
    x = x - ((x >> 1) & 0x5555555555555555)
    x = (x & 0x3333333333333333) + ((x >> 2) & 0x3333333333333333)
    x = (x + (x >> 4)) & 0x0F0F0F0F0F0F0F0F
    x = x + (x >> 8)
    x = x + (x >> 16)
    x = x + (x >> 32)
    return x & 0x7F
```

```
def popcount_csa(data: List[int]) -> int:
```

```
    ones = twos = fours = eights = 0
    i = 0
    n = len(data)
    while i + 7 < n:
        t1, o1 = csa(data[i], data[i+1], data[i+2])
        t2, o2 = csa(data[i+3], data[i+4], data[i+5])
        f1, t3 = csa(t1, t2, data[i+6])
        f2, ones = csa(o1, o2, data[i+7])
        fours, twos = csa(f1, f2, twos)
        eights, fours = csa(eights, fours, 0)
        i += 8
```

```
    TOTAL = (
        8 * popcount(eights) +
        4 * popcount(fours) +
        2 * popcount(twos) +
        1 * popcount(ones) )
```

```
    while i < n:
        TOTAL += popcount(data[i])
        i += 1
    return TOTAL
```

Algorithm 5 Population count using VTPC method

```

def popcount_VTPC(n : int) -> int :
    mask_size = 4
    mask = [0,1,1,2]
    m = 0
    i = 0
    d =  $\lceil \log_{\text{mask\_size}}(n) \rceil$ 
    e = mask_size
    r = n
    while d  $\geq$  0 :
        i =  $\lfloor (r/e) \times \text{mask\_size} \rfloor$ 
        r = r -  $(\frac{e}{\text{mask\_size}} \times i)$ 
        e =  $\frac{e}{\text{mask\_size}}$ 
        m = m + mask[i]
        d = d - 1
    return m

```

Algorithm 6 Population count using Divide and Conquer adder

```

def popcount_divconq(NUM : int) -> int :
    NUM = (NUM & 0x55555555) + ((NUM  $\gg$  1) & 0x55555555)
    NUM = (NUM & 0x33333333) + ((NUM  $\gg$  2) & 0x33333333)
    NUM = (NUM + (NUM  $\gg$  4)) & 0x0F0F0F0F
    NUM = (NUM * 0x01010101)  $\gg$  24
    return NUM

```

Algorithm 7 Population count using RF method

```

def countRF(N):
    TEMP = N
    COIN = 2
    prev = 1

    COUNTER = 0

    while TEMP :
        ss = (N+1) // COIN
        v1 = (N+1) % COIN
        if v1 <= prev:
            REM = 0
        else:
            REM = (N+1) % prev
        COUNTER += REM + ss*prev
        prev = prev << 1
        COIN = COIN << 1
        TEMP = TEMP  $\gg$  1

    return COUNTER

def RF(N):
    return countRF(N) - countRF(N-1)

```

3. Complexity Comparison

Under this section we discuss and show the complexity analysis of various algorithms, by taking LOG_{10} OF (N) : Input number, as this value ranges from [1,1000000], it's imperative to transform it into

a scale that's quite implementable on the graph of Time to compute vs Input Size. Hence we applied the following transformations :

Table 2. The transformation analysis of plotting functions over the graph from old v/s new perspective , where N : number till we checked Natural numbers for the several algorithms, the results T : Time taken to compute these results (ms).

S.No	Old parameter	Transformed parameter
1	N	$\log_{10} N$
2	T	$\ln T$

Table 3. An outline in comparison of several software algorithmic implementations of techniques in computing the population count for binary sequences, where M : Number of bits used to implement the input Number N : Input Number L : Number of elements in addressing each element of Lookup Table S : Size of Mask Vector.

S. No	Method	Average Time Complexity	Space Complexity
1	Naive Approach	$O(M)$	$O(1)$
2	Brian Kerningham	$O(\log_2 (N))$	$O(1)$
3	Lookup Table	$O(M / L)$	$O(256)$
4	CSA	$O(\log_2 (M))$	$O(M)$
5	VTPC	$O(\log_5 (N))$	$O(S)$
6	Divide and Conquer	$O(\log_2 (M))$	$O(1)$
7	RF Algorithm	$O(\log_2 N)$	$O(1)$

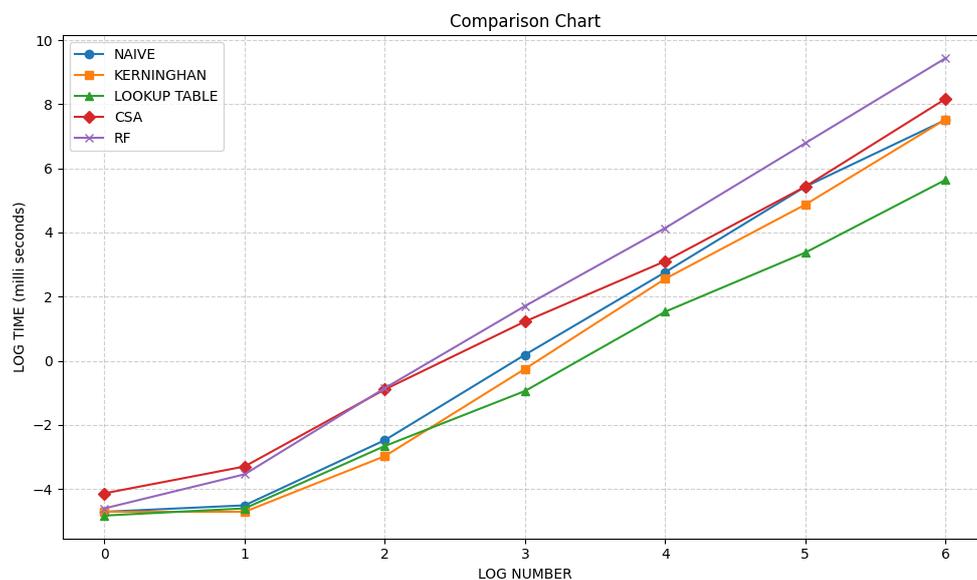


Figure 1. The plot of natural LOG T : Time taken against \log_{10} (NUMBER) of Natural Numbers till count of population count is taken.

Table 4. Time Taken to Compute Bit's population count till given Number - NUM by various Algorithms, with time taken to compute is in milli seconds (ms).

NUM	NAIVE	KERNINGHAN	LOOKUP	CSA	RF
1	0.009	0.009	0.008	0.016	0.01
10	0.011	0.009	0.01	0.037	0.029
100	0.084	0.051	0.07	0.409	0.426
1000	1.206	0.776	0.389	3.403	5.509
10000	15.816	12.887	4.616	22.309	62.769
100000	227.941	130.344	29.188	228.655	886.574
1000000	1823.992	1858.631	281.325	3518.617	12526.57

4. Conclusion

Through the current work, we presented multiple software-oriented algorithms that are pivotal in computing the population count of ones in a binary sequence. The results were developed as a result of and are dependent upon the programming language, compiler/interpreter version, also the server's memory and priority pooling allocation for the google colab notebook, that was used to run and implement these algorithms.

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